## CARNEGIE MELLON UNIVERSITY DEPARTMENT OF COMPUTER SCIENCE 15-445/645 – DATABASE SYSTEMS (FALL 2017) PROF. ANDY PAVLO

## Homework 4 (by Sivaprasad Sudhir) Due: Wednesday Oct 11, 2017 @ 11:59pm

## **IMPORTANT:**

- Upload this PDF with your answers to Gradescope by 11:59pm on Wednesday Oct 11, 2017.
- **Plagiarism**: Homework may be discussed with other students, but all homework is to be completed **individually**.
- You have to use this PDF for all of your answers.

For your information:

- Graded out of **100** points; **2** questions total
- Rough time estimate:  $\approx 1 2$  hours (0.5 1 hours for each question)

Revision : 2017/10/08 22:58

Question	Points	Score
Sorting	40	
Join Algorithms	60	
Total:	100	

- - (a) [10 points] What is the smallest number of buffers B, that can sort the file with N = 1,000,000 pages, in 2 passes?
    □ 32 □ 33 □ 34 □ 99 □ 100 □ 101 □ 102 □ 999 □ 1,000 □ 1,000 □ 1,000,001
  - (b) [10 points] What is the smallest number of buffers B, that can sort the file with N = 1,000,000 pages, in 3 passes?
    □ 32 □ 33 □ 34 □ 99 □ 100 □ 101 □ 102 □ 999 □ 1,000 □ 1,001 □ 1,000,000 □ 1,000,001
  - (c) [10 points] How many passes are needed to sort the file with N = 1,000,000 pages with 6 buffers?

 $\Box 7 \quad \Box 8 \quad \Box 9 \quad \Box 10 \quad \Box 11$ 

(d) **[10 points]** What is the total I/O cost to sort the file with N = 1,000,000 pages with 6 buffers?

 $\Box 14,000,000 \Box 8,000,000 \Box 18,000,000 \Box 10,000,000 \Box 22,000,000$ 

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Consider relations R(x, y) and S(x, z) to be joined on the common attribute x. Assume that there are no indexes.

- There are B = 30 pages in the buffer
- Table R spans M = 1800 pages with 50 tuples per page
- Table S spans N = 500 pages with 100 tuples per page

What are the I/O costs for the following joins?

- Assume the simplest cost model, where pages are read and written one at a time
- Assume that you will need one buffer block to hold the evolving output block and one input block to hold the current input block of the inner relation
- Ignore the cost of the final writing of the results
- (a) **[10 points]** Block nested loop join with R as the outer relation and S as the inner relation  $\Box$  30,000  $\Box$  31,800  $\Box$  32,900  $\Box$  33,300  $\Box$  34,300
- (b) **[5 points]** Block nested loop join with S as the outer relation and R as the inner relation  $\Box$  29,300  $\Box$  31,100  $\Box$  31,800  $\Box$  32,900  $\Box$  34,300
- (c) Hash join with S as the outer relation and R as the inner relation (Ignore recursive partitioning and partially filled blocks)
  - i. [5 points]
     Cost of partition phase

     □
     2,300
     □
     4,600
     □
     6,900
     □
     3,600
     □
     1,000

     ii. [5 points]
     Cost of probing phase
     □
     2,300
     □
     4,600
     □
     6,900
     □
     3,600
     □
     1,000
- (d) Sort-merge join
  - i. [**10 points**] Cost of sorting R □ 8966 □ 7934 □ 6578 □ 1828 □ 2204
  - ii. **[5 points]** Cost of sorting S □ 8966 □ 7934 □ 6578 □ 1828 □ 2204
  - iii. **[10 points]** Cost of merge assuming no duplicates in the join attribute  $\Box$  2,300  $\Box$  4,600  $\Box$  6,900  $\Box$  154  $\Box$  77
  - iv. **[10 points]** Cost of merge in the worst case □ 6,900 □ 31,800 □ 33,300 □ 900,000 □ 1,800,000