Carnegie Mellon University

Distributed OLTP Databases (Part II)

Lecture #23

Database Systems
15-445/15-645
Fall 2018

AP
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Computer Science
Carnegie Mellon Univ.

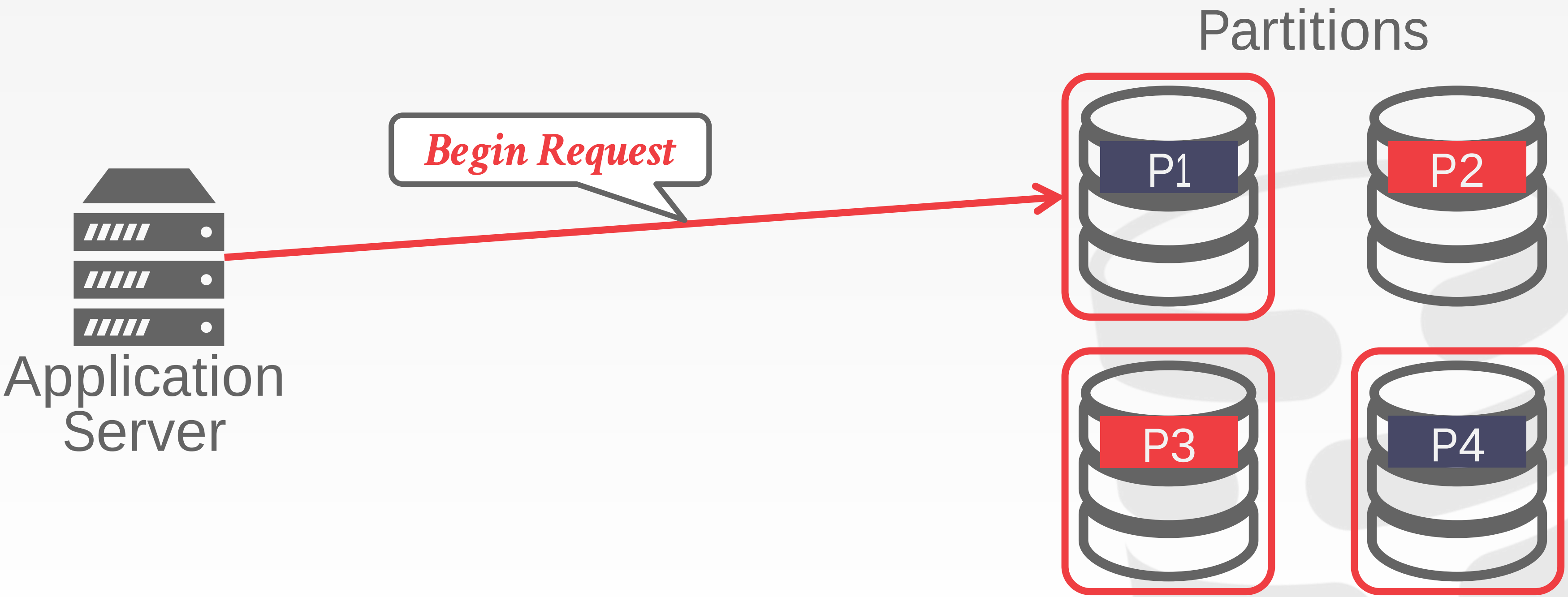
LAST CLASS

System Architectures
→ Shared-Memory, Shared-Disk, Shared-Nothing

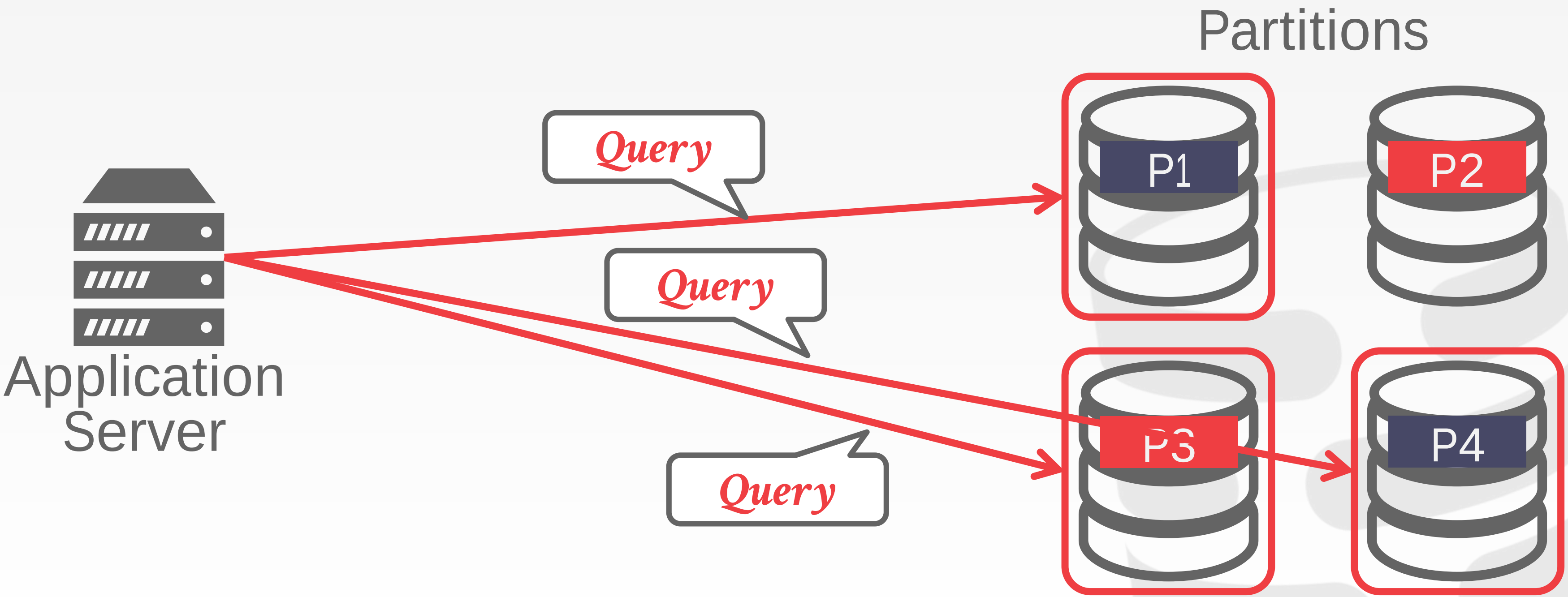
Partitioning/Sharding
→ Hash, Range, Round Robin

Transaction Coordination
→ Centralized vs. Decentralized

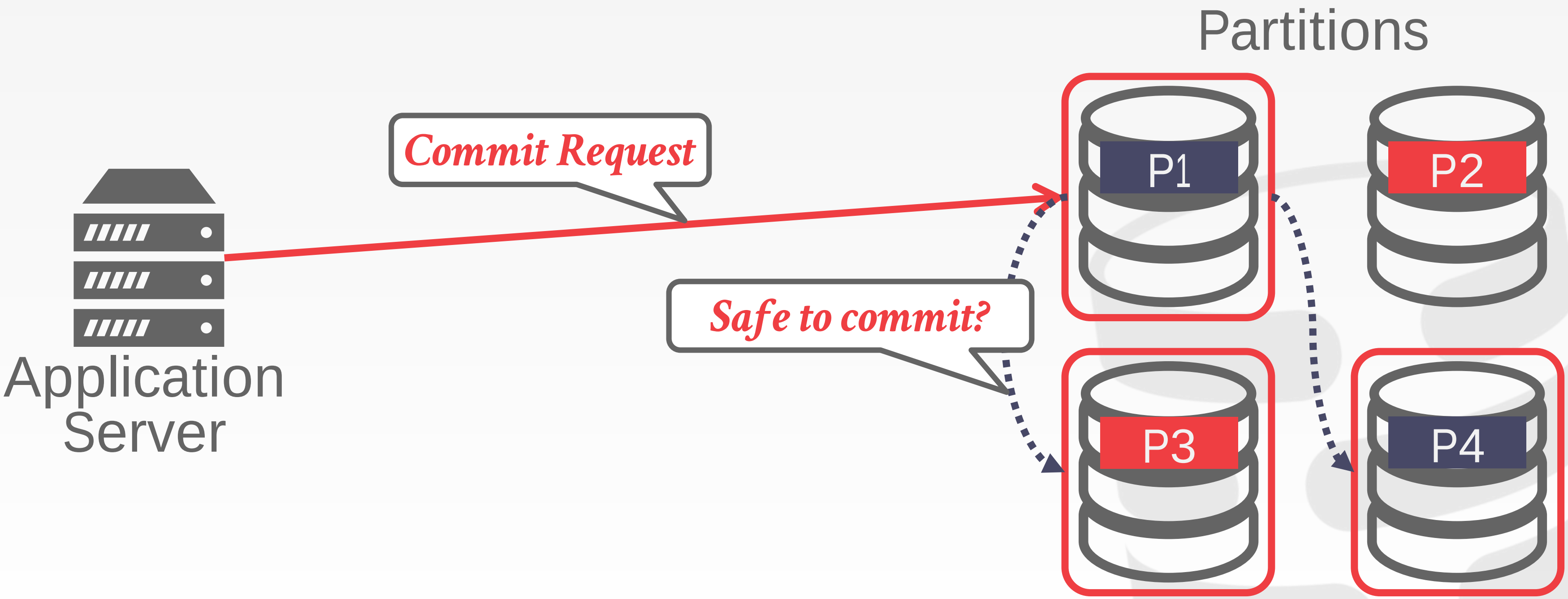
DECENTRALIZED COORDINATOR



DECENTRALIZED COORDINATOR



DECENTRALIZED COORDINATOR



OBSERVATION

We have not discussed how to ensure that all nodes agree to commit a txn and then to make sure it does commit if we decide that it should.

→ What happens if a node fails?
→ What happens if our messages show up late?
→ What happens if we don't wait for every node to agree?

TODAY'S AGENDA

Atomic Commit Protocols
Replication
Consistency Issues (CAP)
Federated Databases

ATOMIC COMMIT PROTOCOL

When a multi-node txn finishes, the DBMS needs to ask all of the nodes involved whether it is safe to commit.

Examples:

→ Two-Phase Commit
→ Three-Phase Commit (not used)
→ Paxos
→ Raft
→ ZAB (Apache Zookeeper)
→ Viewstamped Replication

TWO-PHASE COMMIT (SUCCESS)

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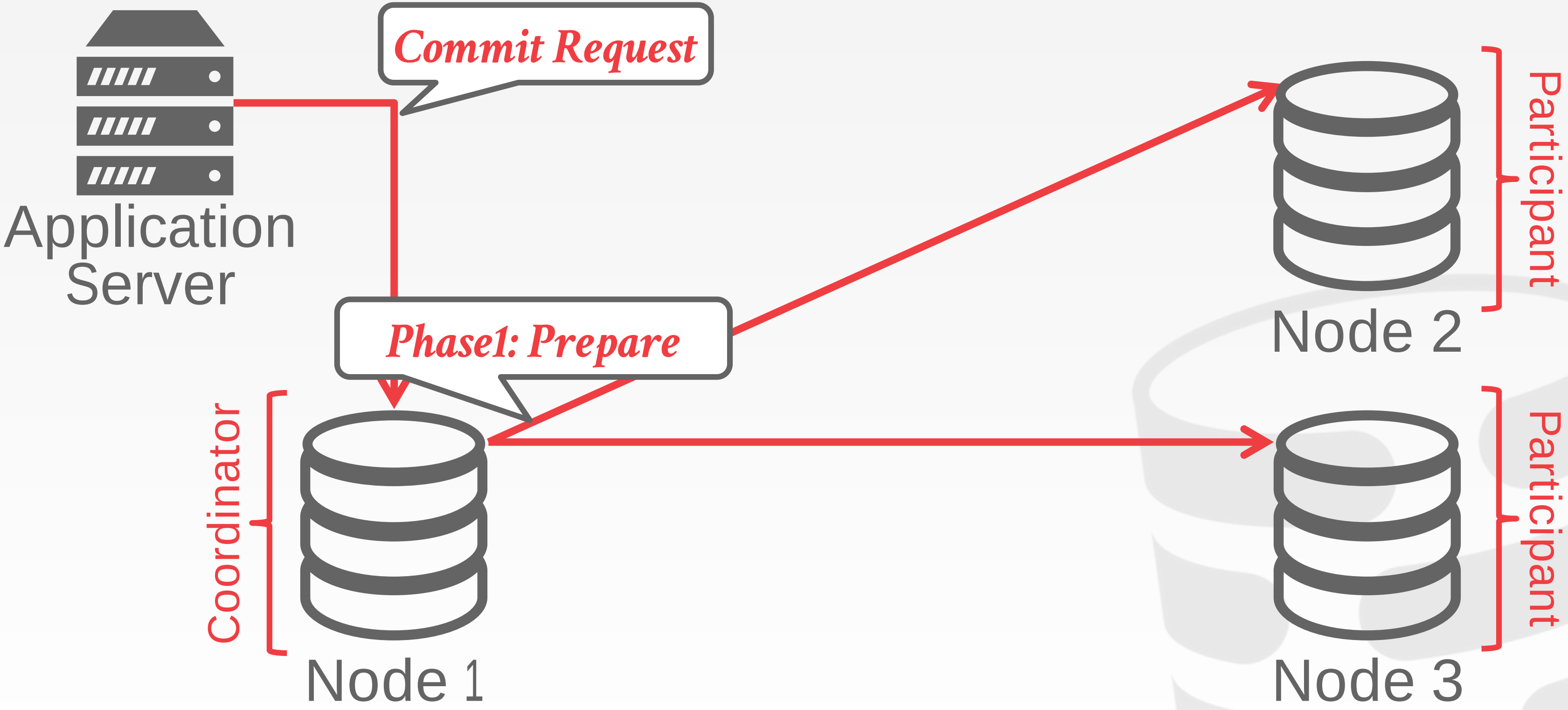

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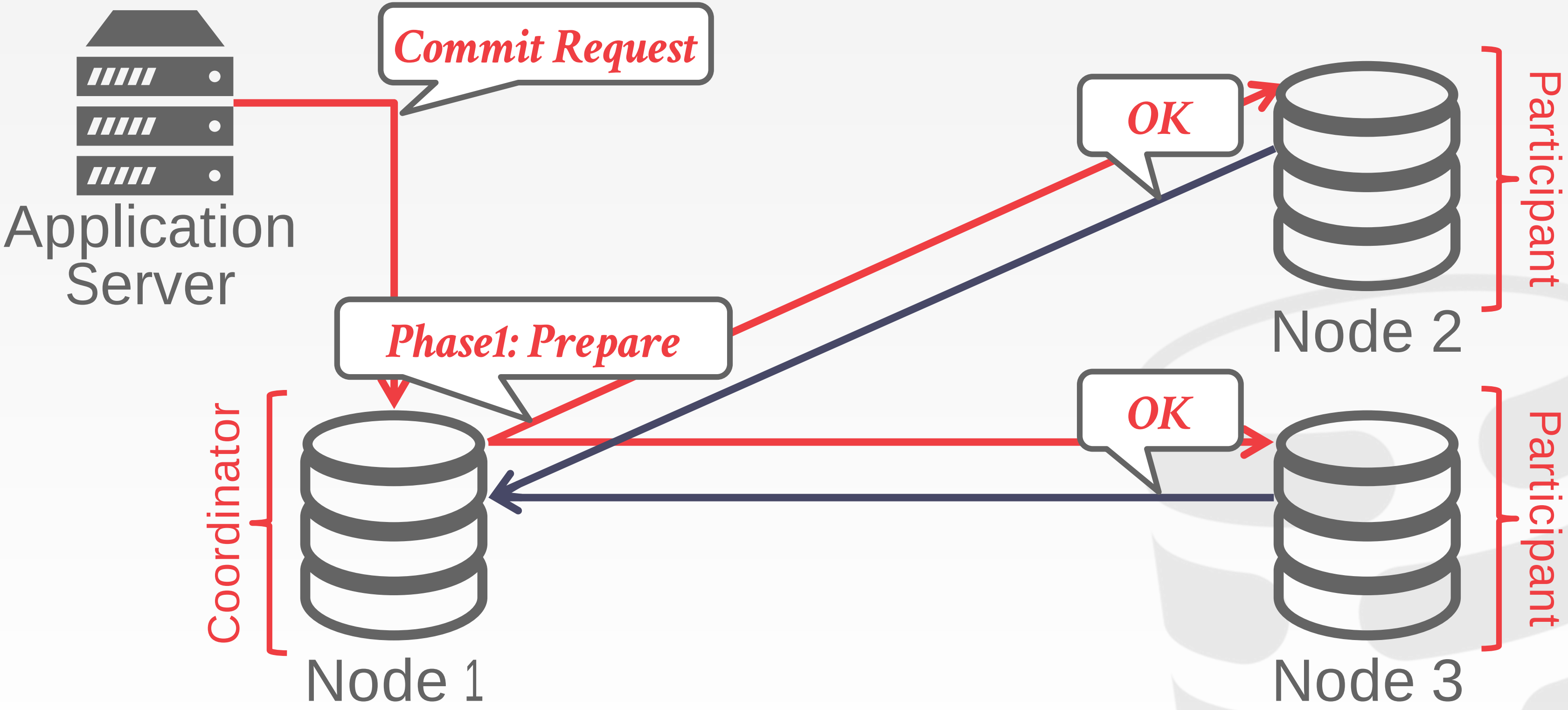

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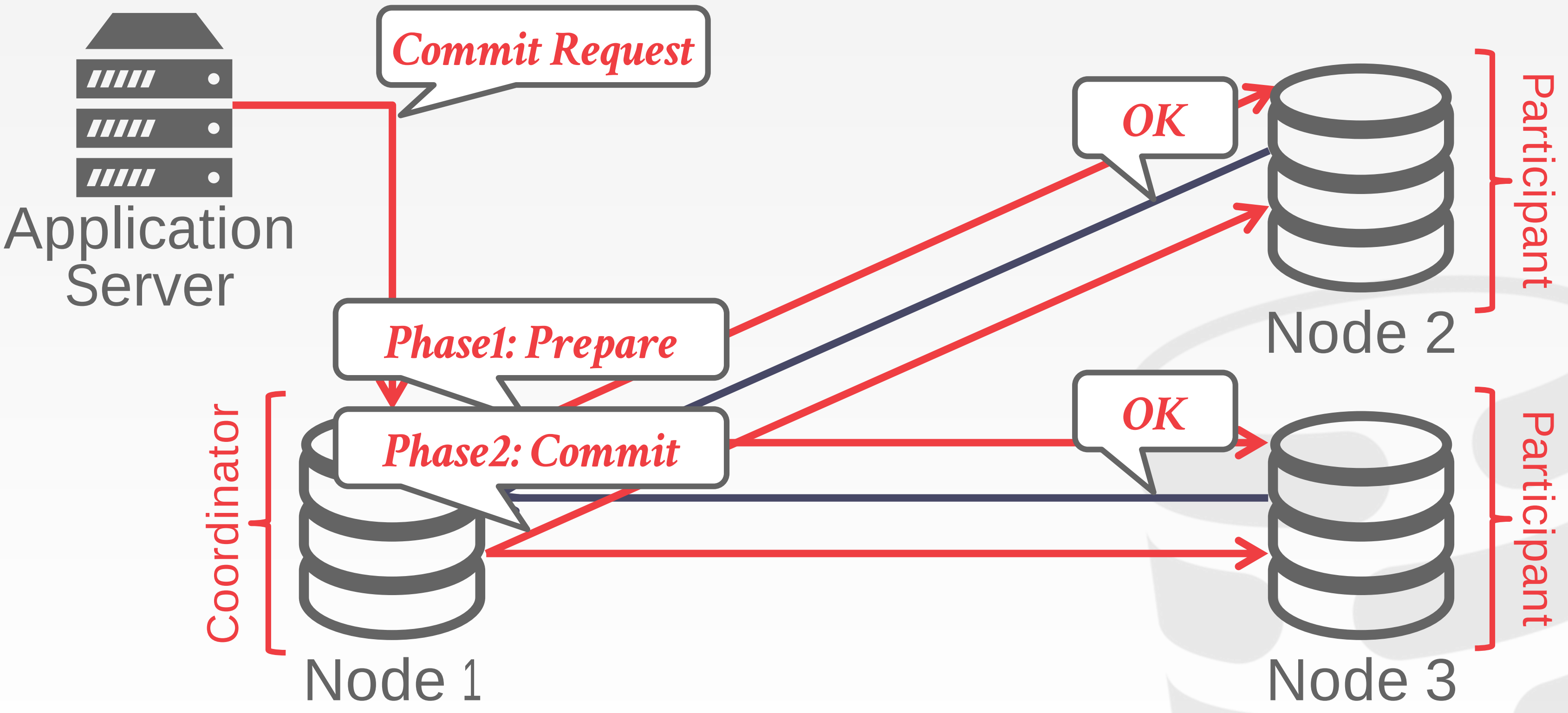

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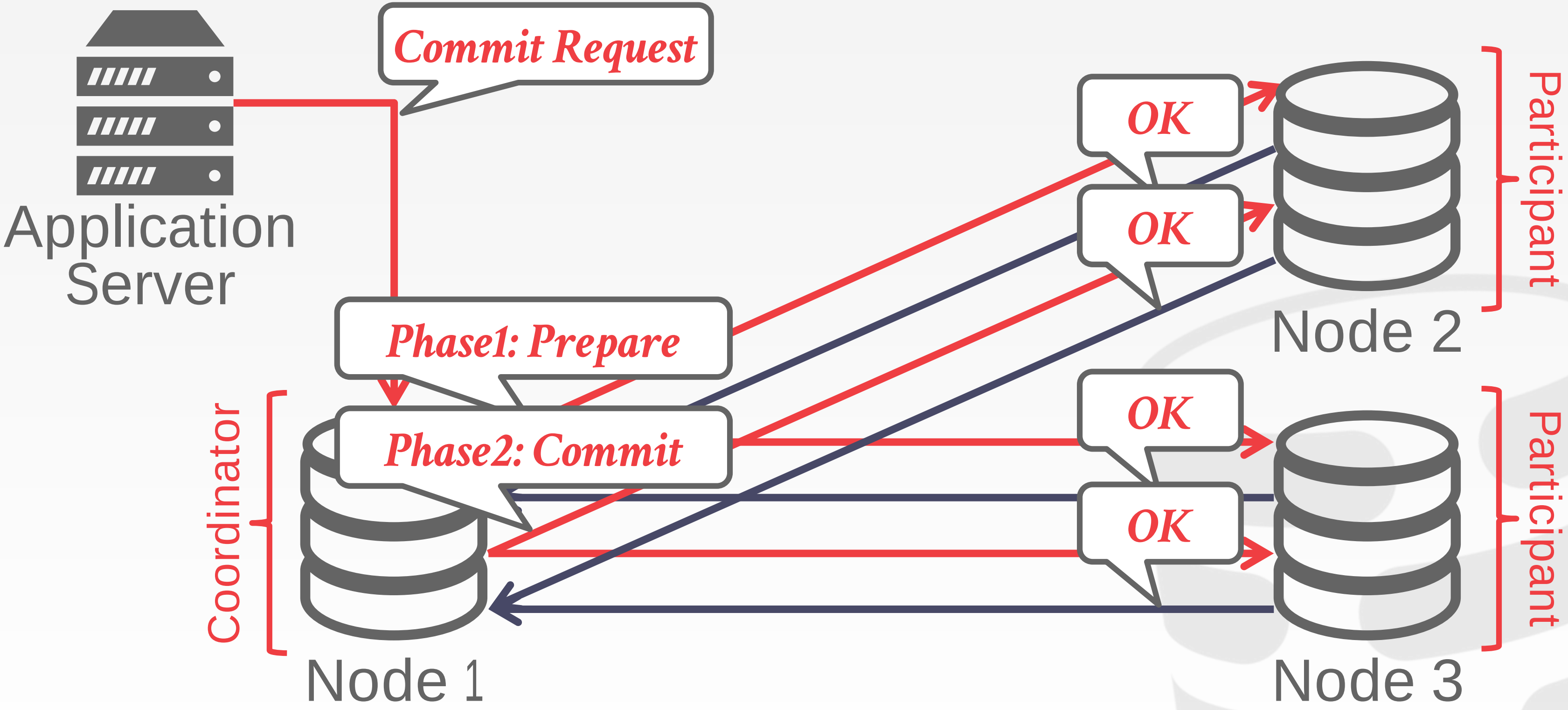
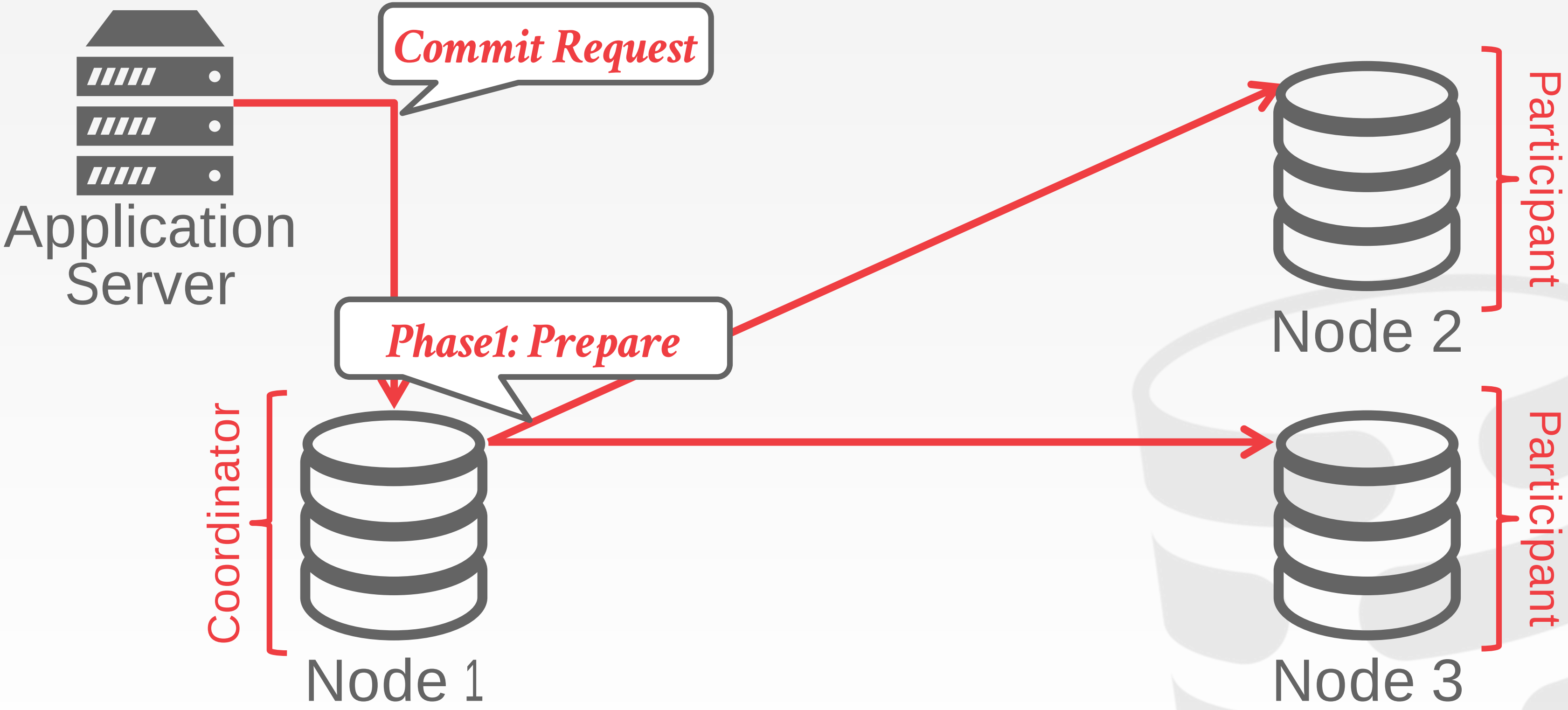


TWO-PHASE COMMIT (SUCCESS)

TWO-PHASE COMMIT (ABORT)

TWO-PHASE COMMIT (ABORT)



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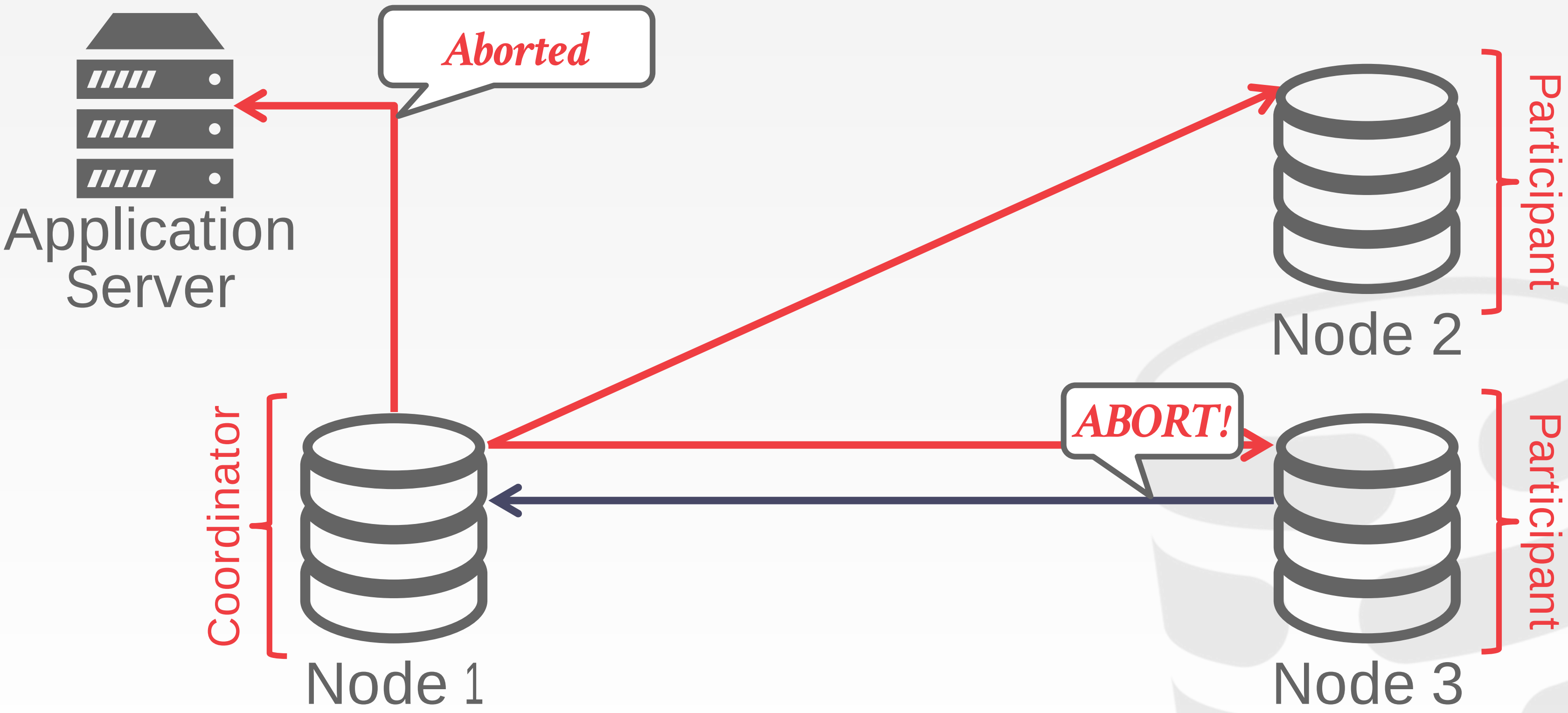

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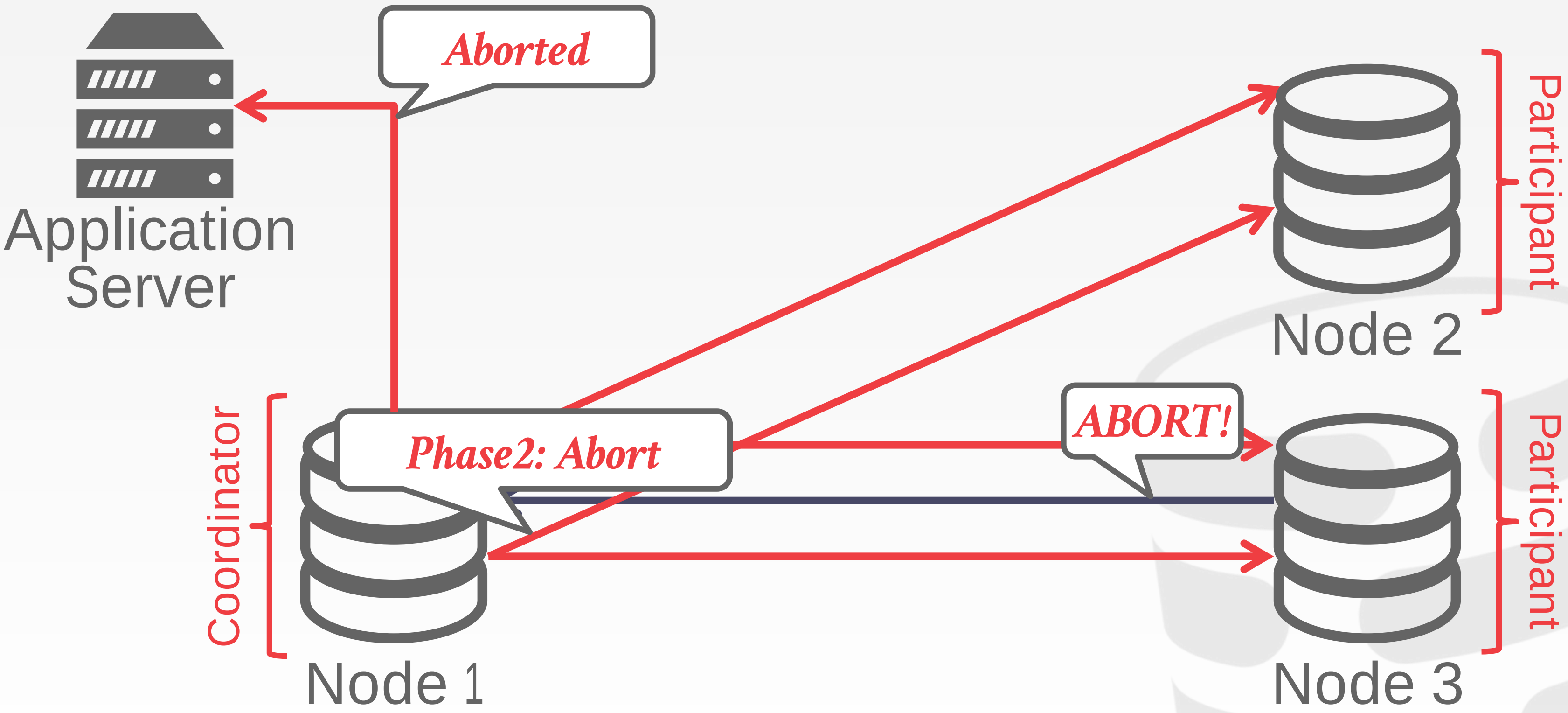
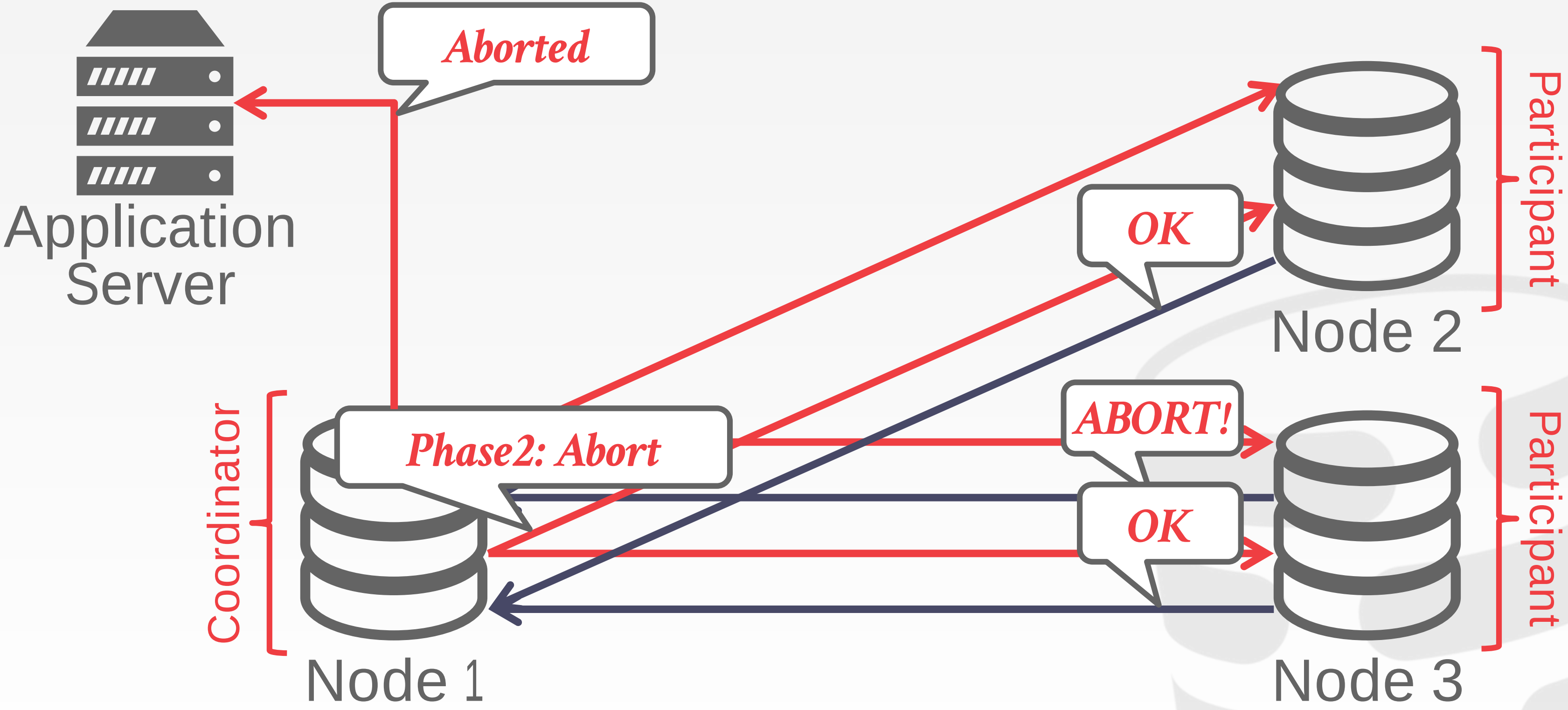


TWO-PHASE COMMIT (ABORT)



2PC OPTIMIZATIONS

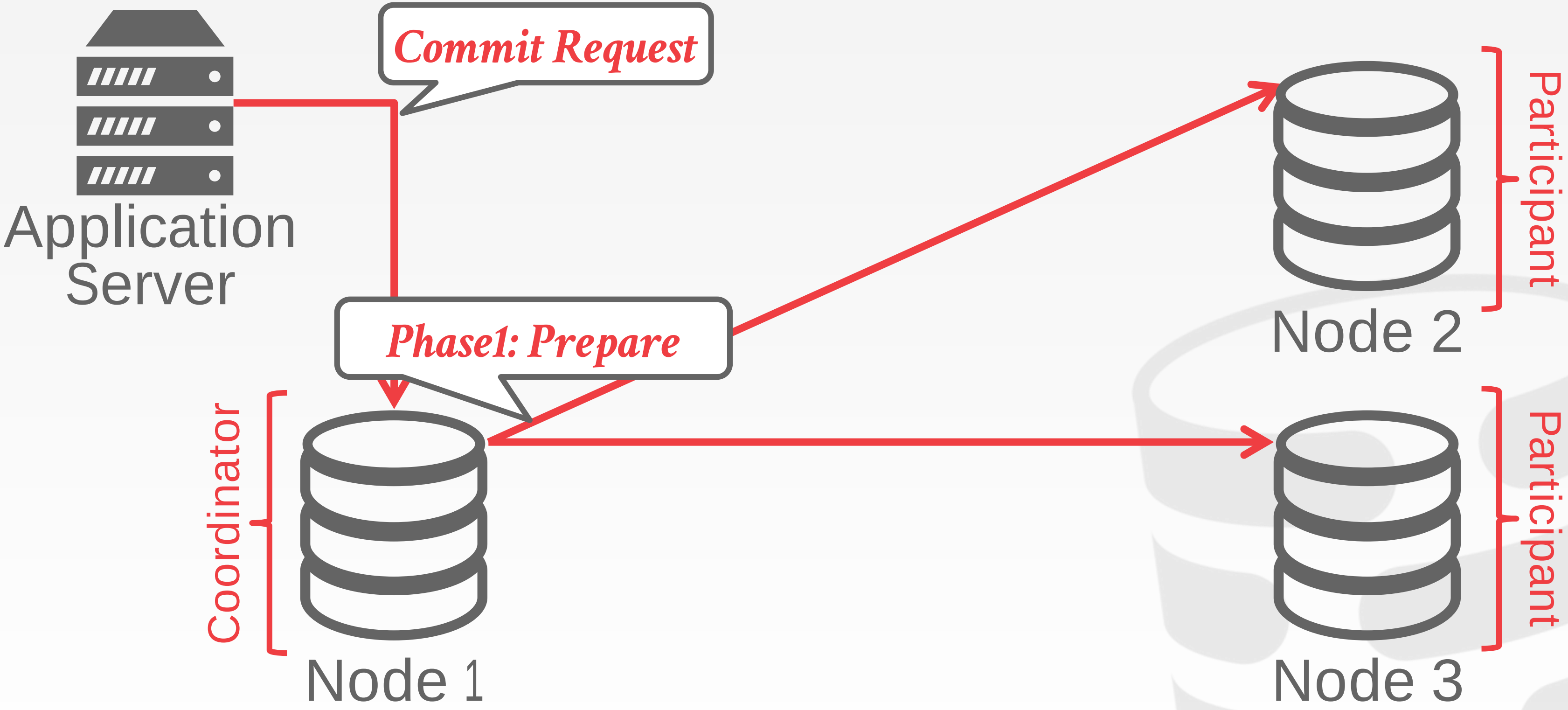
Early Prepare Voting

→ If you send a query to a remote node that you know will be the last one you execute there, then that node will also return their vote for the prepare phase with the query result.

Early Acknowledgement After Prepare

→ If all nodes vote to commit a txn, the coordinator can send the client an acknowledgement that their txn was successful before the commit phase finishes.

EARLY ACKNOWLEDGEMENT



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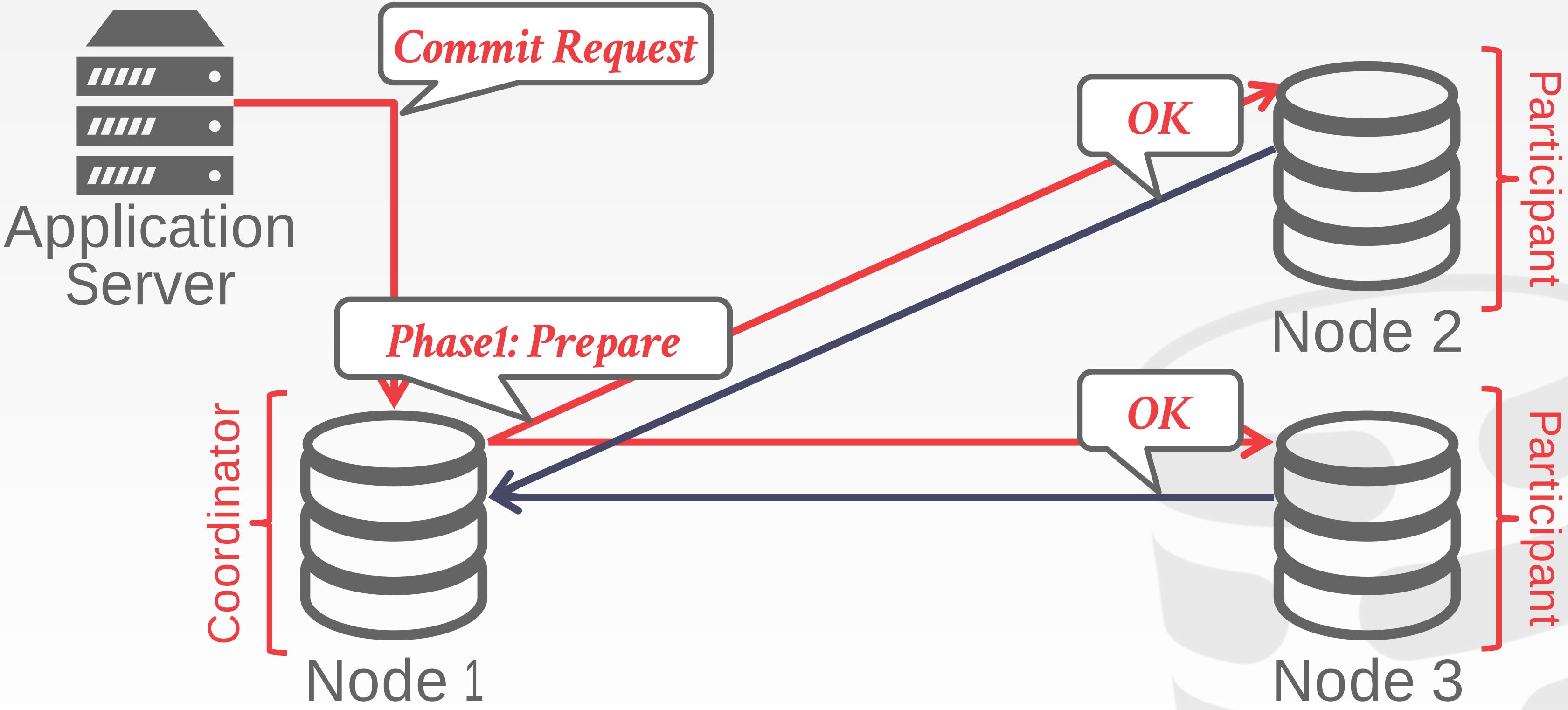

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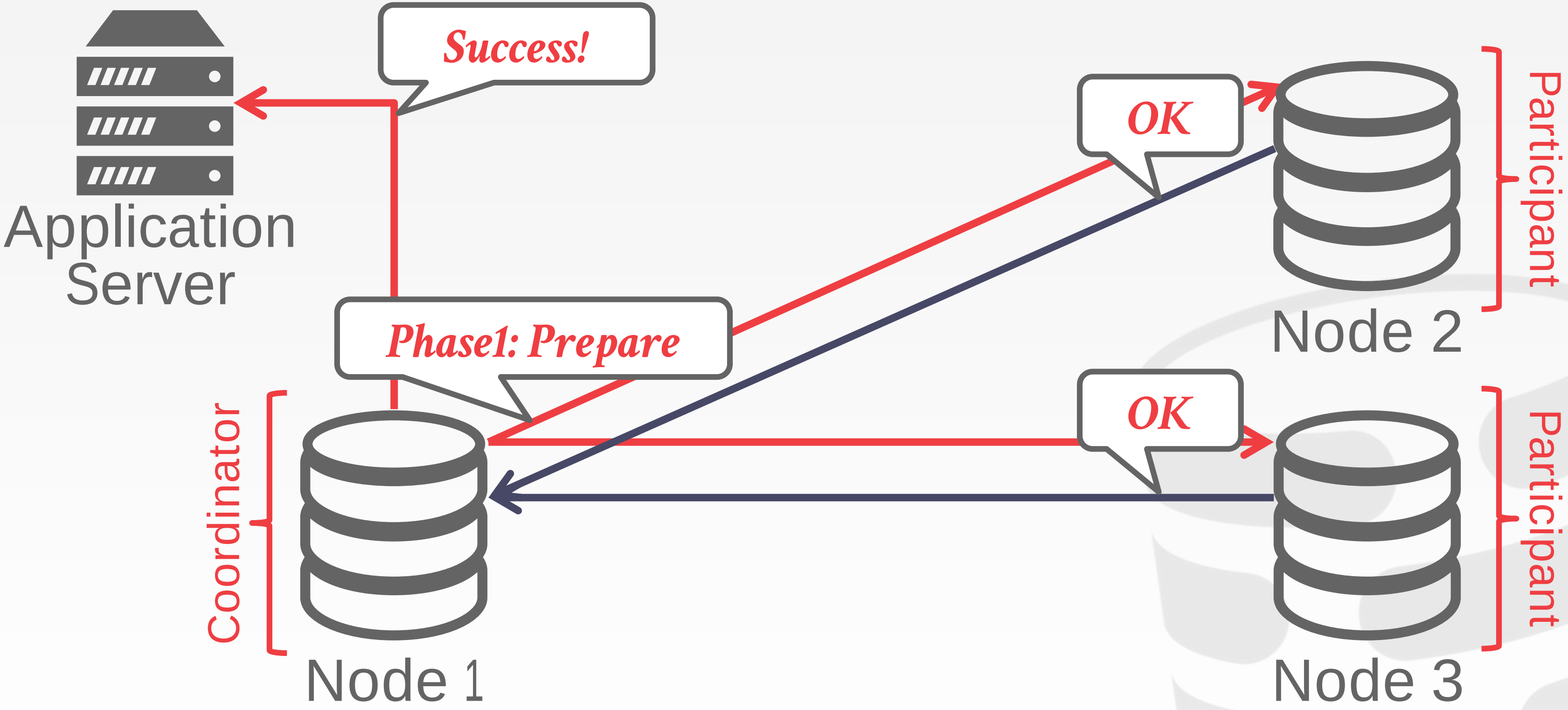

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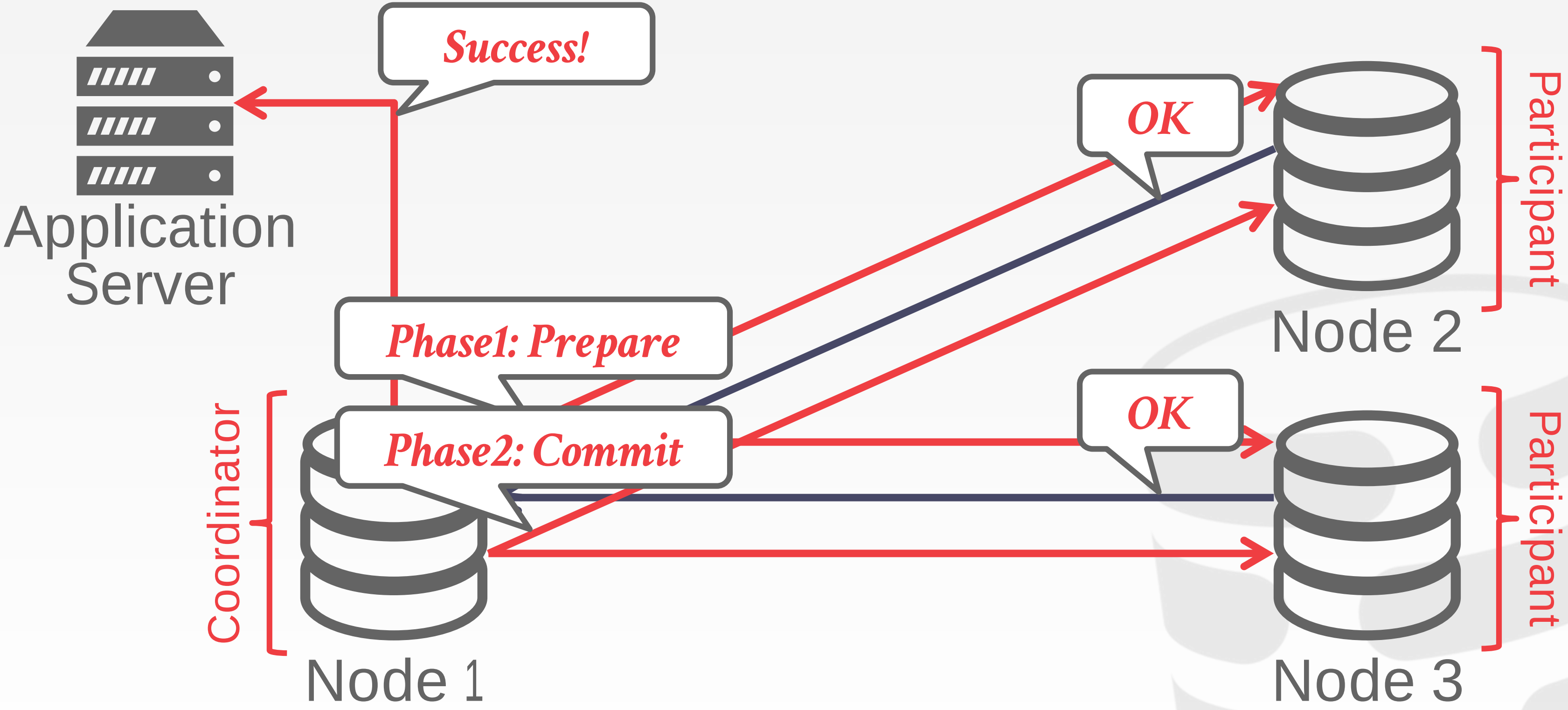

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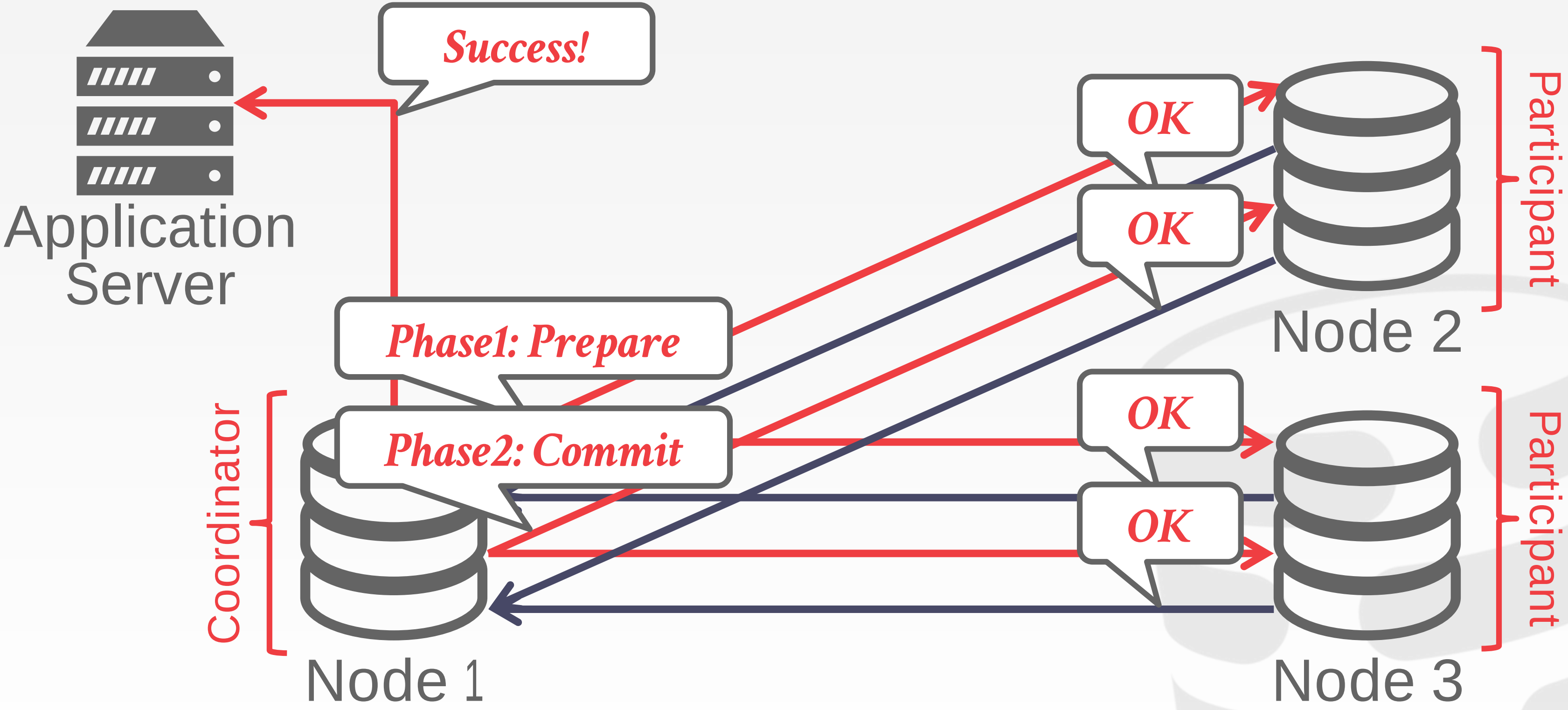


TWO-PHASE COMMIT

Each node has to record the outcome of each phase in a stable storage log.

What happens if coordinator crashes?

→ Participants have to decide what to do.

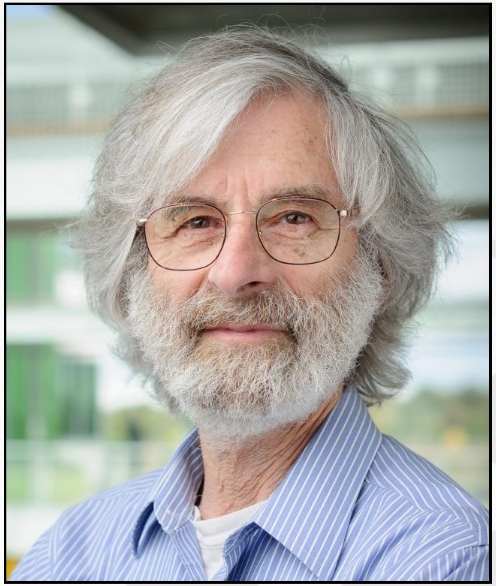
What happens if participant crashes?

→ Coordinator assumes that it responded with an abort if it hasn't sent an acknowledgement yet.

PAXOS

Consensus protocol where a coordinator proposes an outcome (e.g., commit or abort) and then the participants vote on whether that outcome should succeed.

Does not block if a majority of participants are available and has provably minimal message delays in the best case.



Lamport

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The Part-Time Parliament

LESLIE LAMPORT
Digital Equipment Corporation

Recent archaeological discoveries on the island of Paxos reveal that the parliament functioned despite the peripatetic propensity of its part-time legislators. The legislators maintained consistent copies of the parliamentary record, despite their frequent forays from the chamber and the forgetfulness of their messengers. The Paxon parliament's protocol provides a new way of implementing the state-machine approach to the design of distributed systems.

Categories and Subject Descriptors: C2.4 [**Computer-Communications Networks**]: Distributed Systems—*Network operating systems*; D4.5 [**Operating Systems**]: Reliability—*Fault-tolerance*; J.1 [**Administrative Data Processing**]: Government

General Terms: Design, Reliability

Additional Key Words and Phrases: State machines, three-phase commit, voting

This submission was recently discovered behind a filing cabinet in the *TOCS* editorial office. Despite its age, the editor-in-chief felt that it was worth publishing. Because the author is currently doing field work in the Greek isles and cannot be reached, I was asked to prepare it for publication.

The author appears to be an archeologist with only a passing interest in computer science. This is unfortunate; even though the obscure ancient Paxon civilization he describes is of little interest to most computer scientists, its legislative system is an excellent model for how to implement a distributed computer system in an asynchronous environment. Indeed, some of the refinements the Paxons made to their protocol appear to be unknown in the systems literature.

The author does give a brief discussion of the Paxon Parliament's relevance to distributed computing in Section 4. Computer scientists will probably want to read that section first. Even before that, they might want to read the explanation of the algorithm for computer scientists by Lampson [1996]. The algorithm is also described more formally by De Prisco et al. [1997]. I have added further comments on the relation between the ancient protocols and more recent work at the end of Section 4.

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University of California, San Diego

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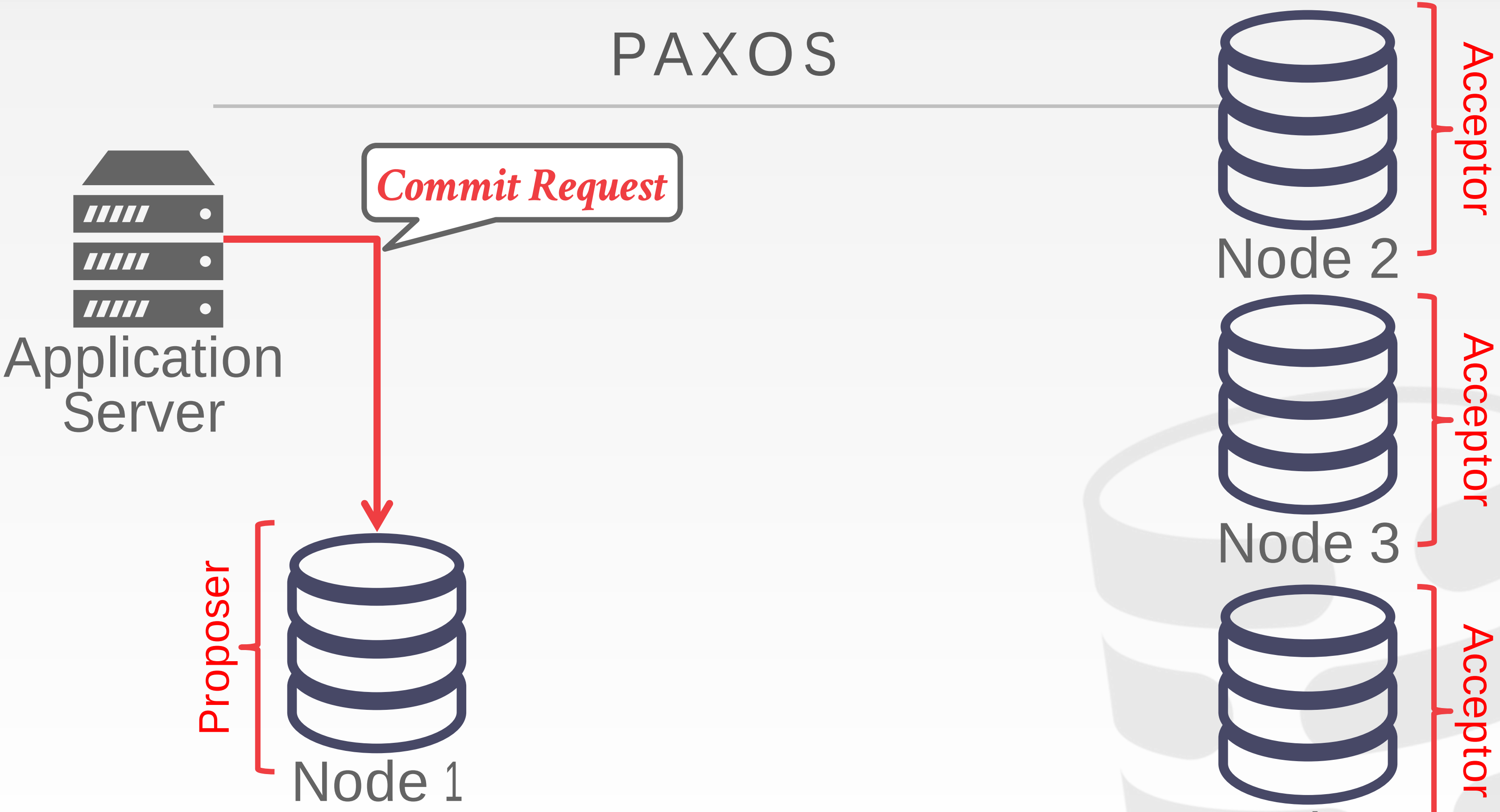
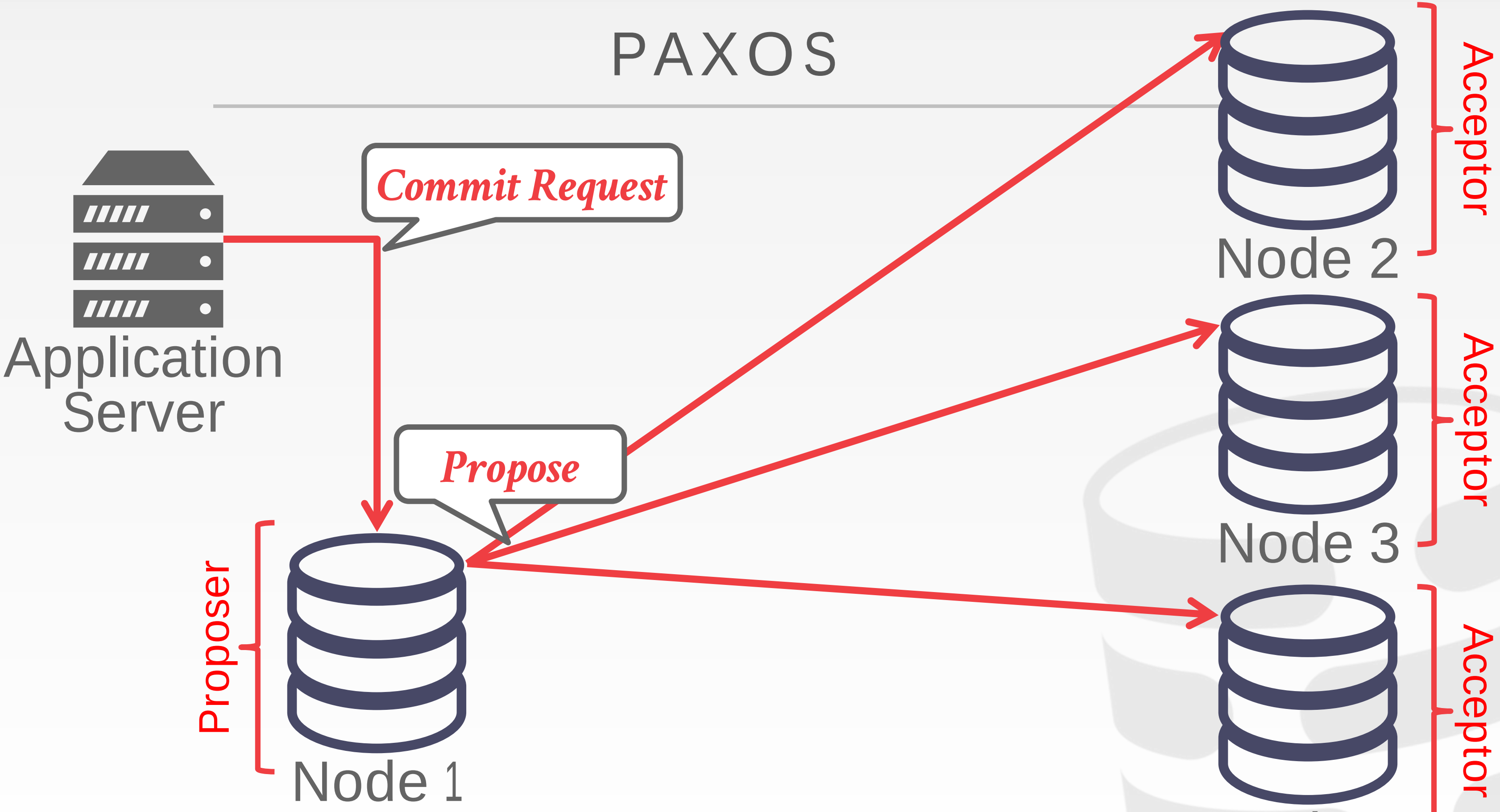
PAXOS


Commit Request
Application Server
Proposer
Node 1
Acceptor
Node 2
Acceptor
Node 3
Acceptor
Node 4

PAXOS



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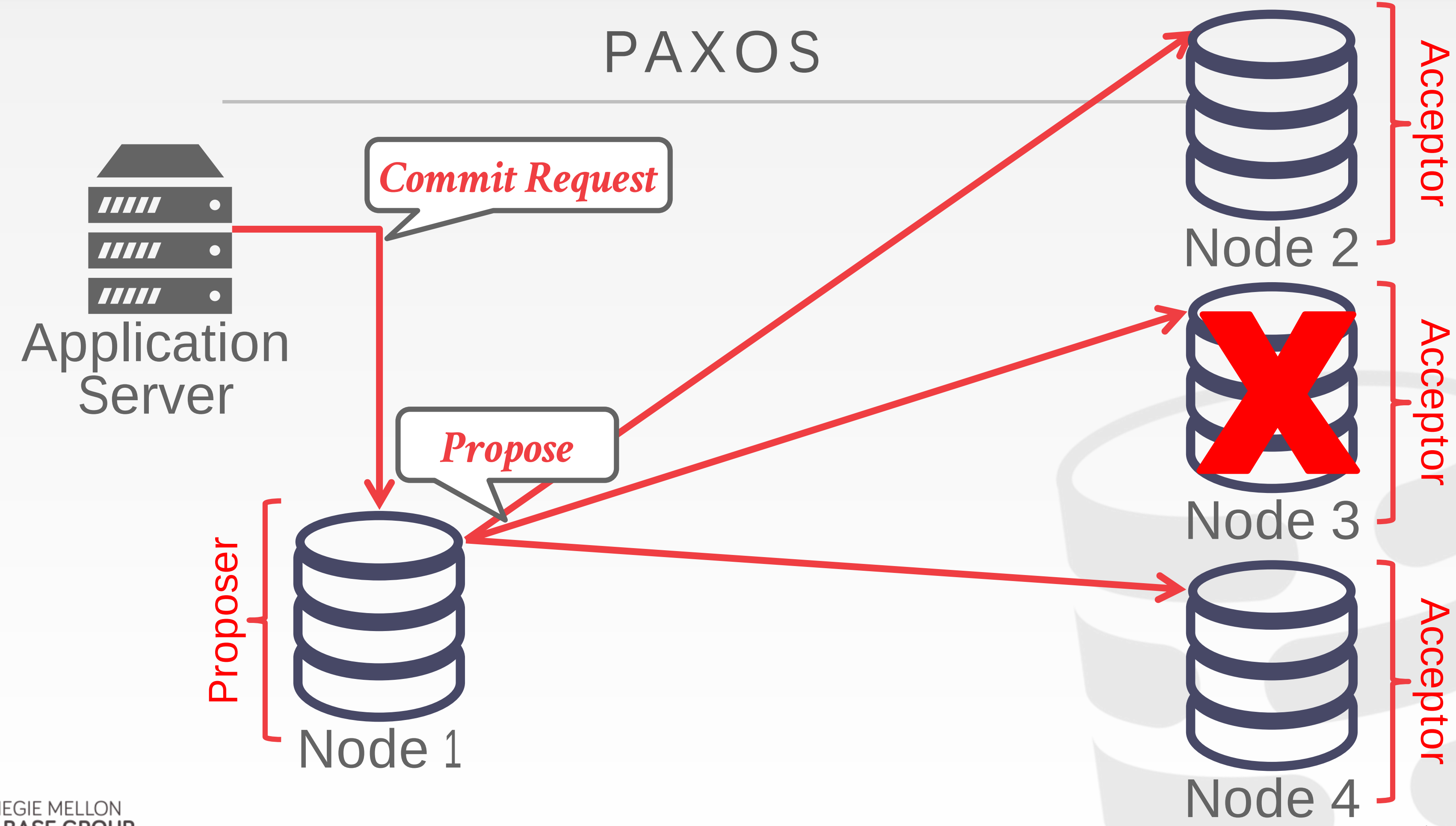
Application Server
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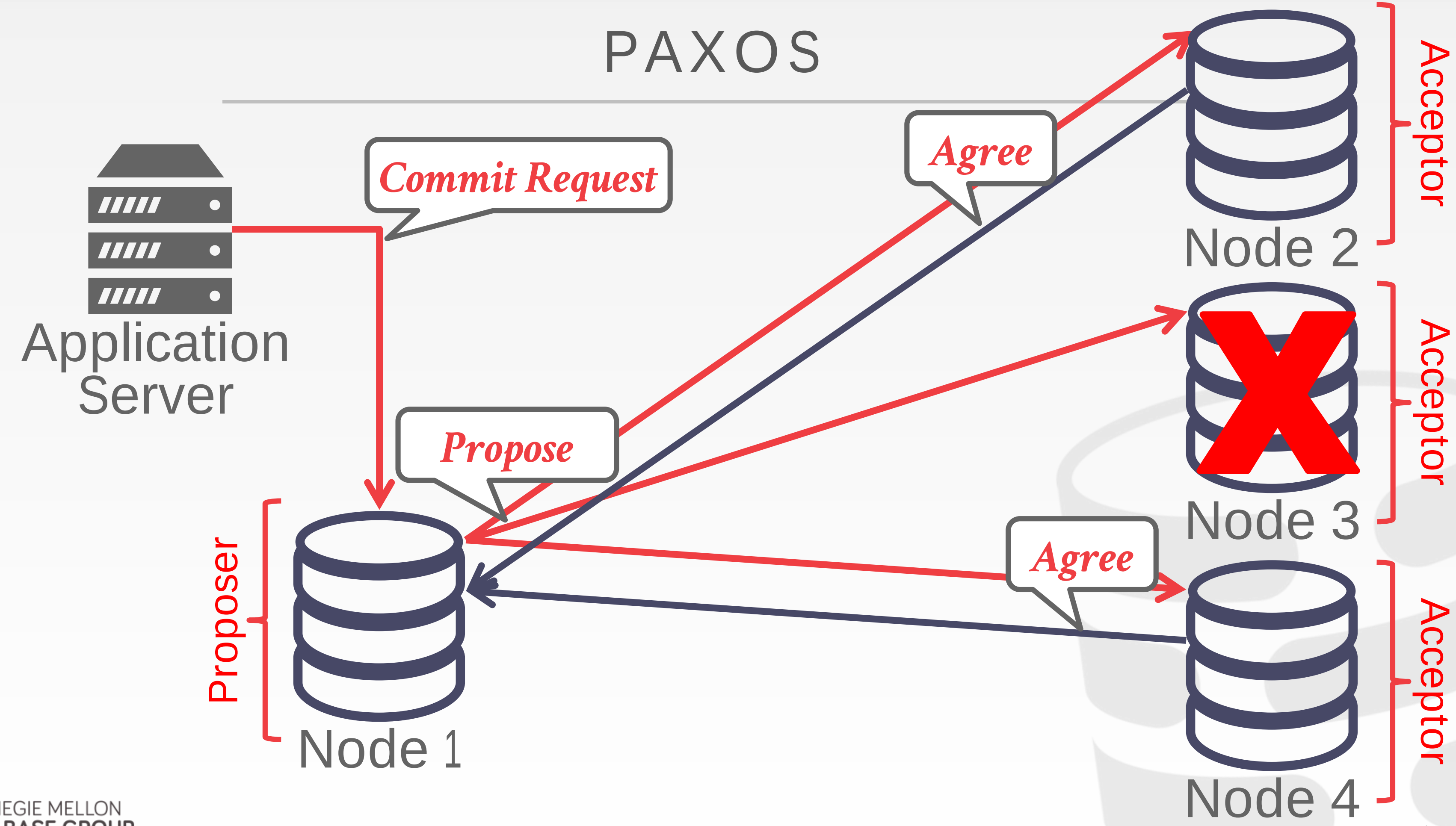
PAXOS


Application Server
Commit Request
Propose
Agree
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Node 1
Node 2
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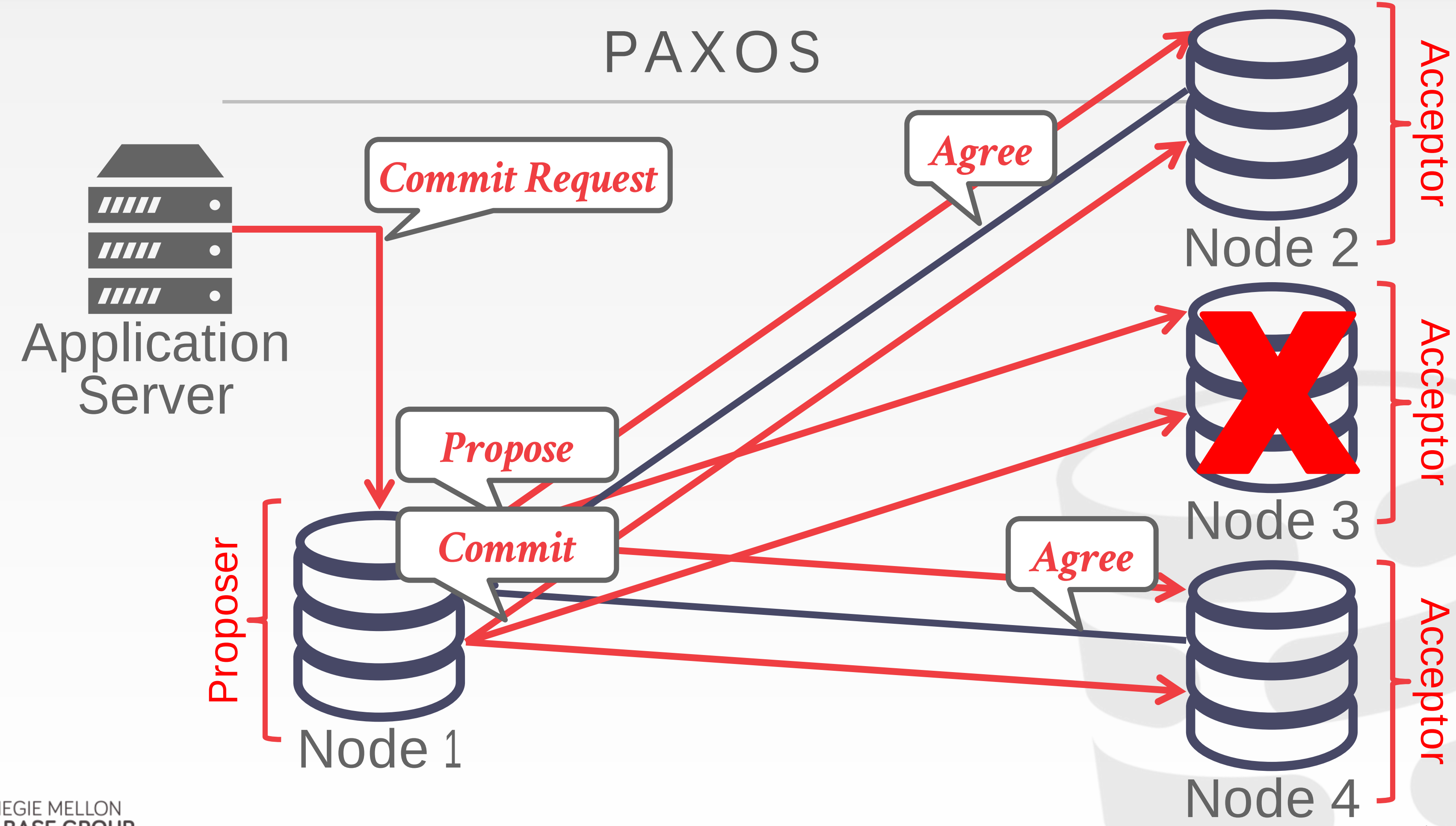
PAXOS


Application Server
Commit Request
Propose
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Proposer
Node 1
Node 2
Node 3
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Acceptor
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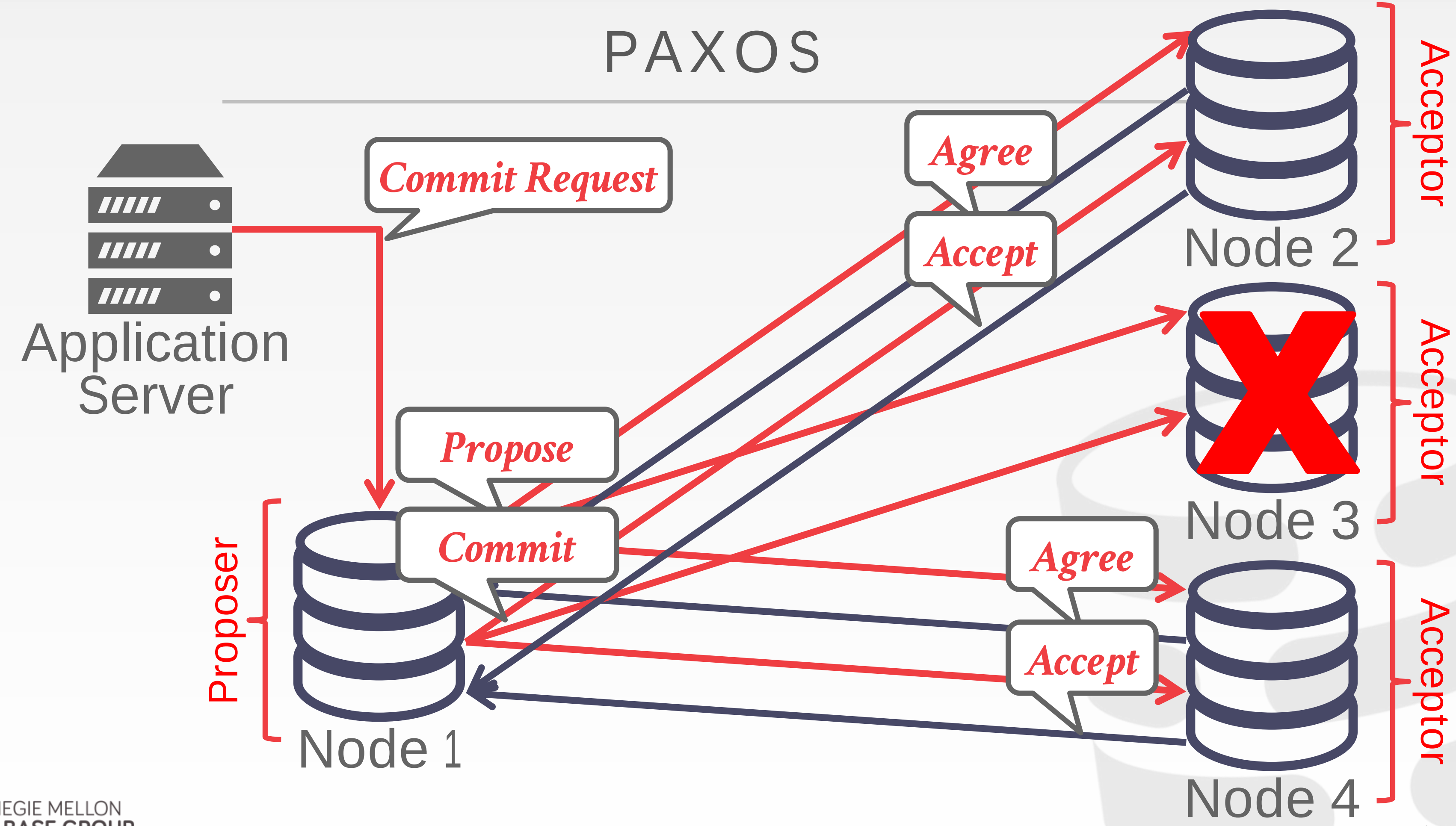
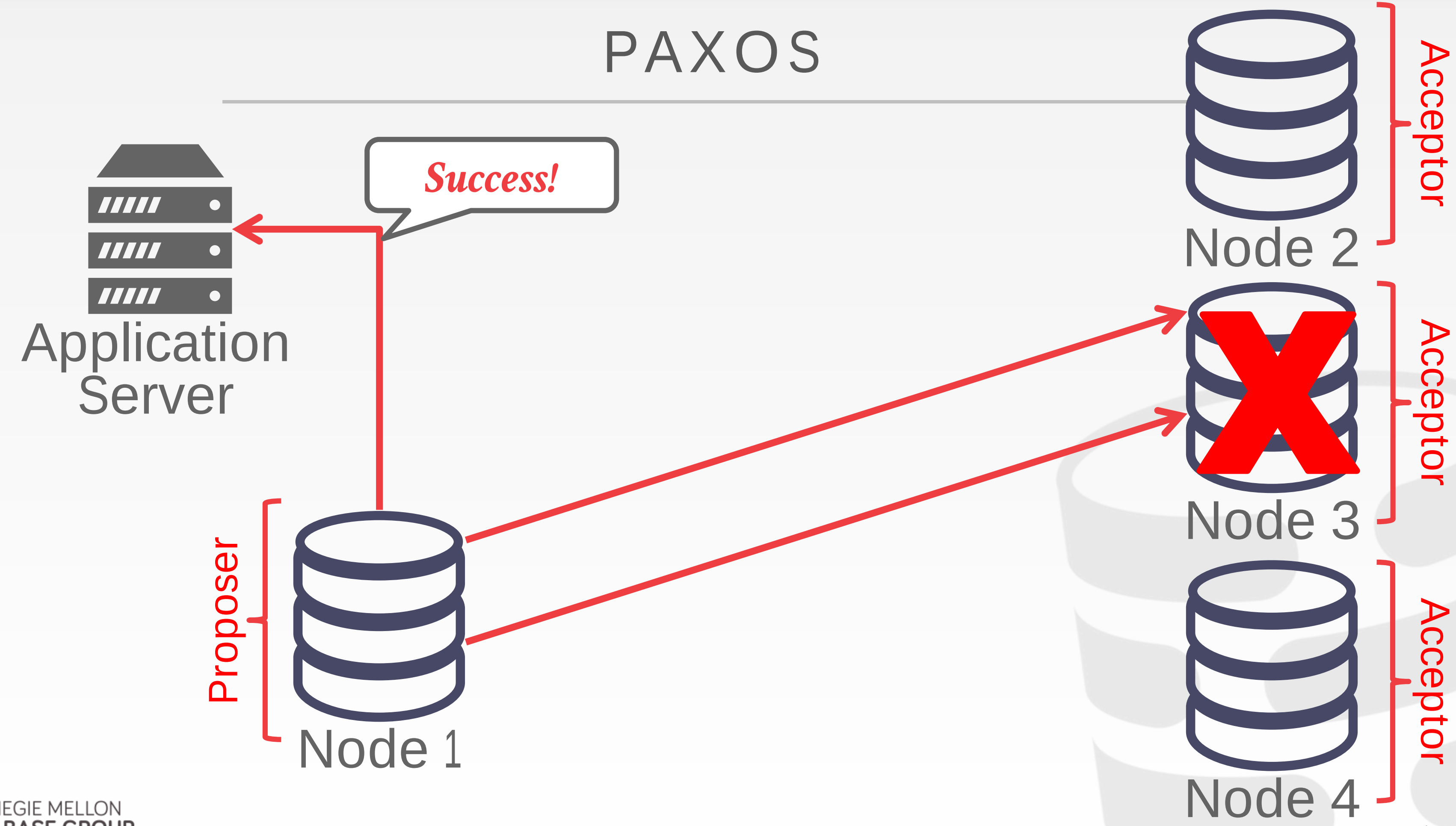
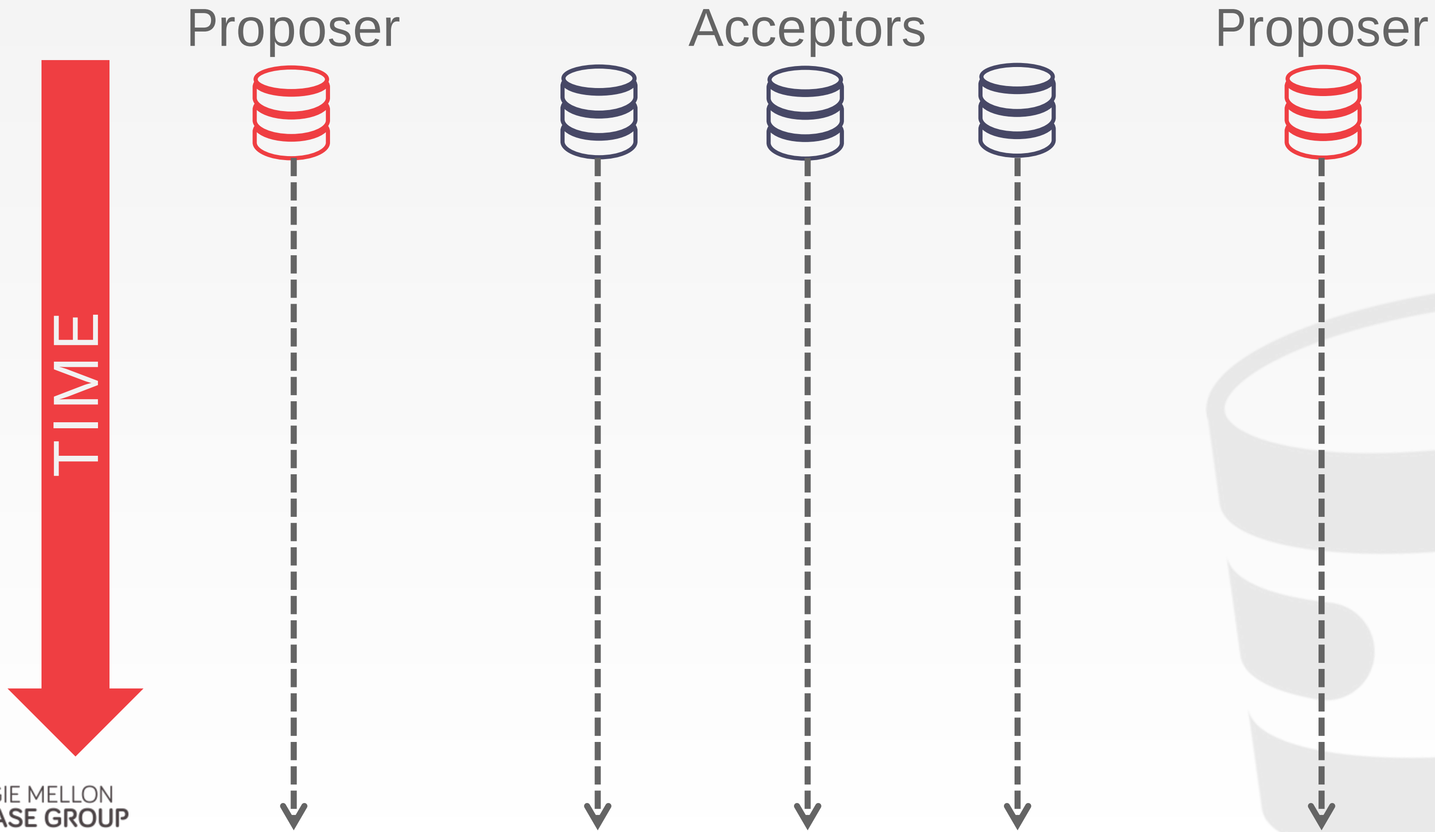
PAXOS


Application Server
Commit Request
Propose
Commit
Agree
Accept
Agree
Accept
Proposer
Node 1
Node 2
Node 3
Node 4
Acceptor
Acceptor
Acceptor

PAXOS


Success!
Application Server
Proposer
Node 1
Acceptor
Node 2
Acceptor
Node 3
Acceptor
Node 4

PAXOS



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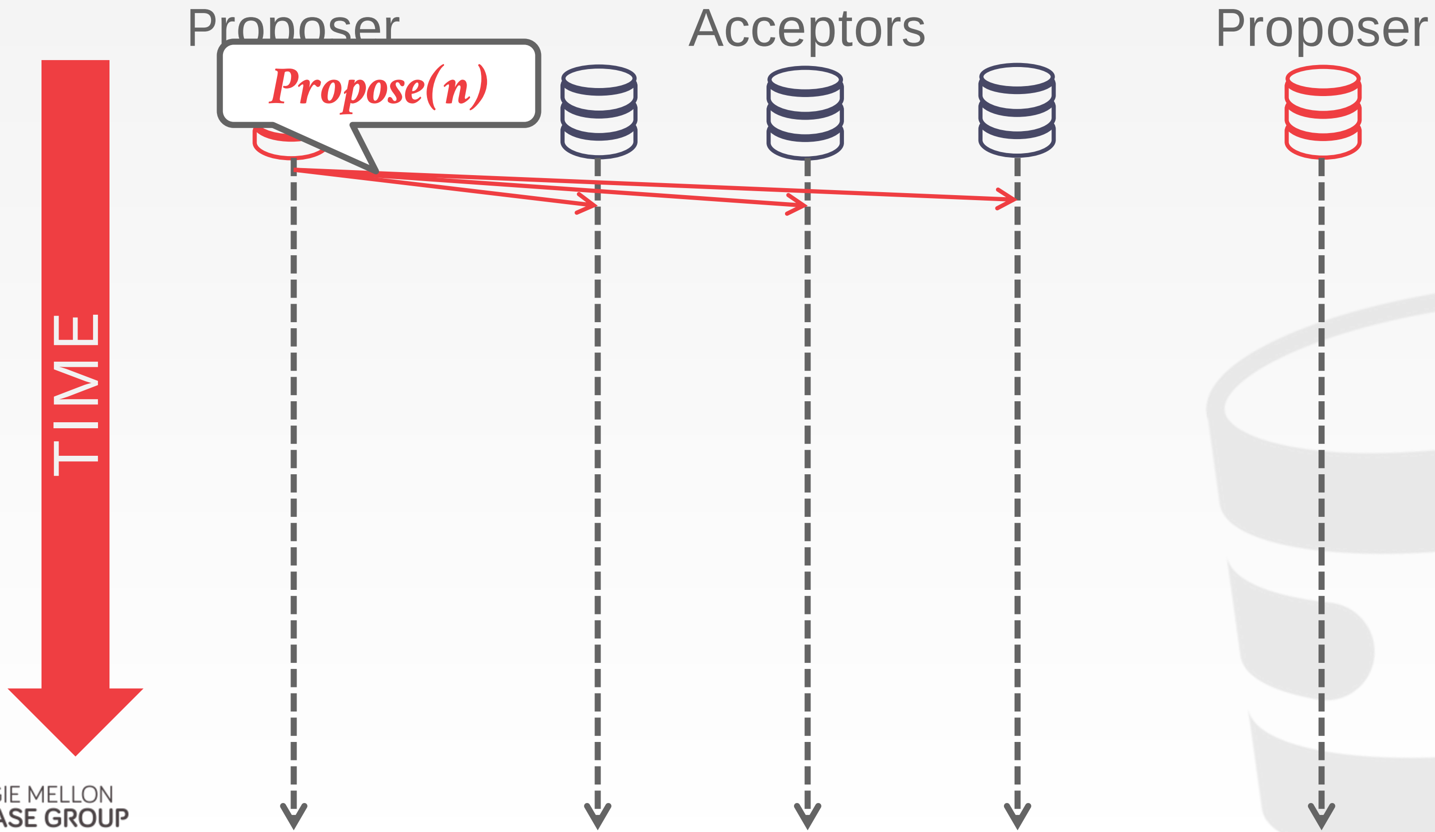

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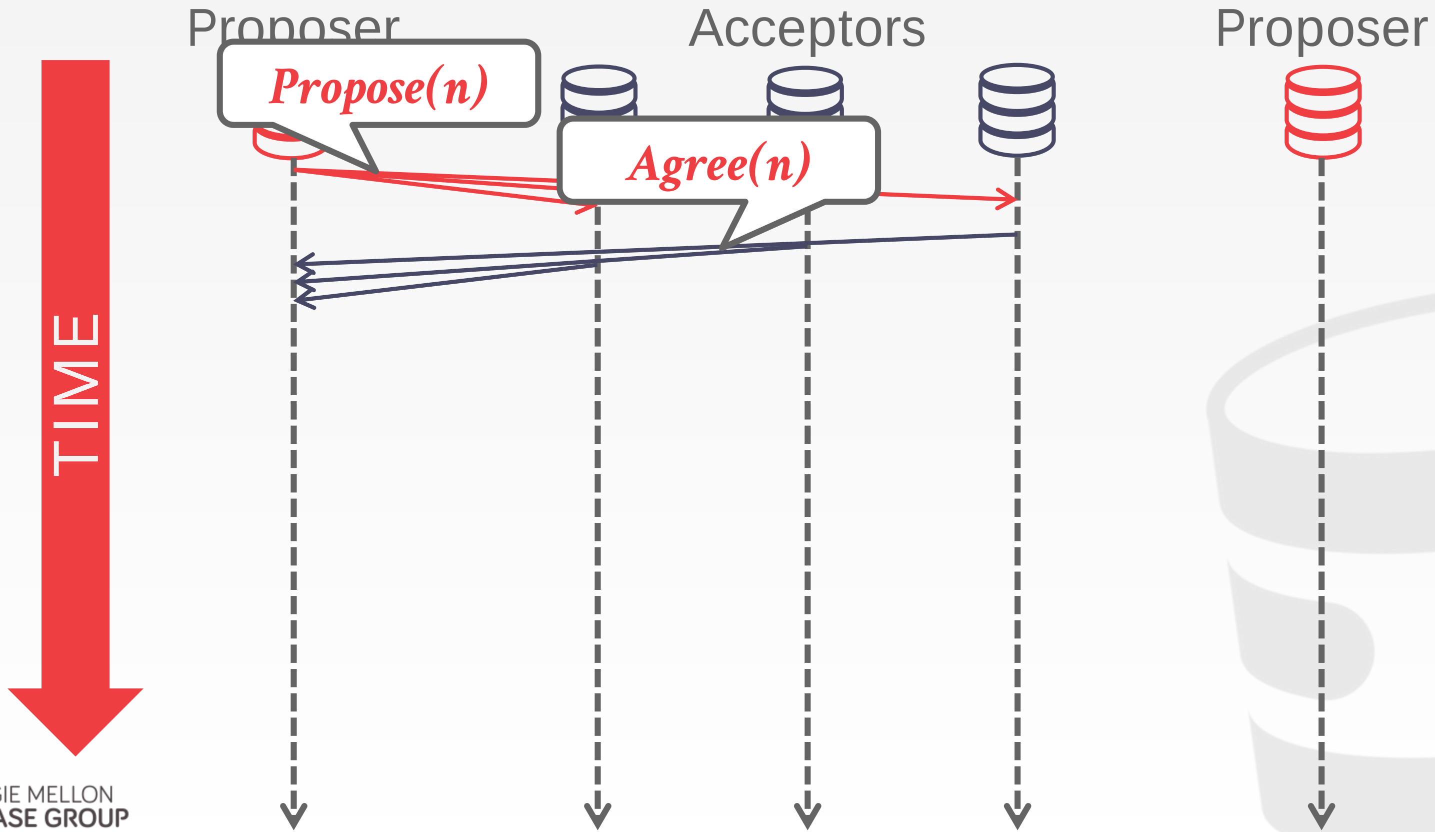

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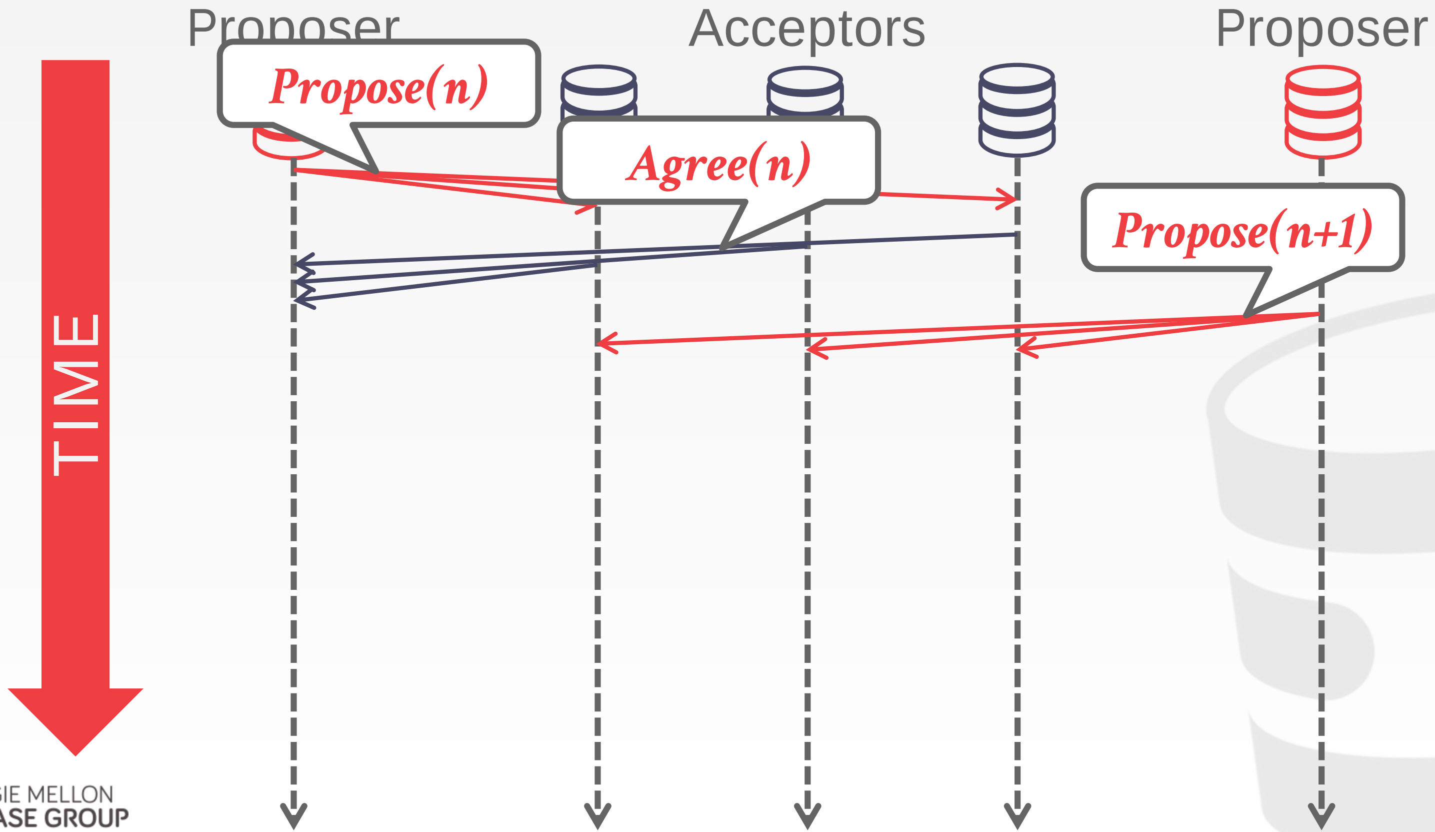
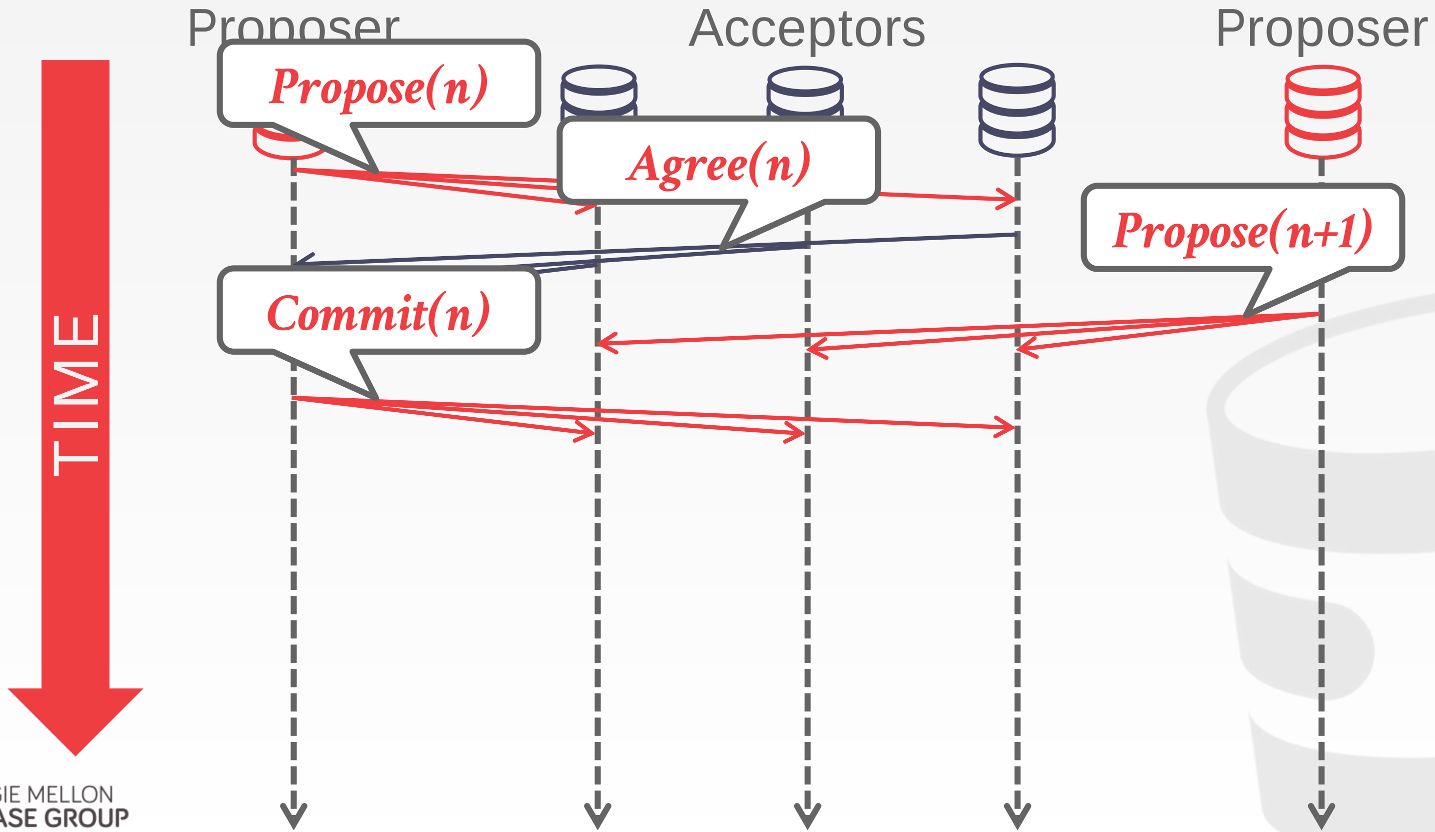
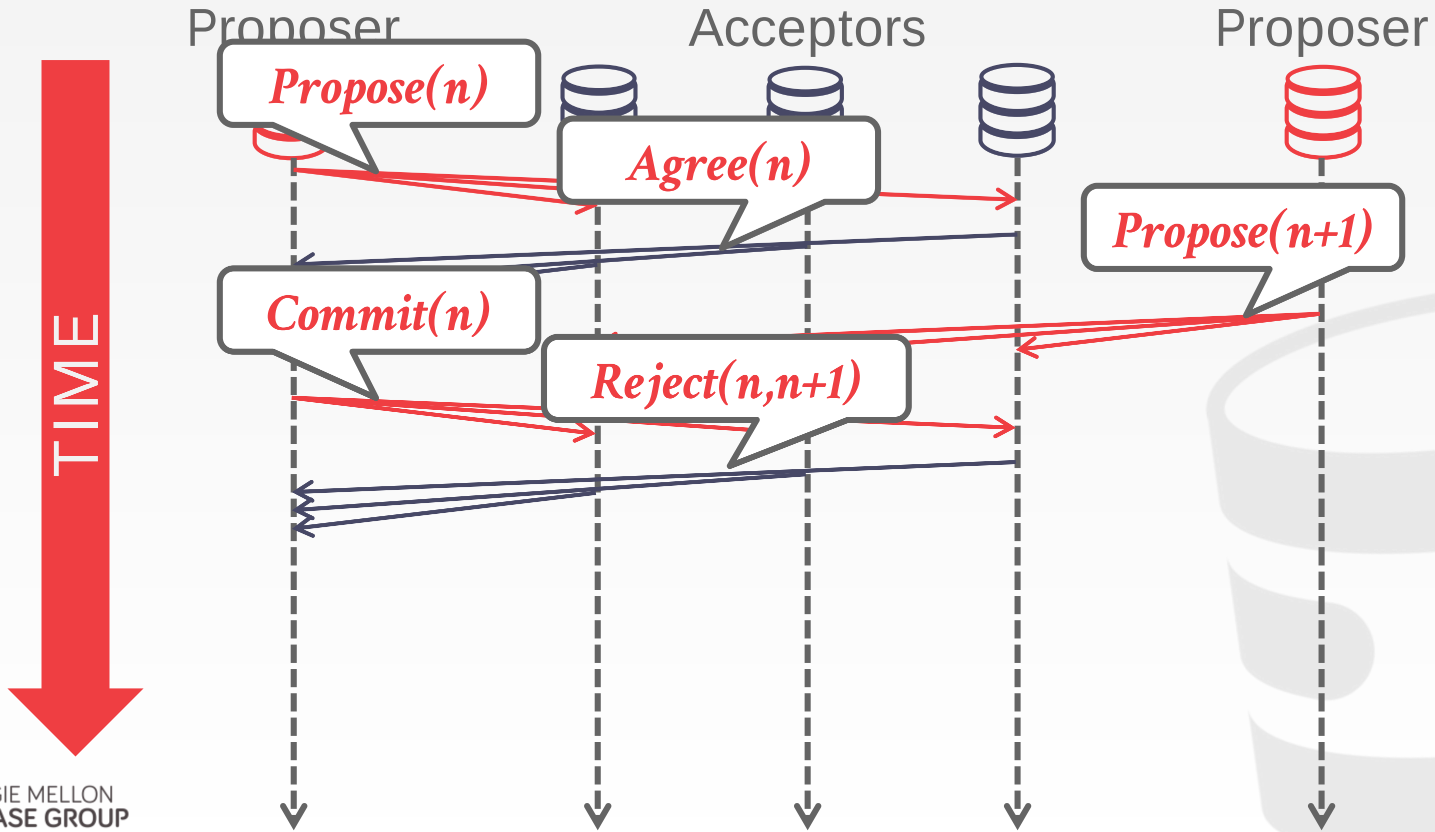

Proposer
Acceptors
Proposer
TIME
Propose(n)
Agree(n)
Propose(n+1)

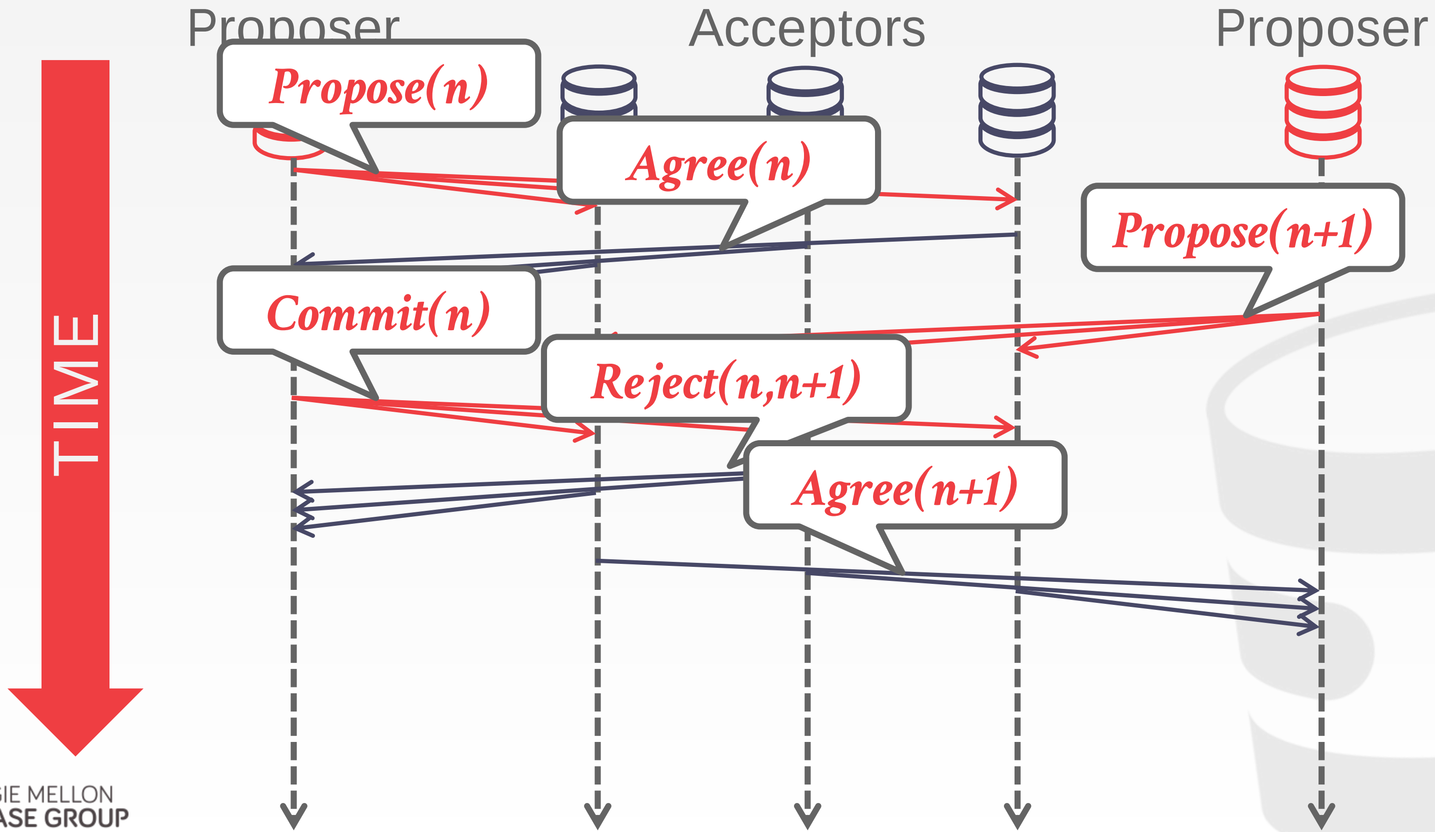
PAXOS


Proposer
Acceptors
Proposer
Propose(n)
Agree(n)
Propose(n+1)
Commit(n)
TIME

PAXOS


Proposer
Acceptors
Proposer
TIME
Propose(n)
Agree(n)
Propose(n+1)
Commit(n)
Reject(n,n+1)

PAXOS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Acceptors
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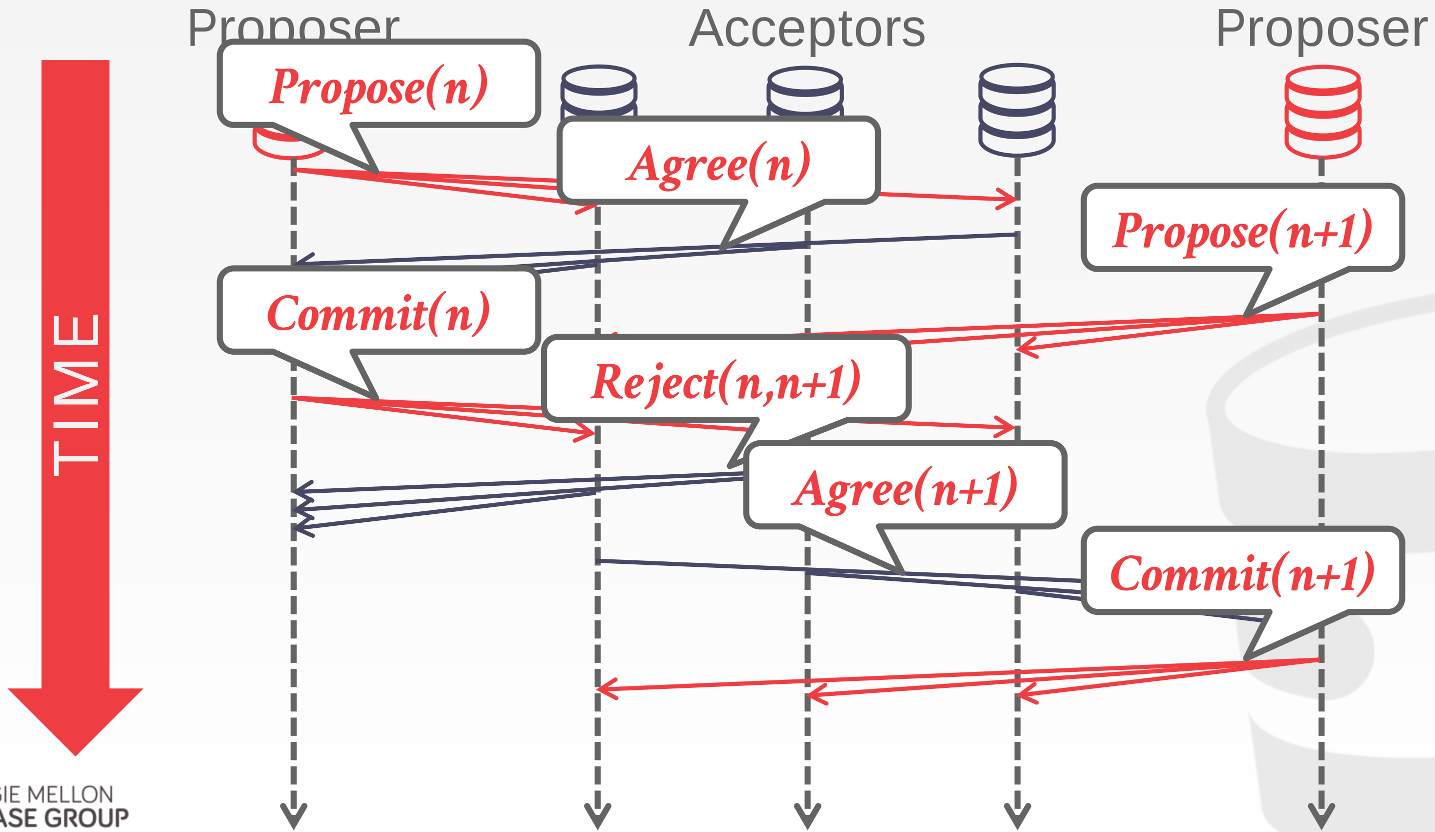
Propose(n)
Agree(n)
Propose(n+1)
Commit(n)
Reject(n,n+1)
Agree(n+1)

PAXOS



Proposer
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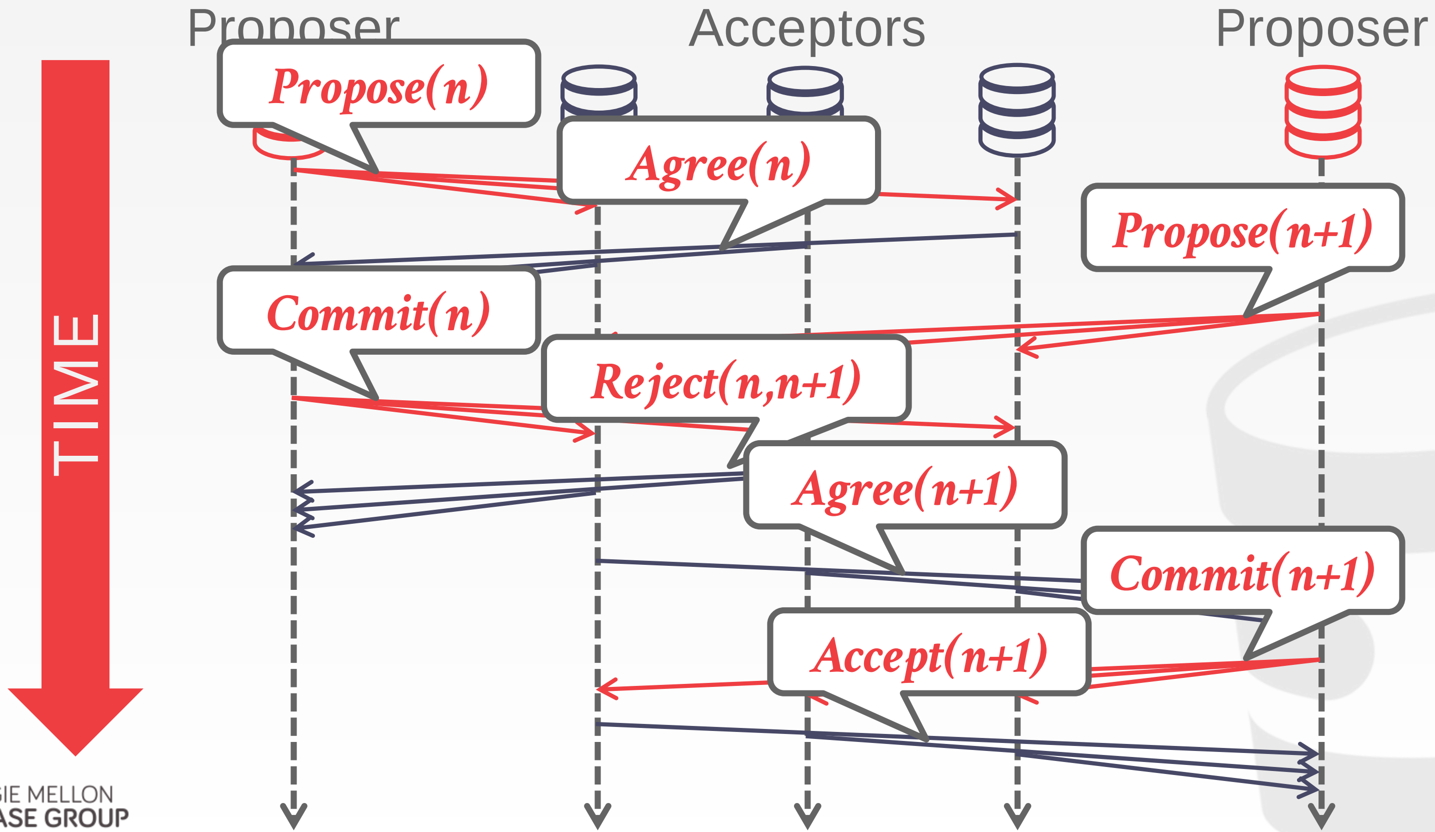
PAXOS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TIME
Propose(n)
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Propose(n+1)
Commit(n)
Reject(n,n+1)
Agree(n+1)
Commit(n+1)
Accept(n+1)

MULTI-PAXOS

If the system elects a single leader that is in charge of proposing changes for some period of time, then it can skip the **PREPARE** phase.

→ Fall back to full Paxos whenever there is a failure.

The system has to periodically renew who the leader is.

2PC VS. PAXOS

Two-Phase Commit

→ Blocks if coordinator fails after the prepare message is sent, until coordinator recovers.

Paxos

→ Non-blocking as long as a majority participants are alive, provided there is a sufficiently long period without further failures.

REPLICATION

The DBMS can replicate data across redundant nodes to increase availability.

Design Decisions:

→ Replica Configuration
→ Propagation Scheme
→ Propagation Timing

REPLICA CONFIGURATIONS

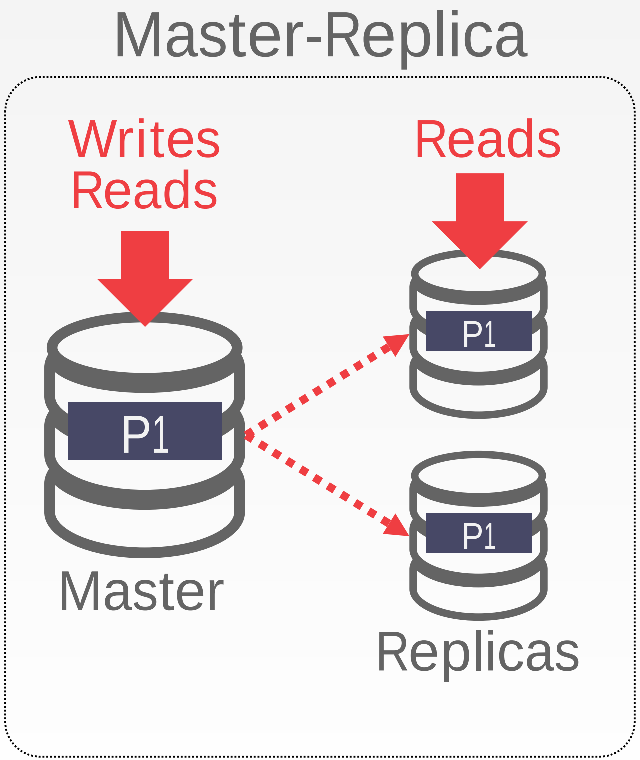
Approach #1: Master-Replica

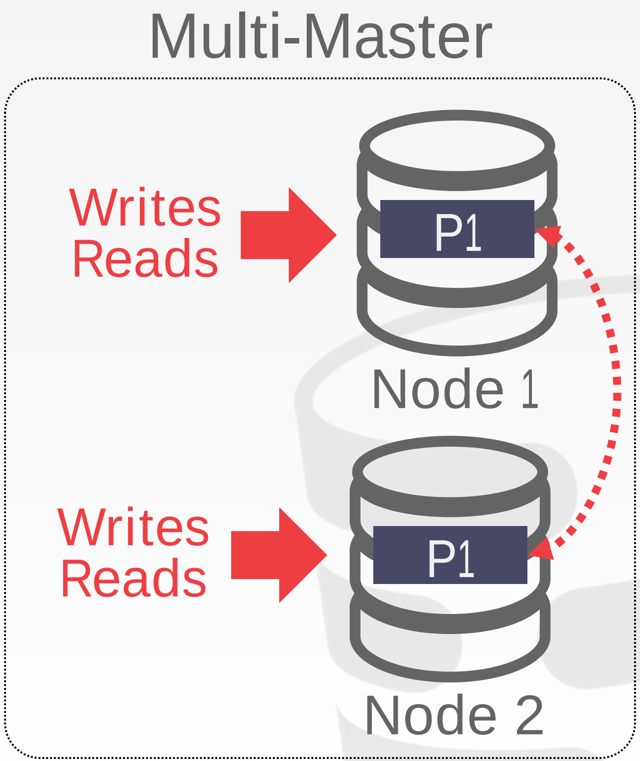
→ All updates go to a designated master for each object.
→ The master then propagates those updates to its replicas.
→ Read-only txns may be allowed to access replicas.
→ If the master goes down, then hold an election to select a new master.

Approach #2: Multi-Master

→ Txns can update data objects at any replica.
→ Replicas synchronize with each other.

REPLICA CONFIGURATIONS





K-SAFETY

***K*-safety is a threshold for determining the fault tolerance of the replicated database.**

The value ***K*** represents the number of replicas per data object that must exist at all times.

If the number of replicas goes below this threshold, then the DBMS halts execution and takes itself offline.

PROPAGATION SCHEME

When a txn commits on a replicated database, the DBMS has to decide whether it has to wait for that txn's changes to propagate to other nodes before it can send the acknowledgement to application.

Propagation levels:
→ Synchronous
→ Asynchronous
→ Semi-Synchronous

PROPAGATION SCHEME

Approach #1: Synchronous

→ The master sends updates to replicas and then waits for them to acknowledge that they fully applied (i.e., logged) the changes.

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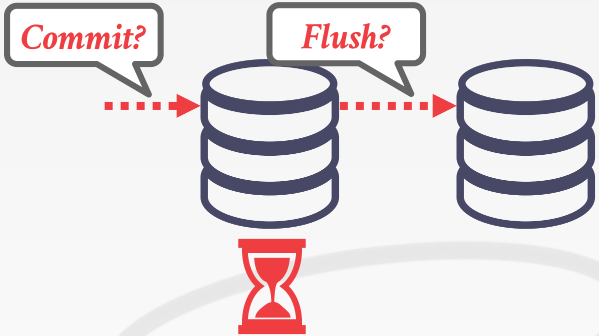

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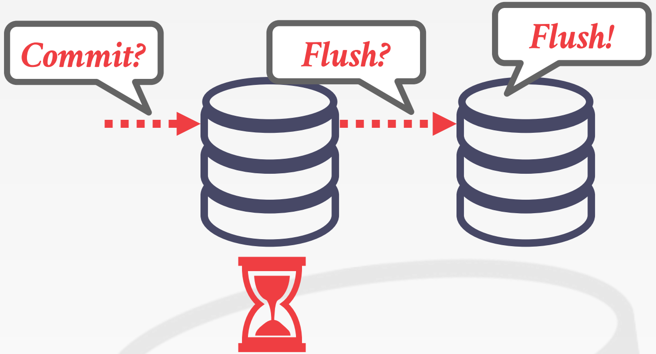

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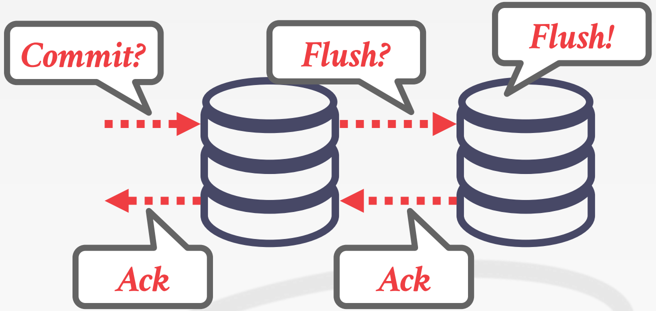

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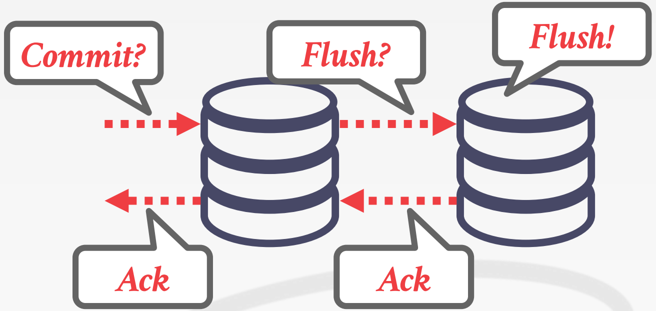

Approach #2: Asynchronous

→ The master immediately returns the acknowledgement to the client without waiting for replicas to apply the changes.

PROPAGATION SCHEME

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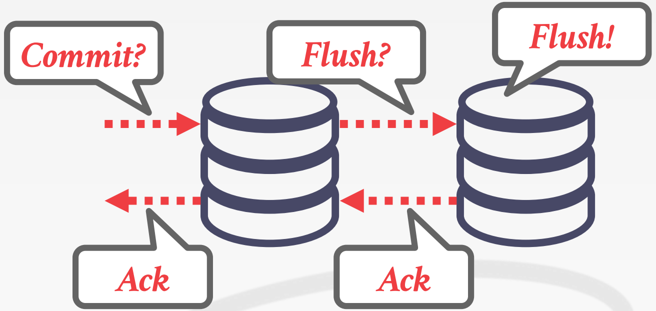

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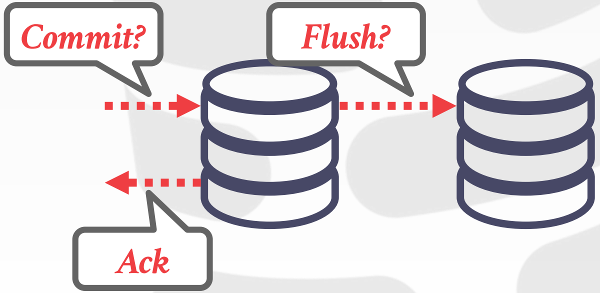

PROPAGATION SCHEME

Approach #3: Semi-Synchronous

→ Replicas immediately send acknowledgements without logging them.

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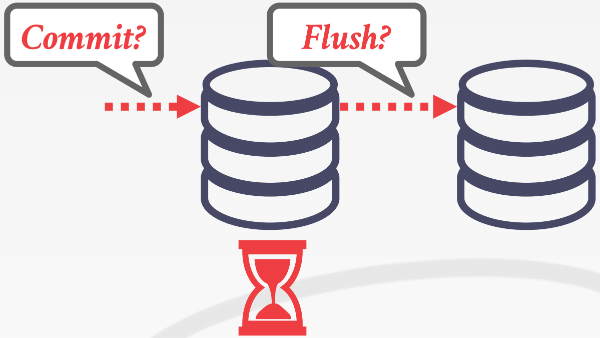

PROPAGATION SCHEME

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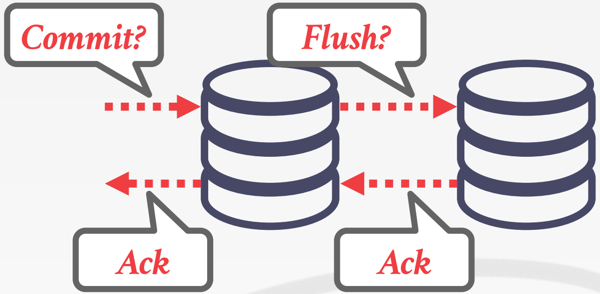

PROPAGATION SCHEME

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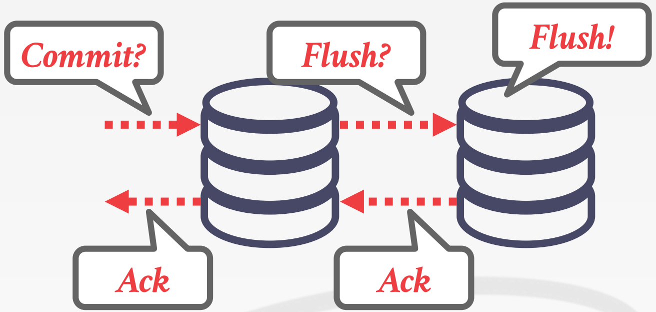


Applications can make trade-offs on protecting the integrity of the database versus performance.

PROPAGATION TIMING

Approach #1: Continuous

→ The DBMS sends log messages immediately as it generates them.
→ Also need to send a commit/abort message.

Approach #2: On Commit

→ The DBMS only sends the log messages for a txn to the replicas once the txn is commits.
→ Do not waste time sending log records for aborted txns.
→ Assumes that a txn's log fits entirely in memory.

ACTIVE VS. PASSIVE

Approach #1: Active-Active

→ A txn executes at each replica independently.
→ Need to check at the end whether the txn ends up with the same result at each replica.

Approach #2: Active-Passive

→ Each txn executes at a single location and propagates the changes to the replica.
→ Not the same as master-replica vs. multi-master

CAP THEOREM

Proposed by Eric Brewer that it is impossible for a distributed system to always be:

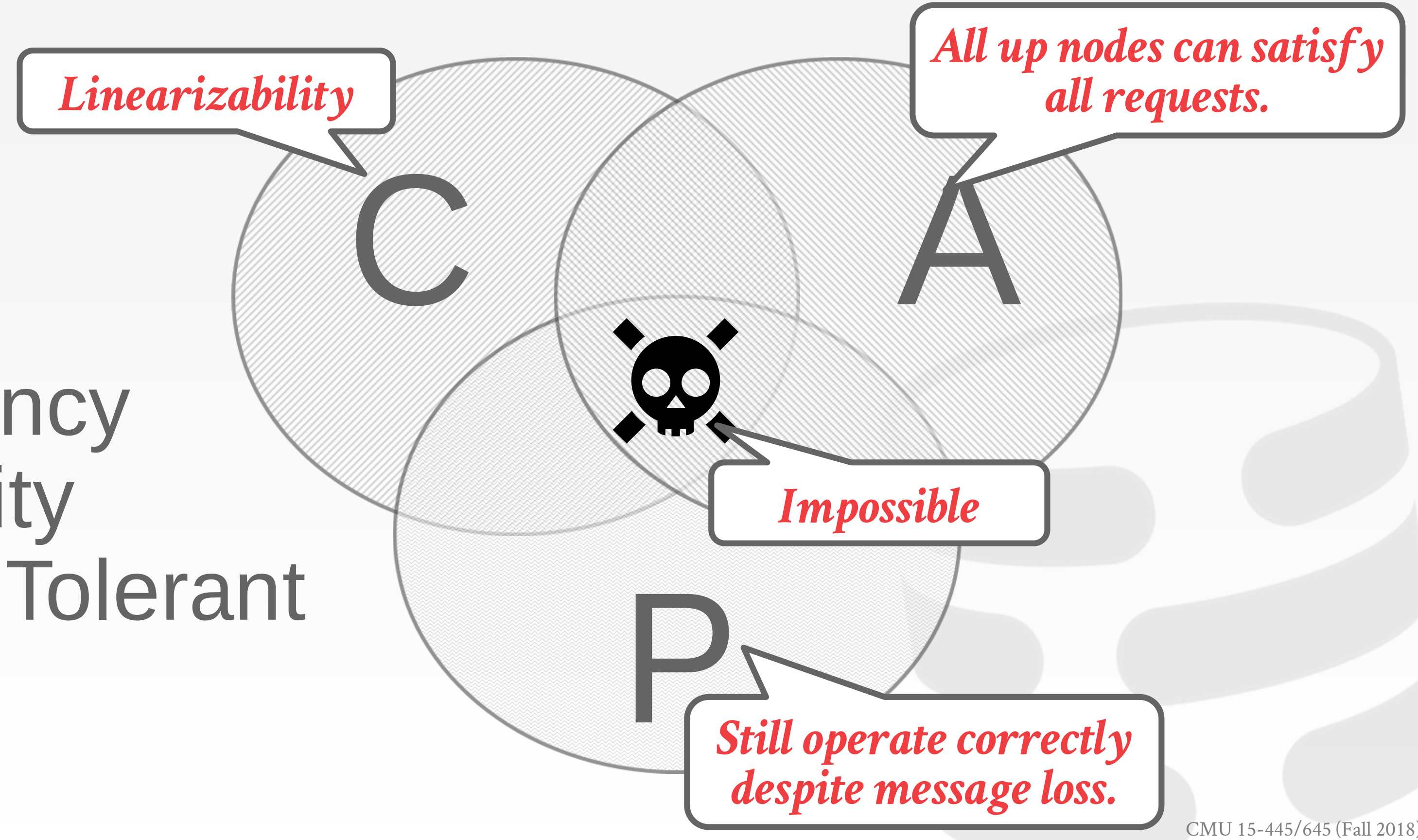
→ Consistent
→ Always Available
→ Network Partition Tolerant

Proved in 2002.

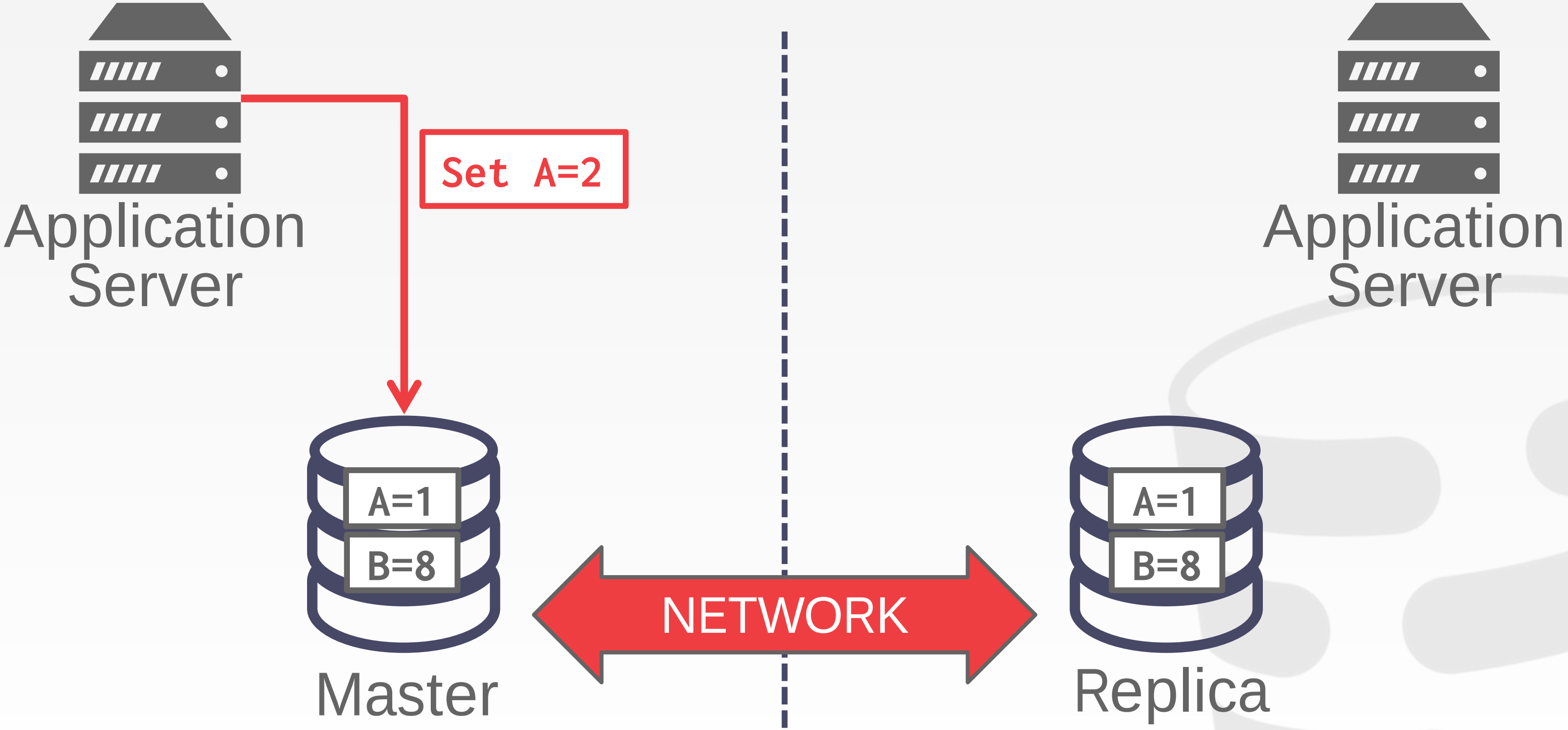
Brewer

CAP THEOREM

Consistency
Availability
Partition Tolerant



CAP – CONSISTENCY



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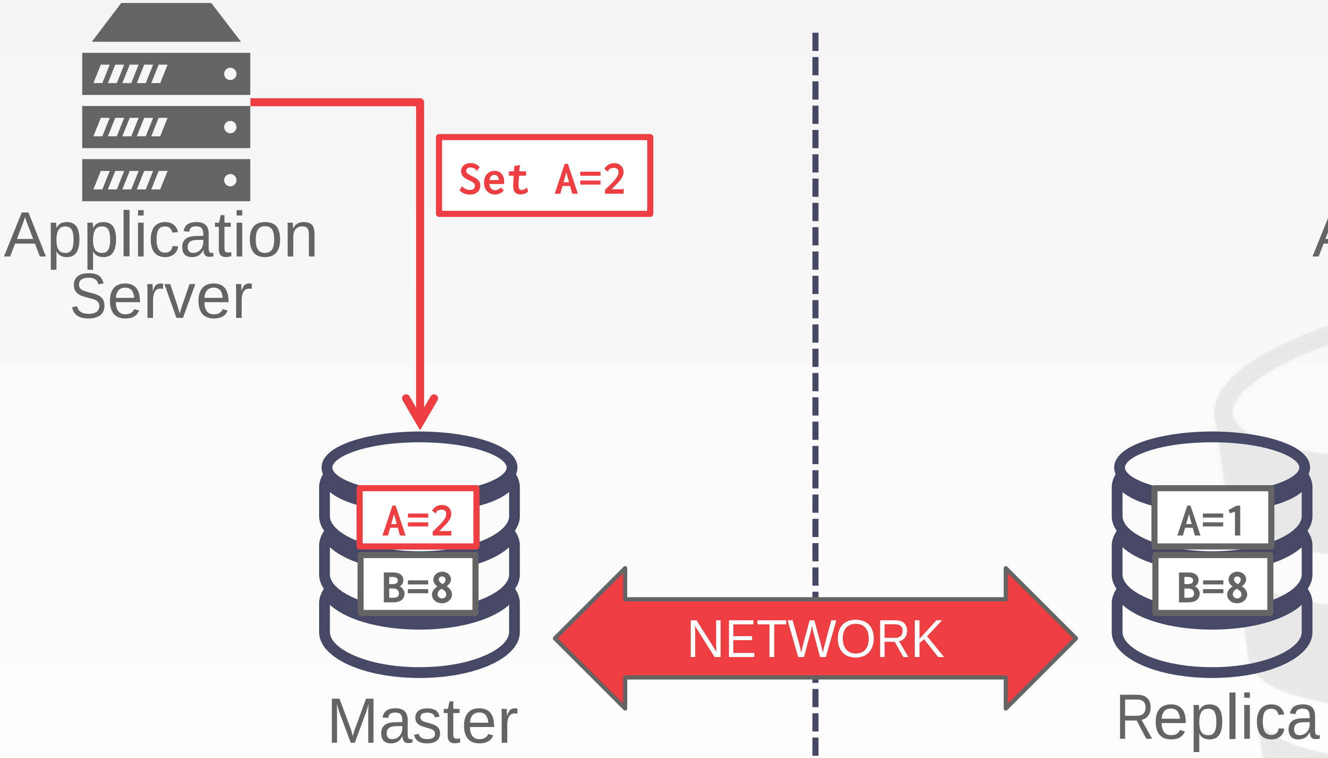

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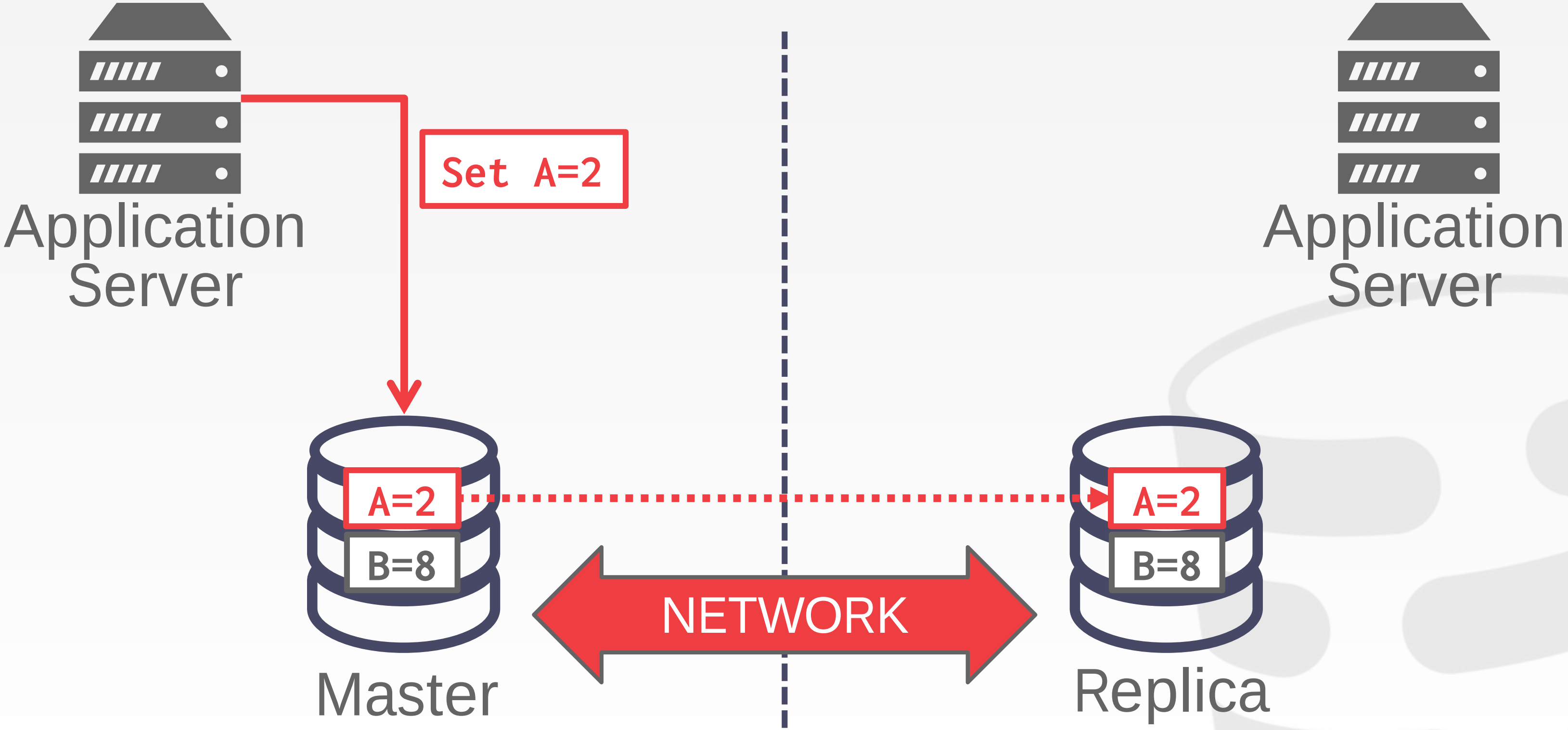

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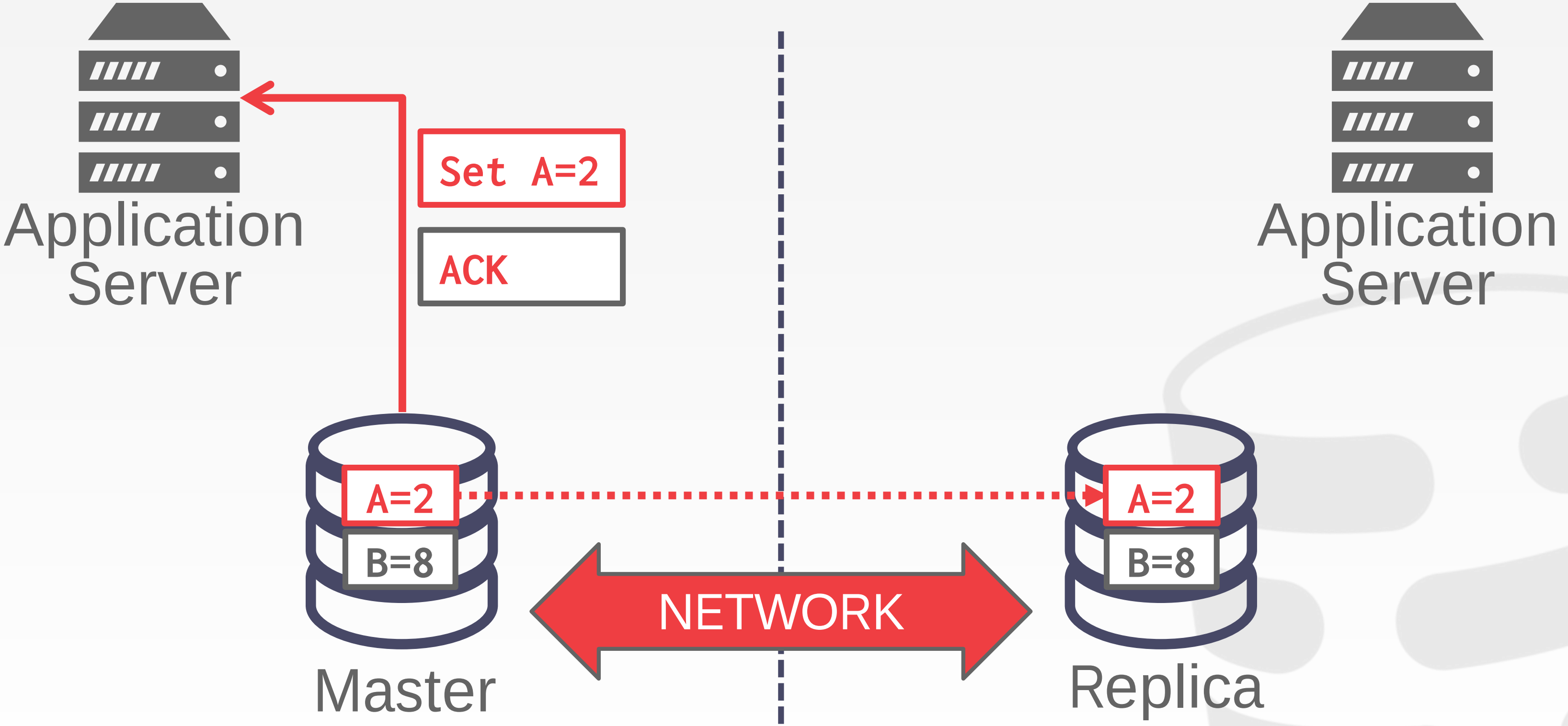

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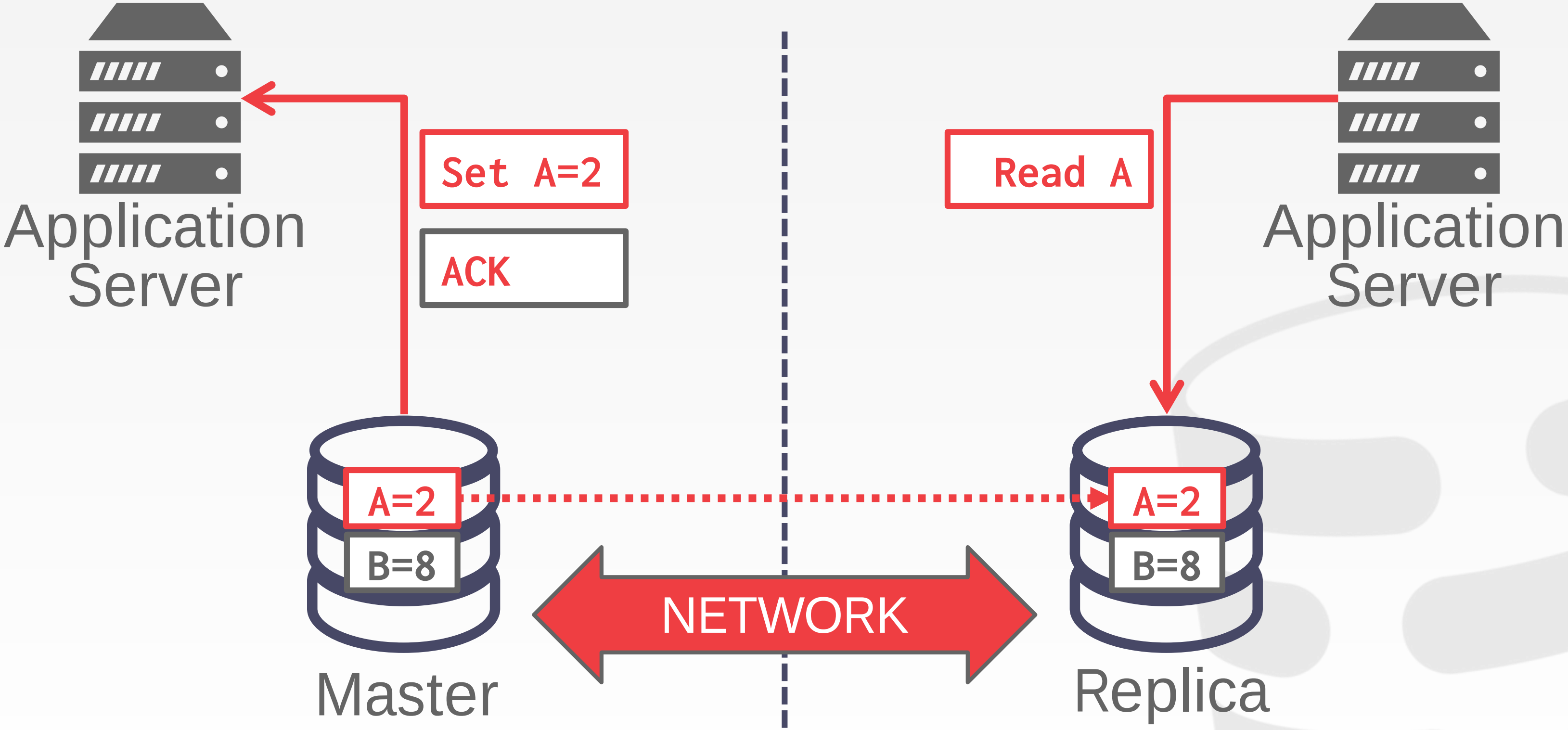

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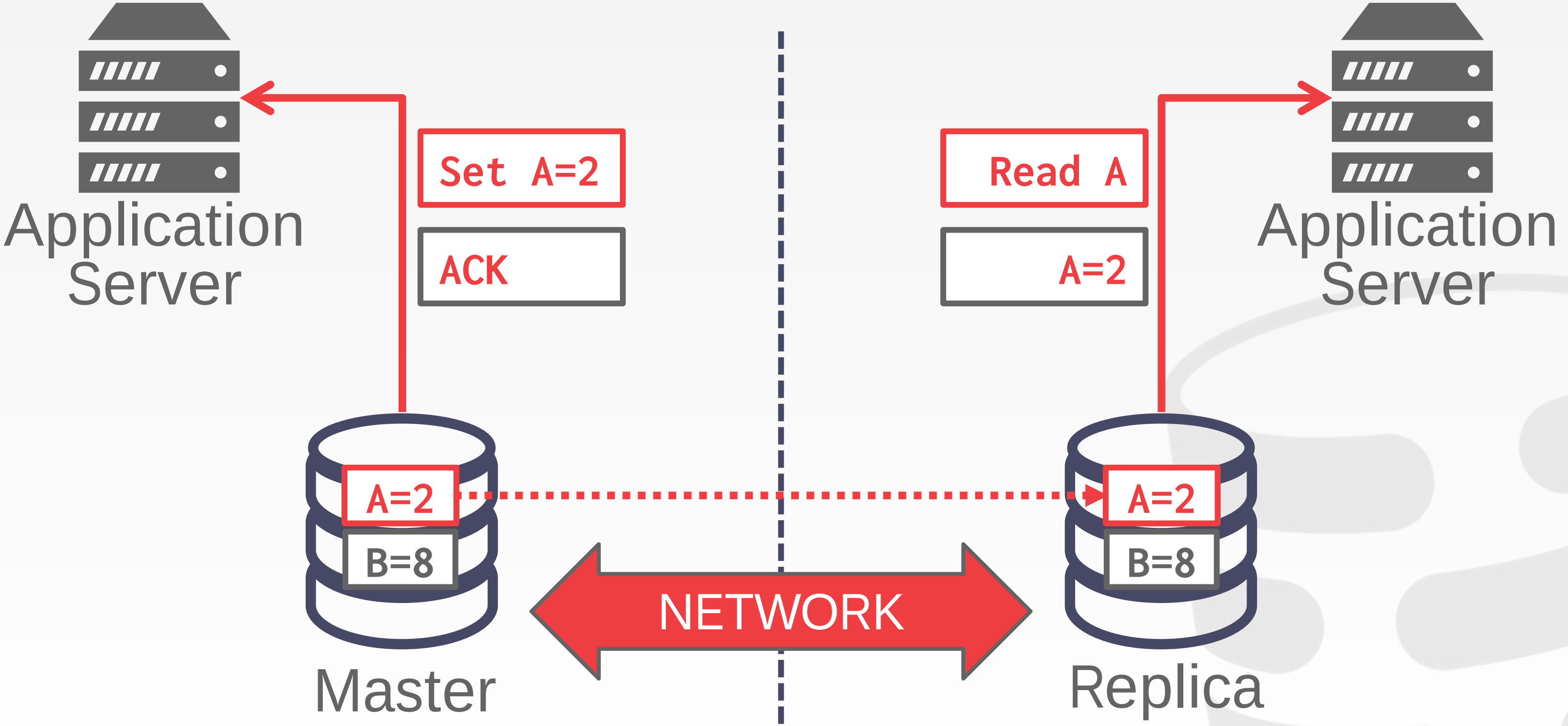
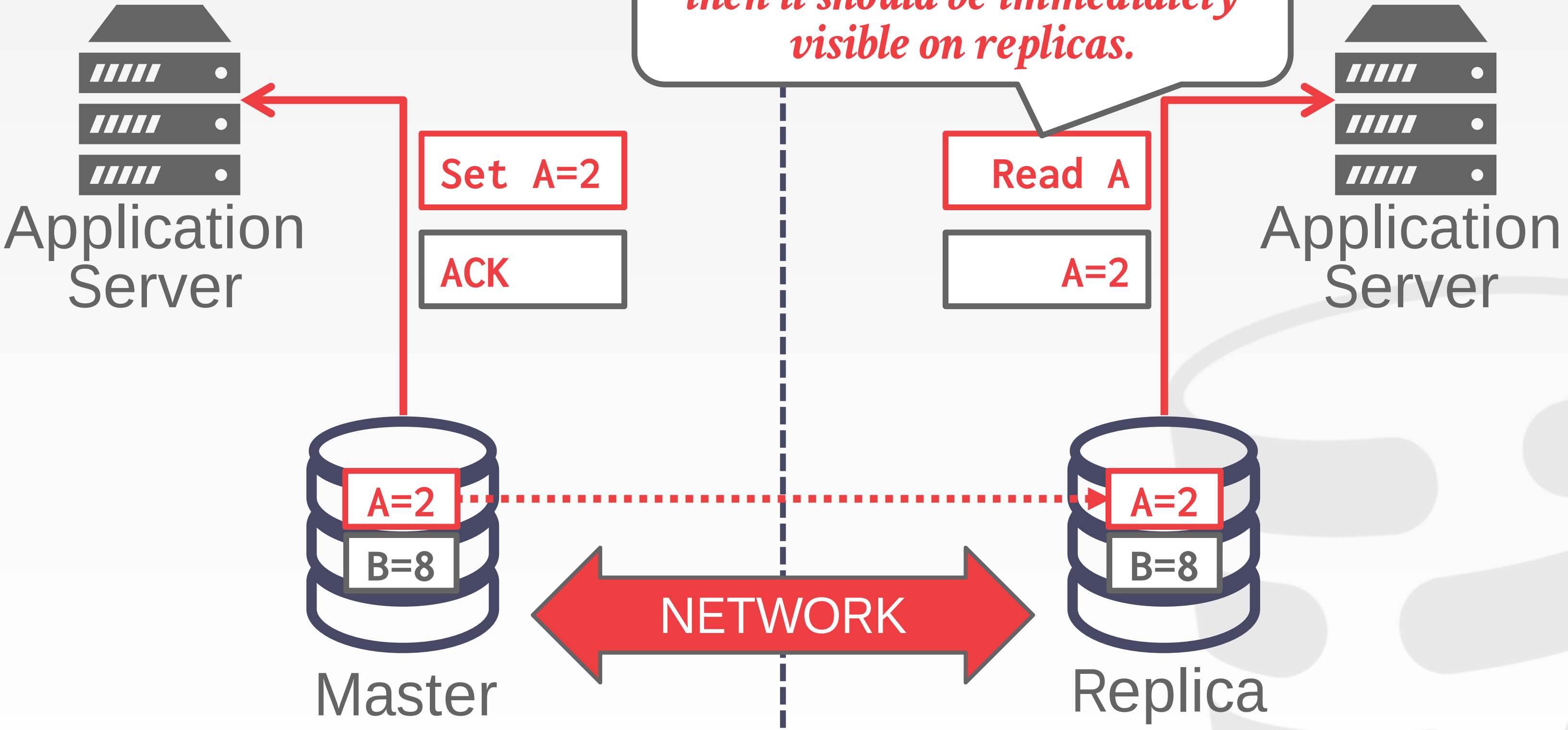
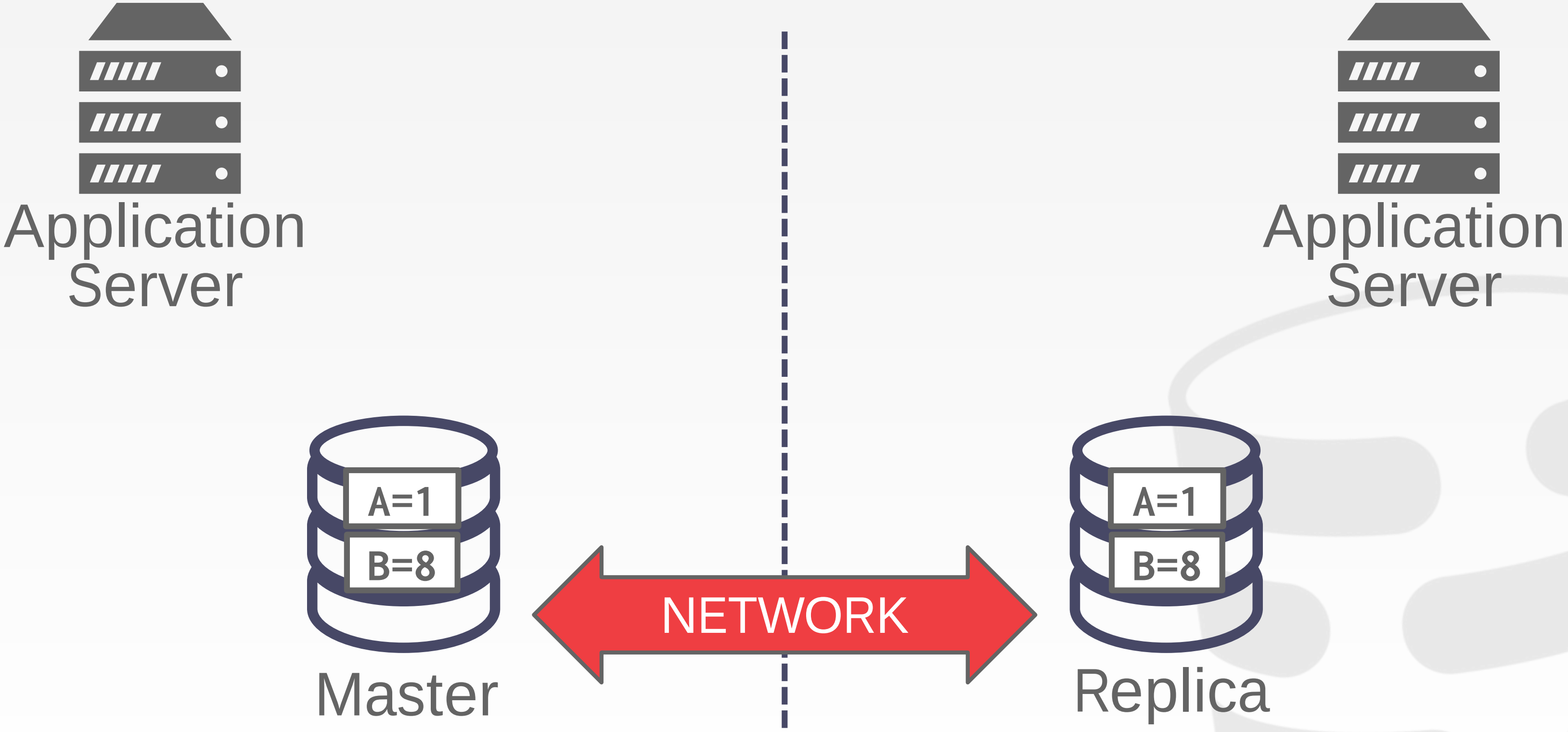


CAP –

If master says the txn committed, then it should be immediately visible on replicas.



CAP – AVAILABILITY



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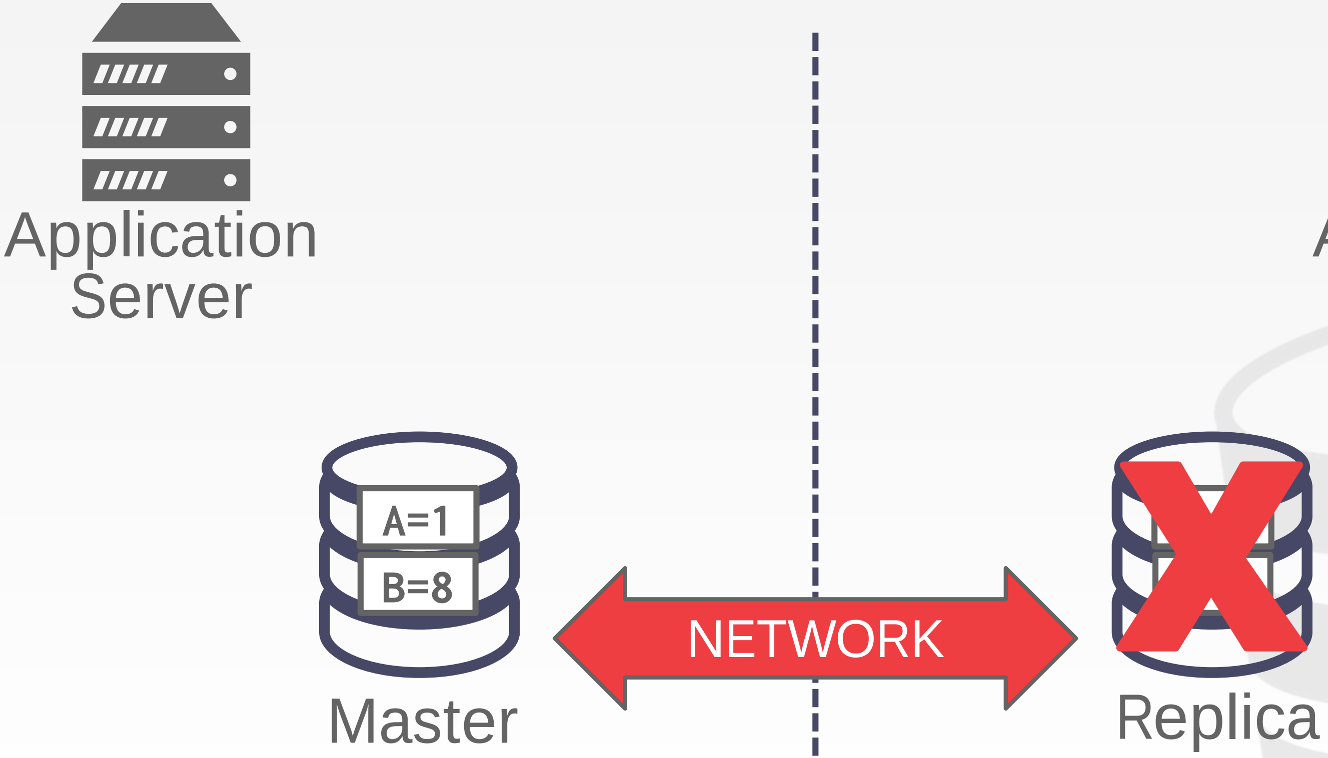

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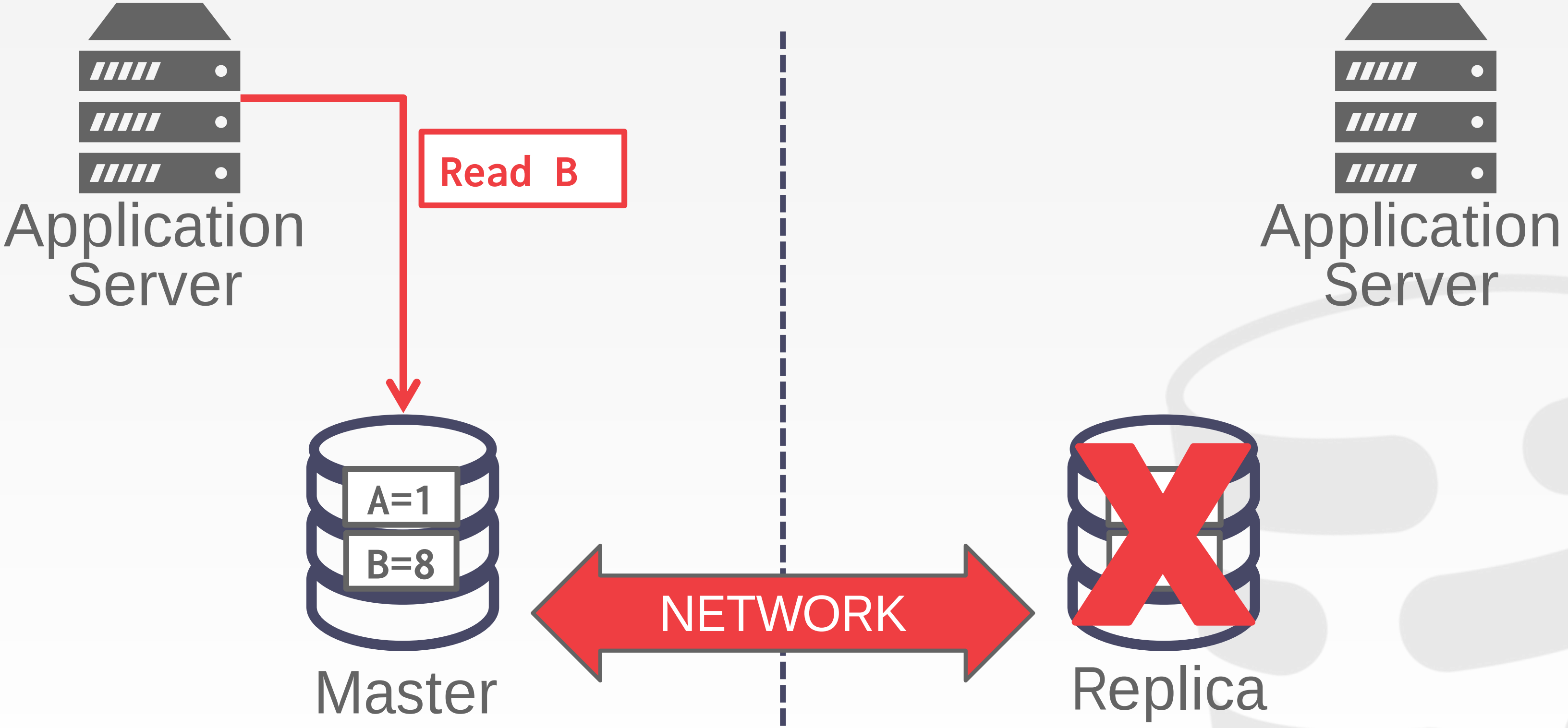

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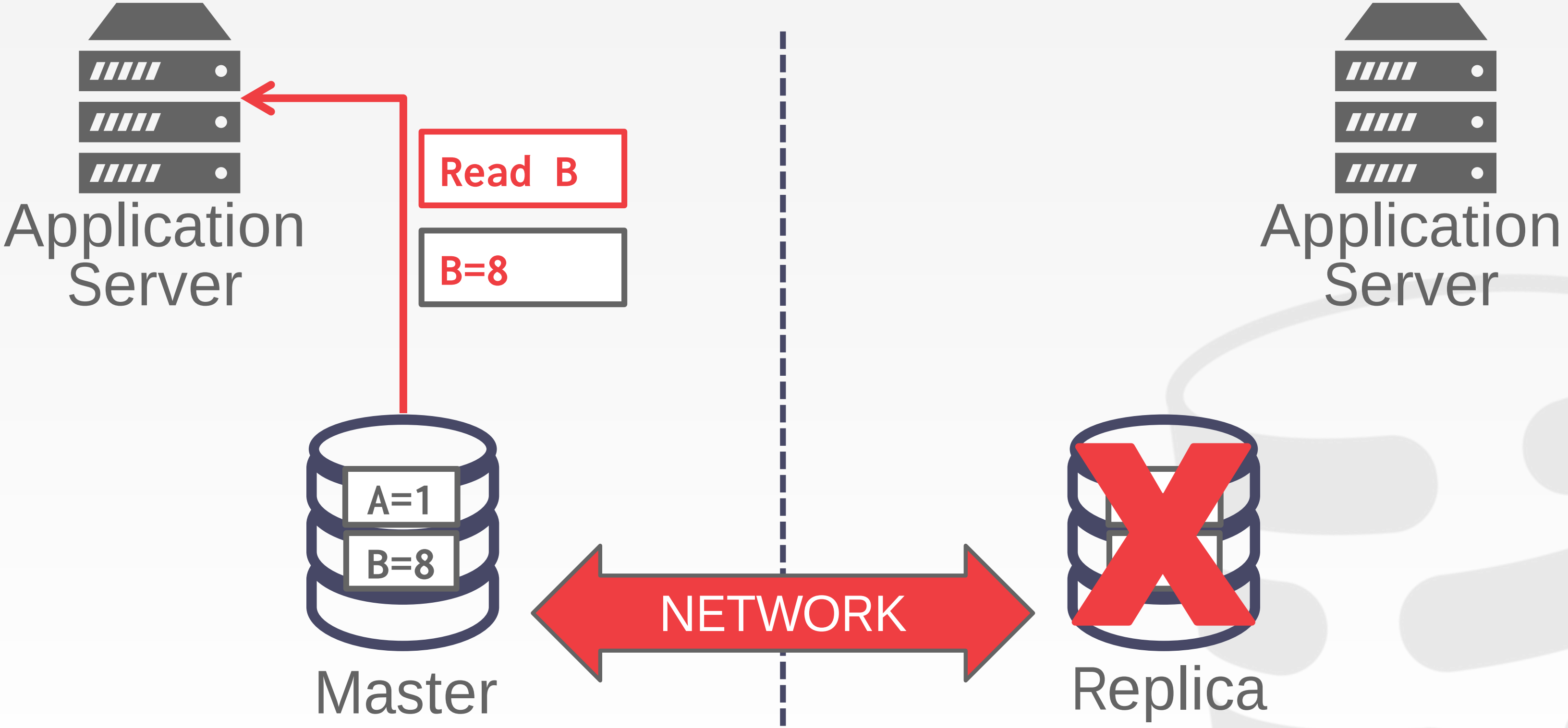

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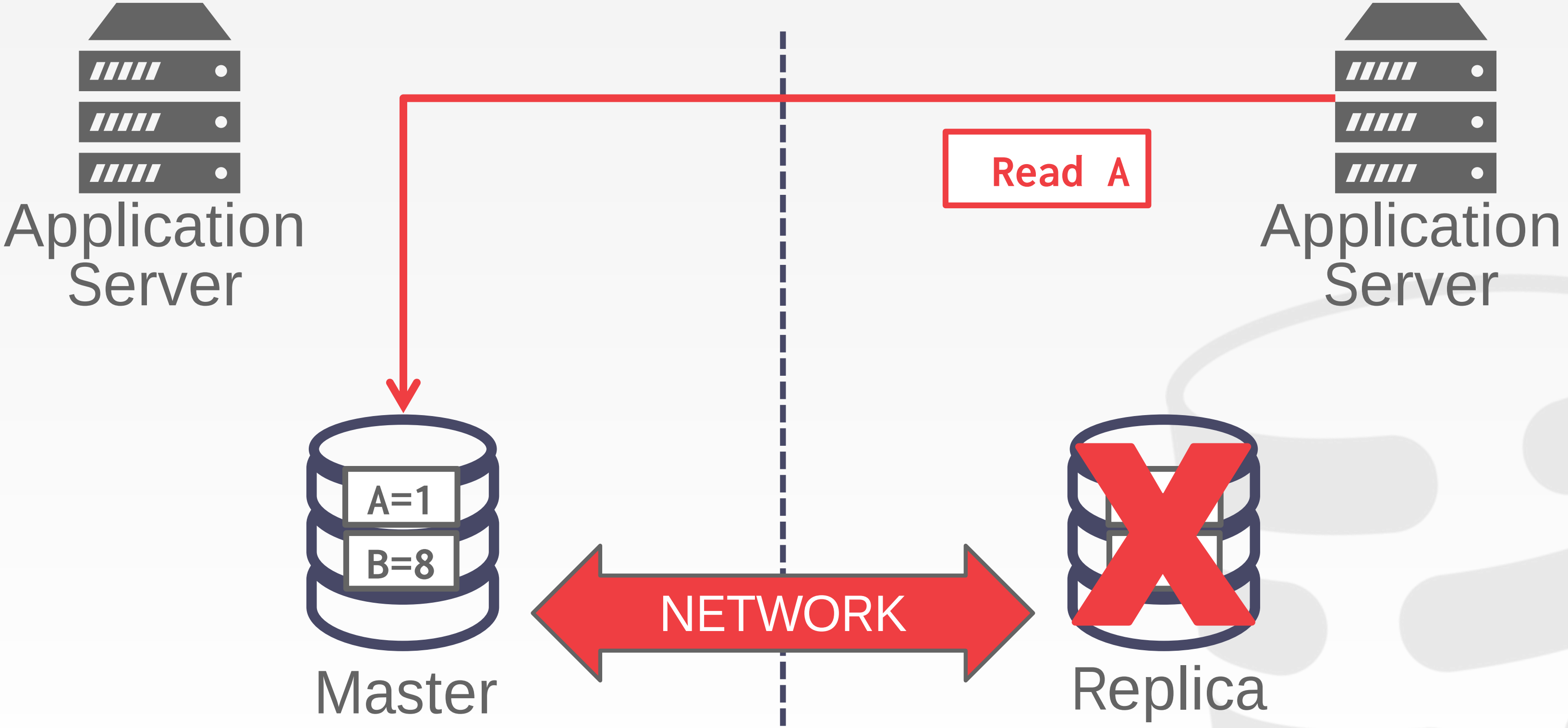

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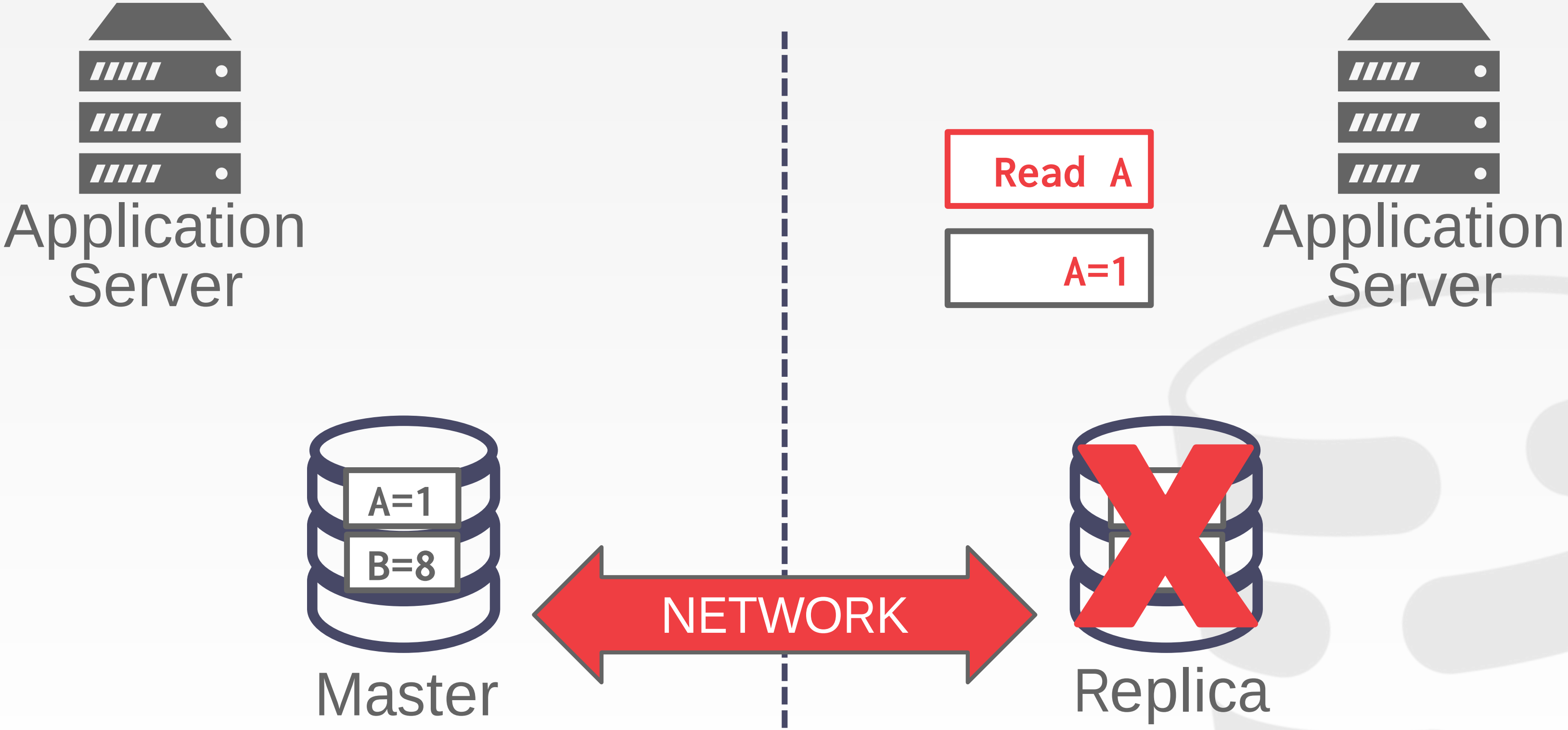
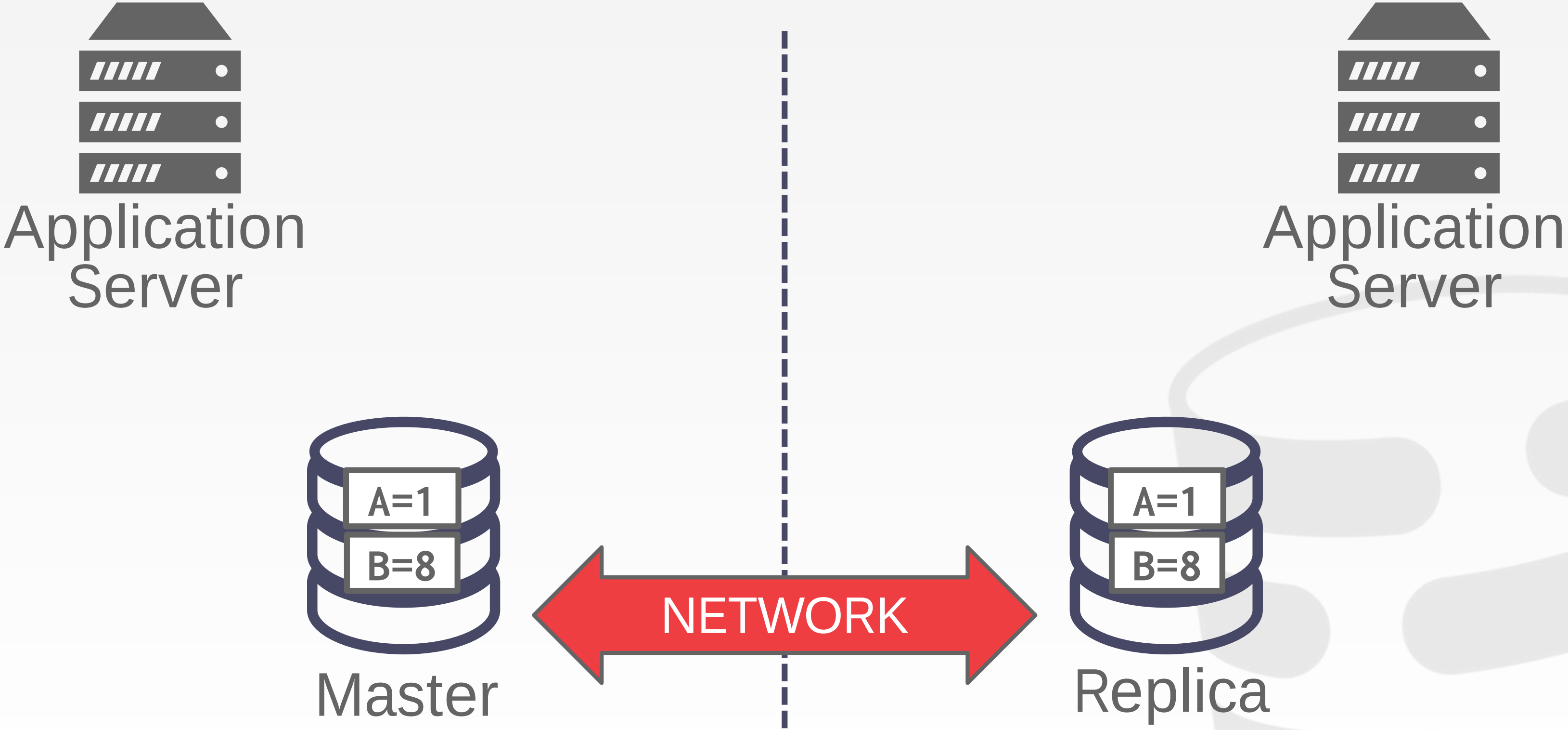


CAP – PARTITION TOLERANCE



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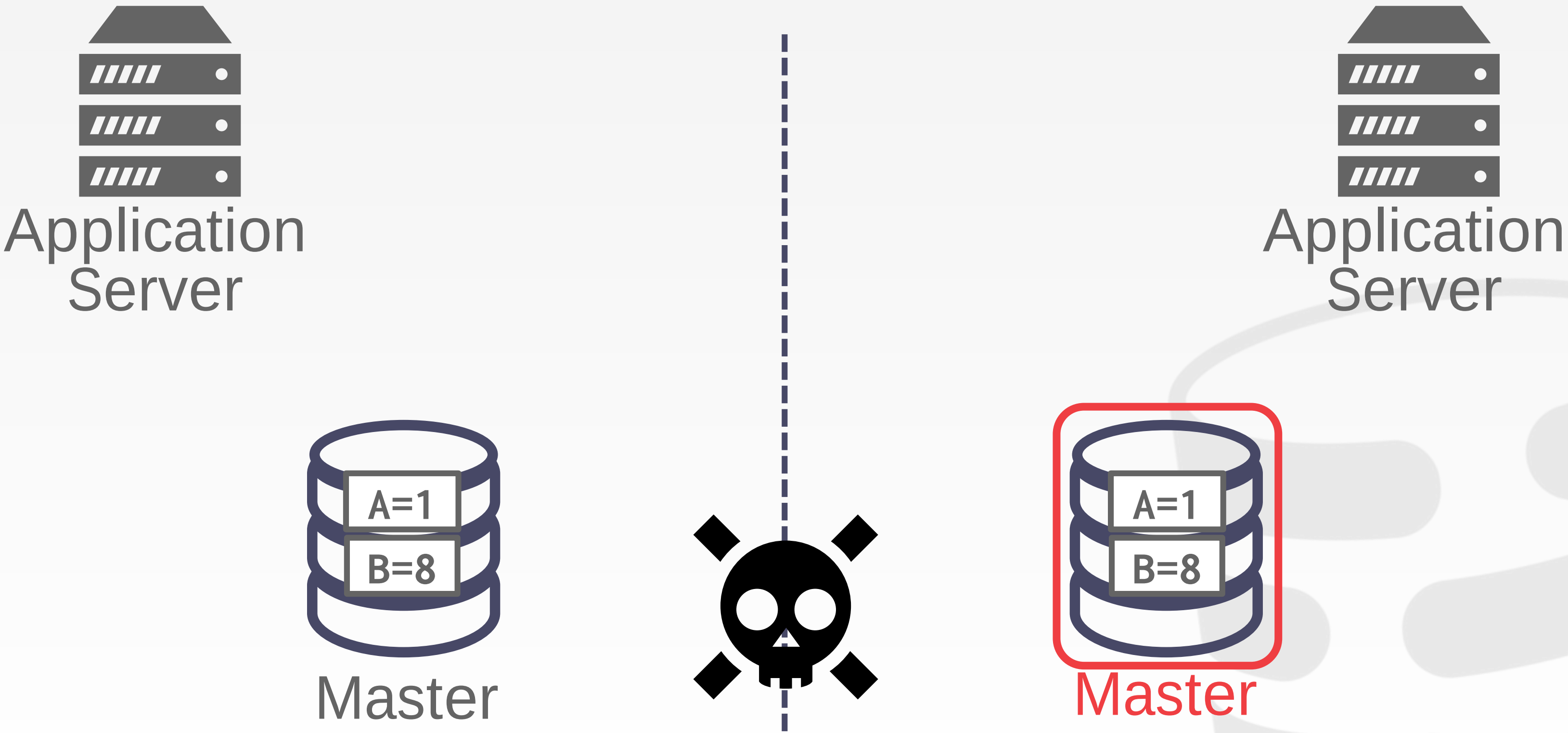

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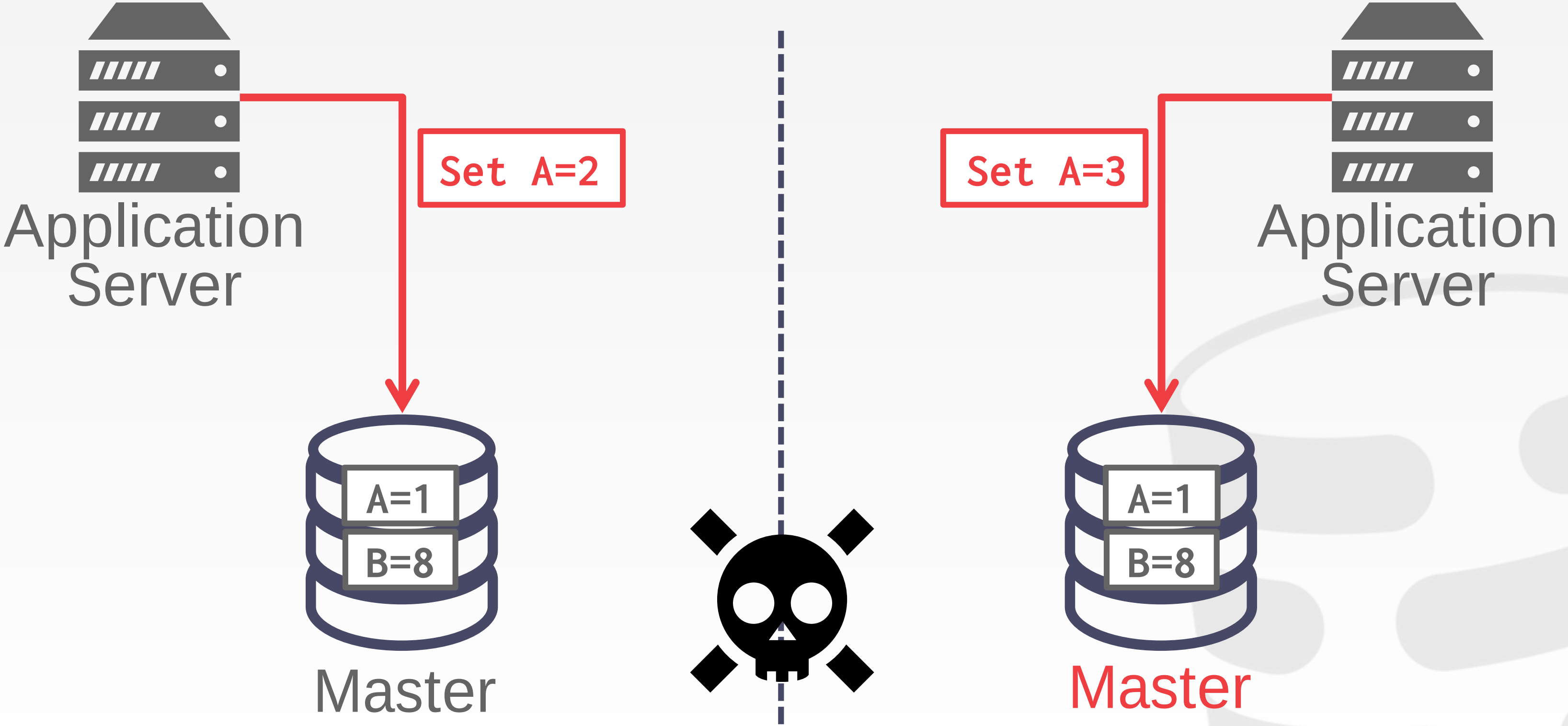

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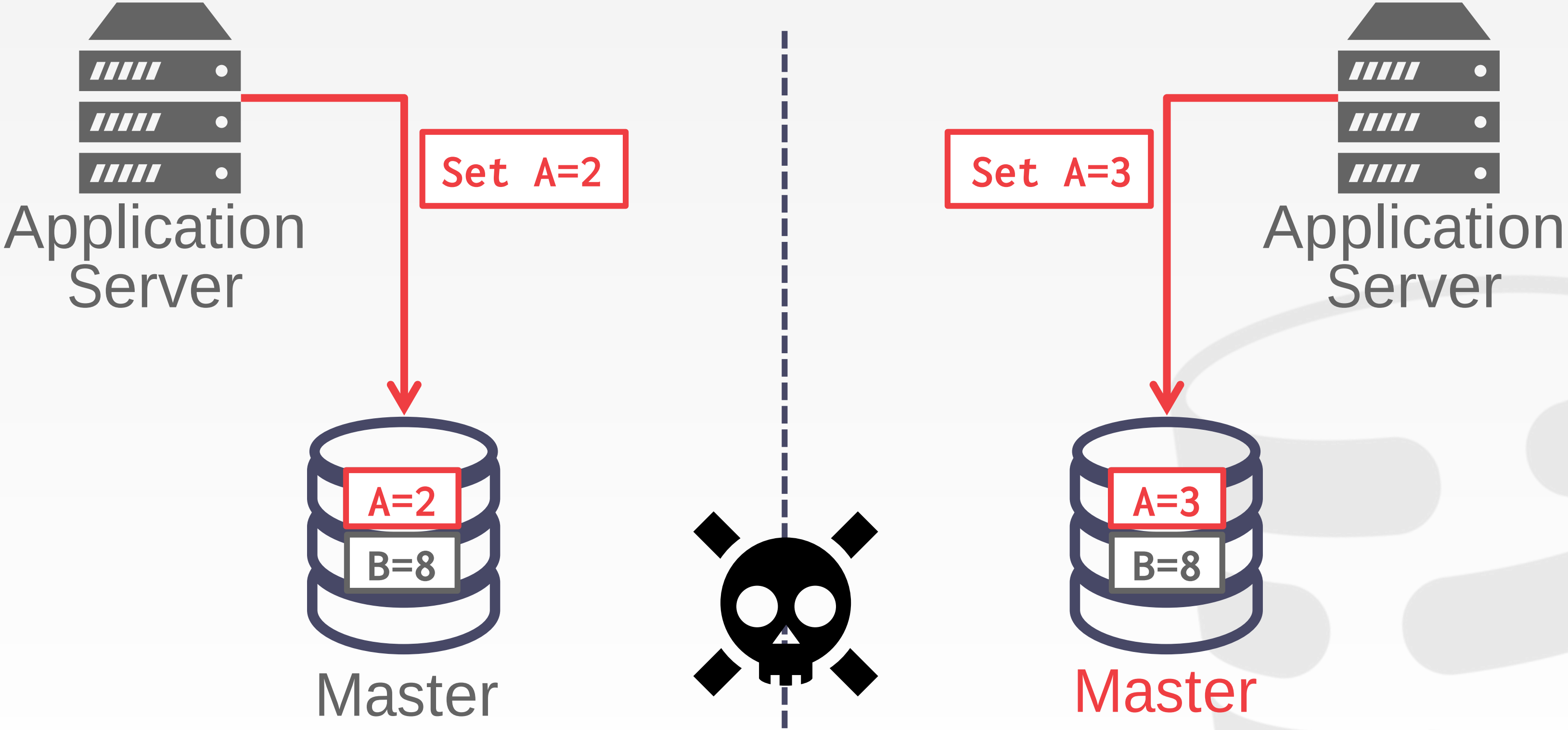

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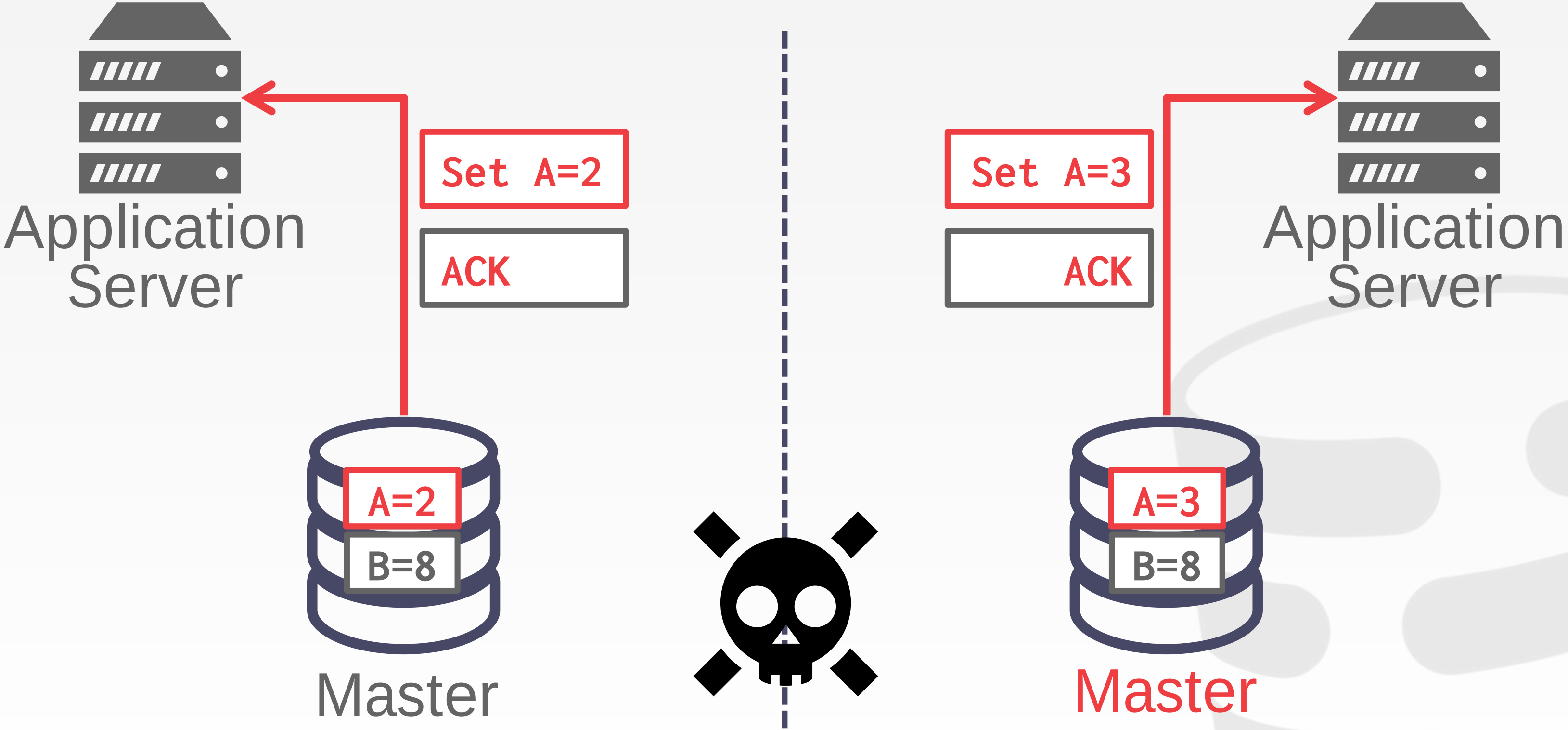
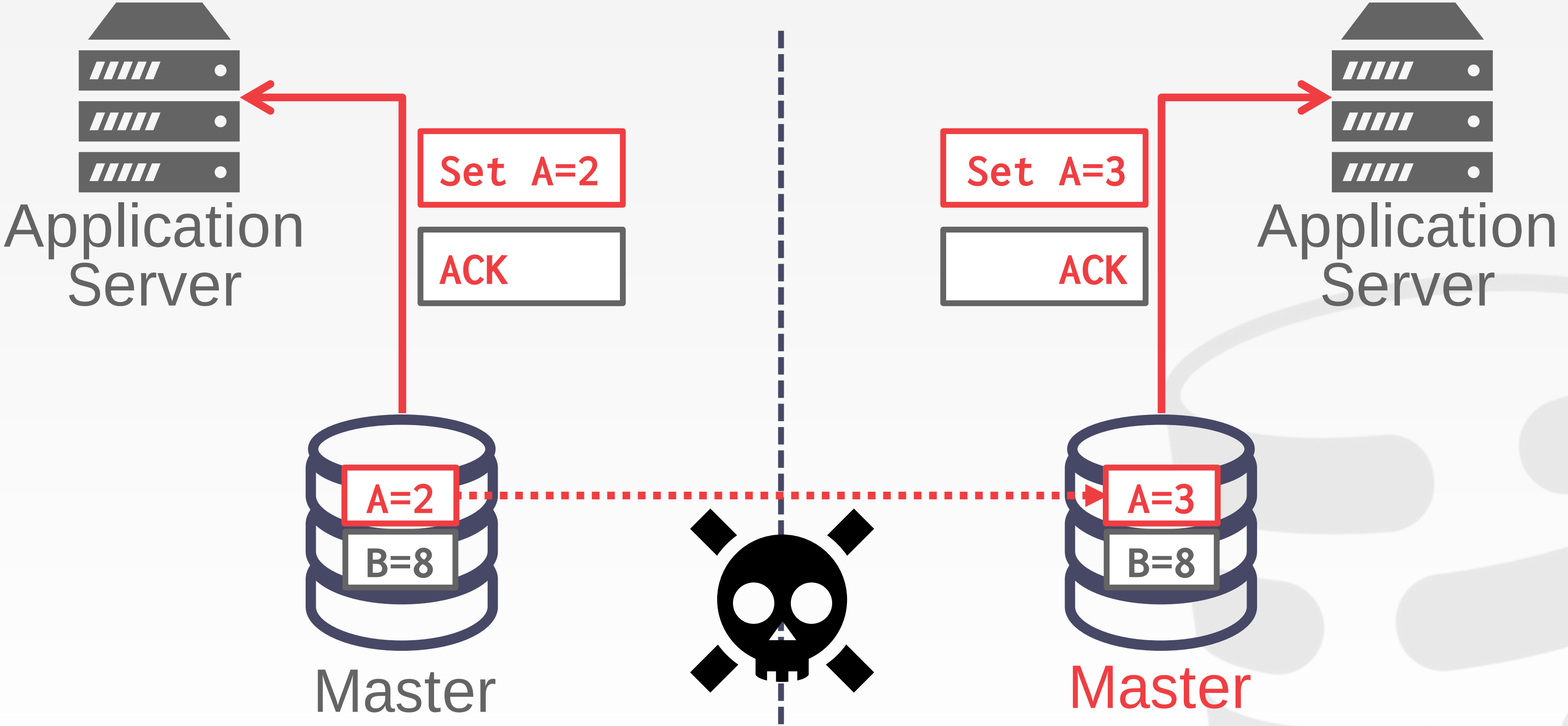


CAP – PARTITION TOLERANCE



CAP FOR OLTP DBMSs

How a DBMS handles failures determines which elements of the CAP theorem they support.

Traditional/NewSQL DBMSs

→ Stop allowing updates until a majority of nodes are reconnected.

NoSQL DBMSs

→ Provide mechanisms to resolve conflicts after nodes are reconnected.

OBSERVATION

We have assumed that the nodes in our distributed systems are running the same DBMS software.

But organizations often run many different DBMSs in their applications.

It would be nice if we could have a single interface for all our data.

FEDERATED DATABASES

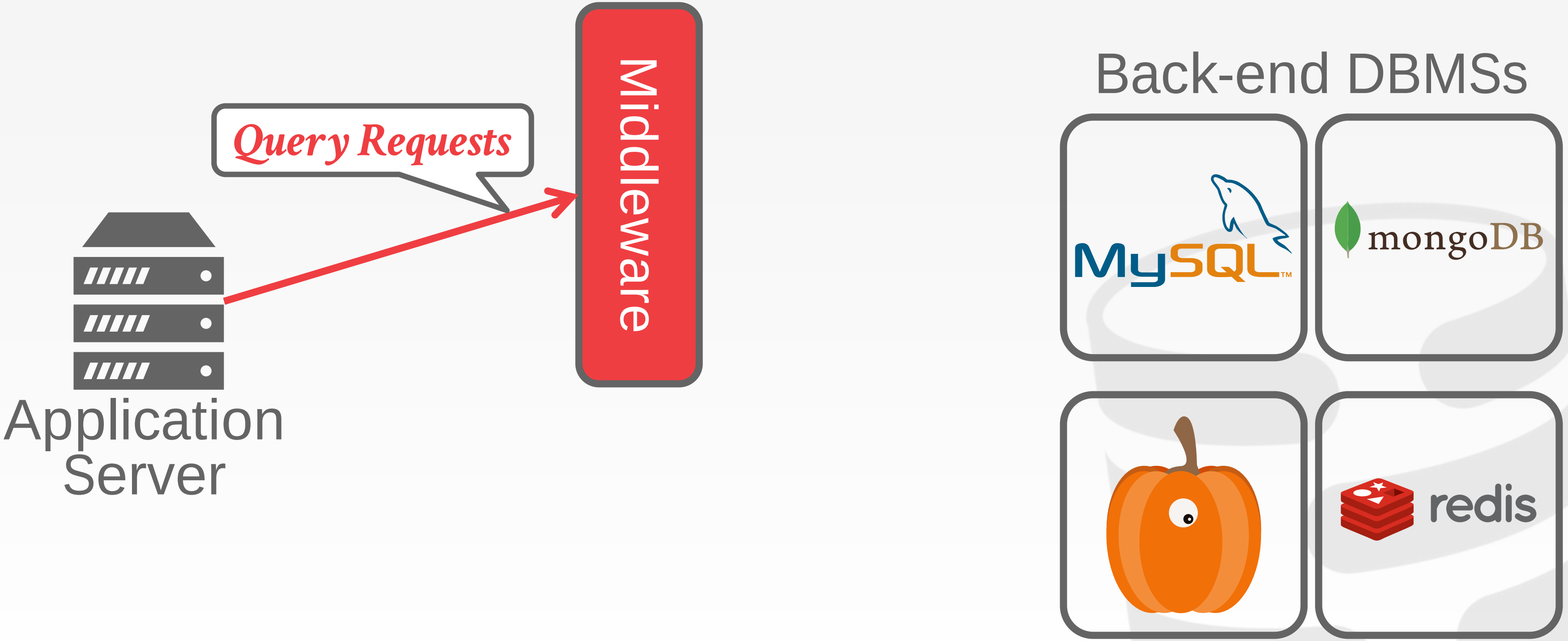
Distributed architecture that connects together multiple DBMSs into a single logical system.

A query can access data at any location.

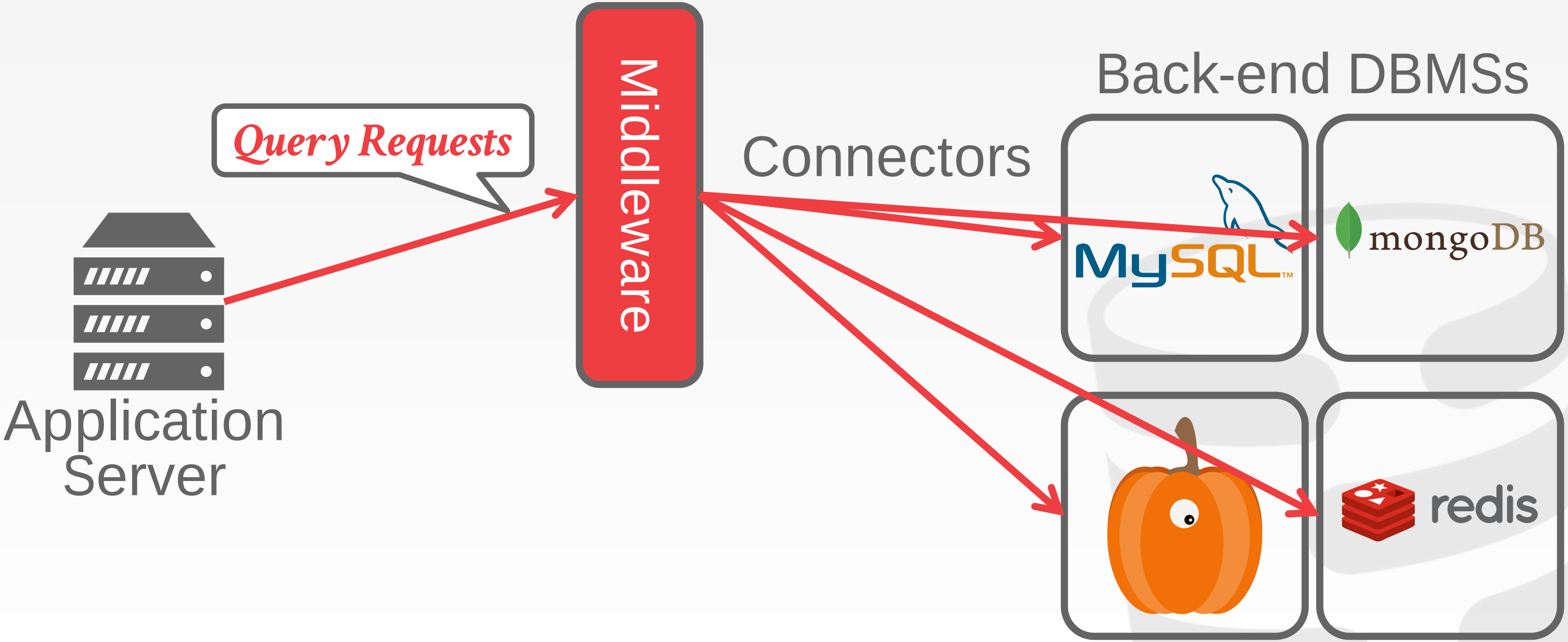
This is hard and nobody does it well

→ Different data models, query languages, limitations.
→ No easy way to optimize queries
→ Lots of data copying (bad).

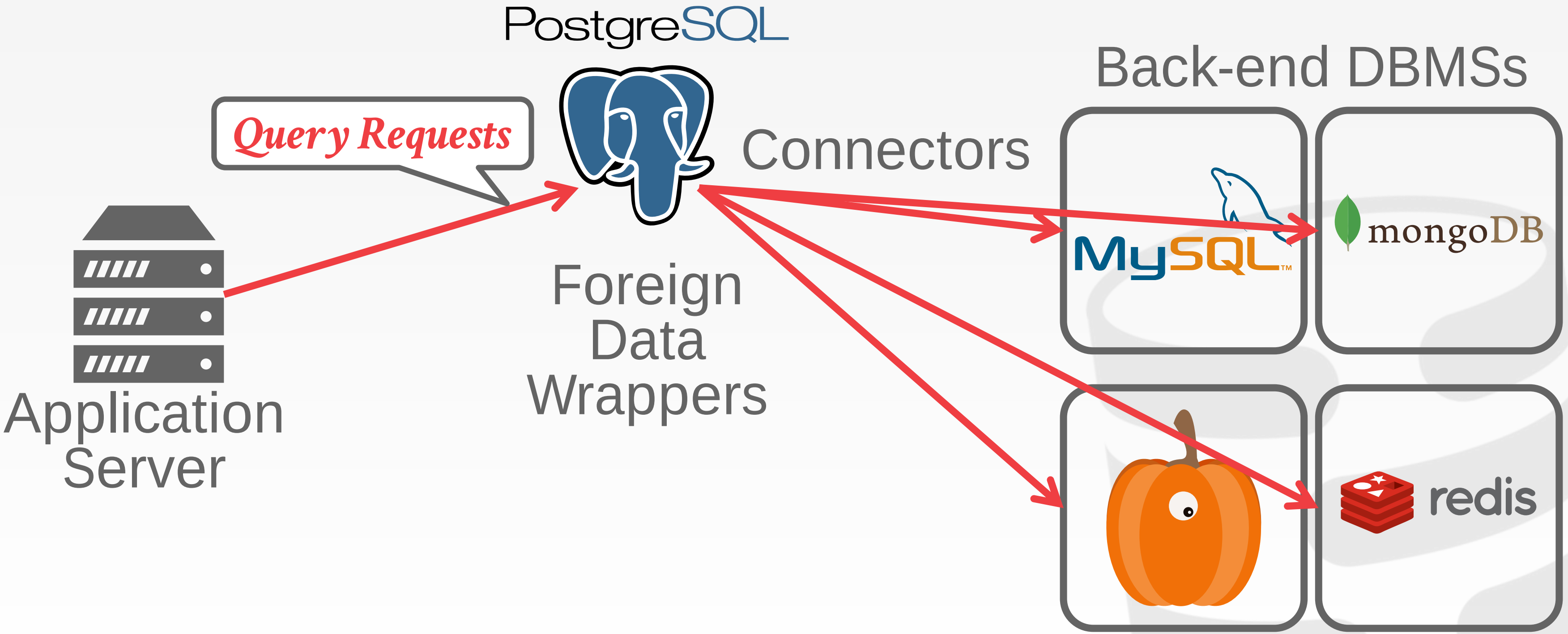
FEDERATED DATABASE EXAMPLE



FEDERATED DATABASE EXAMPLE



FEDERATED DATABASE EXAMPLE



CONCLUSION

We assumed that the nodes in our distributed DBMS are friendly.

Blockchain databases assume that the nodes are adversarial. This means you have to use different protocols to commit transactions.

NEXT CLASS

Distributed OLAP Systems