Homework #4 (by Joyce Liao)
Due: Sunday Nov 13, 2022 @ 11:59pm

IMPORTANT:
- Enter all of your answers into Gradescope by 11:59pm on Sunday Nov 13, 2022.
- Plagiarism: Homework may be discussed with other students, but all homework is to be completed individually.

For your information:
- Graded out of 100 points; 4 questions total
- Rough time estimate: \( \approx 2 - 4 \) hours (0.5 - 1 hours for each question)

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<thead>
<tr>
<th>Question</th>
<th>Points</th>
<th>Score</th>
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<tbody>
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<td>Serializability and 2PL</td>
<td>27</td>
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<td>Deadlock Detection and Prevention</td>
<td>35</td>
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Revision: 2022/10/23 17:34
Question 1: Serializability and 2PL .................... [27 points]

(a) True/False Questions:

i. [2 points] Cascading aborts do not happen under rigorous 2PL.
   - True    False

ii. [2 points] A schedule that is view serializable is also conflict serializable.
    - True    False

iii. [2 points] 2PL is a pessimistic concurrency control protocol.
    - True    False

iv. [2 points] Deadlocks cannot happen in non-rigorous 2PL.
   - True    False

v. [2 points] There are not dirty reads in non-rigorous 2PL.
   - True    False

vi. [2 points] A schedule with an acyclic precedence graph is view serializable.
    - True    False

(b) Serializability:
Consider the schedule given below in Table 1. R(·) and W(·) stand for ‘Read’ and ‘Write’, respectively.

\[\begin{array}{|c|c|c|c|c|c|c|c|c|}
\hline
time & t_1 & t_2 & t_3 & t_4 & t_5 & t_6 & t_7 & t_8 & t_9 & t_{10} \\
\hline
T_1 & & W(E) & R(D) & & & & W(A) & & & \\
\hline
T_2 & W(B) & & R(A) & & & & & & & \\
\hline
T_3 & & & R(E) & W(C) & & & & & & \\
\hline
T_4 & W(D) & & R(B) & R(C) & & & & & & \\
\hline
\end{array}\]

Table 1: A schedule with 4 transactions

i. [2 points] Is this schedule serial?
   - Yes    No

ii. [3 points] Choose the correct dependency graph of the schedule given above. Each edge in the dependency graph looks like this: ‘\(T_x \rightarrow T_y\) with \(Z\) on the arrow indicating that there is a conflict on \(Z\) where \(T_x\) read/wrote on \(Z\) before \(T_y\)’.

\[\begin{array}{c}
\hline
\text{A)} & \\
\hline
\text{B)} & \\
\hline
\end{array}\]
iii. [2 points] Is this schedule conflict serializable?
   □ Yes  □ No

iv. [4 points] What is the minimum number of transactions that need to be removed to produce a conflict serializable schedule?
   □ Two (T1 and T4)
   □ One (T1)
   □ One (T4)
   □ Zero (original schedule is already serializable)

v. [4 points] Which of the following serial schedules is conflict equivalent to the given schedule? Mark all that apply.
   □ T1, T2, T3, T4
   □ T2, T4, T1, T3
   □ T2, T3, T1, T4
   □ None of the above
Question 2: Deadlock Detection and Prevention.................[35 points]

(a) Deadlock Detection:
Consider the following transactions and note that

- S(·) and X(·) stand for ‘shared lock’ and ‘exclusive lock’, respectively.
- \(T_1\) and \(T_2\) represent two transactions.
- Transactions will never release a granted lock until the last operation shown for the transaction is completed.

\(T_1: (a)\) read(A); (b) write(B); (c) read(C); (d) read(B); (e) read(A); (f) write(C);

i. For each position in \(T_1\) above, which lock should be requested?

\(\alpha\) [1 point] At (a): \(\square\) S(A) \(\square\) X(A) \(\square\) No lock needs to be requested

\(\beta\) [1 point] At (b): \(\square\) S(B) \(\square\) X(B) \(\square\) No lock needs to be requested

\(\gamma\) [1 point] At (c): \(\square\) S(C) \(\square\) X(C) \(\square\) No lock needs to be requested

\(\delta\) [1 point] At (d): \(\square\) S(B) \(\square\) X(B) \(\square\) No lock needs to be requested

\(\epsilon\) [1 point] At (e): \(\square\) S(A) \(\square\) X(A) \(\square\) No lock needs to be requested

\(\zeta\) [1 point] At (f): \(\square\) S(C) \(\square\) X(C) \(\square\) No lock needs to be requested

ii. [4 points] Which of the following schedule can cause a deadlock?

\(\square\)

\(\square\)

\(\square\)

(b) Consider the following lock requests in Table 2. And note that

- S(·) and X(·) stand for ‘shared lock’ and ‘exclusive lock’, respectively.
- \(T_1\), \(T_2\), and \(T_3\) represent three transactions.
- \(LM\) stands for ‘lock manager’.
- Transactions will never release a granted lock.

i. For the lock requests in Table 2, determine which lock will be granted or blocked by the lock manager. Please write ‘g’ in the LM row to indicate the lock is granted and ‘b’ to indicate the lock is blocked or the transaction has already been blocked by a former lock request. For example, in the table, the first lock (S(A) at time \(t_1\)) is marked as granted.

\(\square\)

Question 2 continues…
Table 2: Lock requests of three transactions

<table>
<thead>
<tr>
<th>time</th>
<th>$t_1$</th>
<th>$t_2$</th>
<th>$t_3$</th>
<th>$t_4$</th>
<th>$t_5$</th>
<th>$t_6$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$T_1$</td>
<td>S(A)</td>
<td></td>
<td>S(C)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$T_2$</td>
<td></td>
<td>S(B)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$T_3$</td>
<td>X(B)</td>
<td>X(C)</td>
<td>X(A)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>$LM$</td>
<td>g</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

α) [1 point] At $t_2$: □ g □ b

β) [1 point] At $t_3$: □ g □ b

γ) [1 point] At $t_4$: □ g □ b

δ) [1 point] At $t_5$: □ g □ b

e) [1 point] At $t_6$: □ g □ b

ii. [4 points] Mark the correct wait-for graph for the lock requests in Table 2. Each edge in the wait-for graph looks like this: $T_x \rightarrow T_y$ because of $Z$. $Z$ is denoted in the arrow in the figure. (i.e., $T_x$ is waiting for $T_y$ to release its lock on resource $Z$).

[Diagrams of wait-for graphs]

(A)

(B)
iii. [2 points] Determine whether there exists a deadlock in the lock requests in Table 2.

- There is no deadlock
- Cycle (T₁ → T₂ → T₃) exists and schedule deadlocks
- Cycle (T₁ → T₃ → T₁) exists and schedule deadlocks
- Cycle (T₂ → T₃ → T₂) exists and schedule deadlocks
(c) Deadlock Prevention:
Consider the following lock requests in Table 3.
Like before,
• S(·) and X(·) stand for ‘shared lock’ and ‘exclusive lock’, respectively.
• T₁, T₂, T₃, T₄, and T₅ represent five transactions.
• LM represents a ‘lock manager’.
• Transactions will never release a granted lock.

<table>
<thead>
<tr>
<th>time</th>
<th>t₁</th>
<th>t₂</th>
<th>t₃</th>
<th>t₄</th>
<th>t₅</th>
<th>t₆</th>
<th>t₇</th>
<th>t₈</th>
</tr>
</thead>
<tbody>
<tr>
<td>T₁</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>X(C)</td>
</tr>
<tr>
<td>T₂</td>
<td>S(A)</td>
<td>S(B)</td>
<td>S(C)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>T₃</td>
<td>X(A)</td>
<td>X(B)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>T₄</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>X(C)</td>
<td>X(B)</td>
</tr>
<tr>
<td>LM</td>
<td>g</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 3: Lock requests of four transactions

i. To prevent deadlock, we use a lock manager (LM) that adopts the Wait-Die policy. We assume that in terms of priority: T₁ > T₂ > T₃ > T₄. Here, T₁ > T₂ because T₁ is older than T₂ (i.e., older transactions have higher priority). Determine whether the lock request is granted (‘g’), blocked (‘b’), aborted (‘a’), or already dead (‘–’).

α) [1 point] At t₂: □ g □ b □ a □ –
β) [1 point] At t₃: □ g □ b □ a □ –
γ) [1 point] At t₄: □ g □ b □ a □ –
δ) [1 point] At t₅: □ g □ b □ a □ –
ε) [1 point] At t₆: □ g □ b □ a □ –
ζ) [1 point] At t₇: □ g □ b □ a □ –
η) [1 point] At t₈: □ g □ b □ a □ –

Homework #4 continues...
Question 3: Hierarchical Locking ................................. [20 points]

Consider a database (D) consisting of two tables, Release (R) and Artists (A). Specifically,

- Release(rid, name, artist_credit, language, status, genre, year, number_sold), spans 1000 pages, namely $R_1$ to $R_{1000}$
- Artists(id, name, type, area, gender, begin_date_year), spans 50 pages, namely $A_1$ to $A_{50}$

Further, each page contains 100 records, and we use the notation $R_3:20$ to represent the $20^{th}$ record on the third page of the Release table. Similarly, $A_5:10$ represents the $10^{th}$ record on the fifth page of the Artists table.

We use Multiple-granularity locking, with S, X, IS, IX and SIX locks, and four levels of granularity: (1) database-level (D), (2) table-level (R, A), (3) page-level ($R_1 - R_{1000}$, $A_1 - A_{50}$), (4) record-level ($R_1:1 - R_{1000}:100$, $A_1:1 - A_{50}:100$).

For each of the following operations on the database, check all the sequence of lock requests based on intention locks that should be generated by a transaction that wants to efficiently carry out these operations by maximizing concurrency. Please take care of efficiency for e.g., share vs. exclusive lock and granularity.

Please follow the format of the examples listed below:

- mark “IS(D)” for a request of database-level IS lock
- mark “X($A_2:30$)” for a request of record-level X lock for the $30^{th}$ record on the second page of the Artists table
- mark “S($A_2:30 - A_3:100$)” for a request of record-level S lock from the $30^{th}$ record on the second page of the Artists table to the $100^{th}$ record on the third page of the Artists table.

(a) [4 points] Fetch the record that represents the artist with the earliest begin_date_year. Assume there are no ties.
- IX(D), X(A)
- IS(D), S(A)
- S(D)
- IS(D), SIX(A), X(A)

(b) [4 points] Modify the $4^{th}$ record on $R_{45}$.
- IX(D), IX(R), IX($R_{45}$), X($R_{45}:4$)
- IX(D), SIX(R), X($R_{45}$)
- IX(D), IX(R), SIX($R_{45}$), X($R_{45}:4$)
- IS(D), IS(R), IX($R_{45}$), X($R_{45}:4$)

(c) [4 points] Find the average number_sold across all releases.
- S(D)
- S(D), S(R)
- IS(D), X(R)
- IS(D), S(R)

Question 3 continues...
(d) [4 points] Scan all records in R and modify the 1st record on $R_5$
   - IX(D), SIX(R), IX($R_5$), X($R_5 : 1$)
   - IS(D), S(R), IX($R_5 : 1$)
   - S(D), IX(R), X($R_5 : 1$)
   - IS(D), SIX(R), X($R_5$)

(e) [4 points] Update the name of all artists to be in uppercase letters only.
   - SIX(D), X(A)
   - SIX(D), S(A)
   - IX(D), X(A)
   - IX(D), SIX(A)
Question 4: Optimistic Concurrency Control .................. [18 points]
Consider the following set of transactions accessing a database with object A, B, C, D. The questions below assume that the transaction manager is using optimistic concurrency control (OCC). Assume that a transaction switches from the READ phase immediately into the VALIDATION phase after its last operation executes.

Note: VALIDATION may or may not succeed for each transaction. If validation fails, the transaction will get immediately aborted.

You can assume that the DBMS is using the serial validation protocol discussed in class where only one transaction can be in the validation phase at a time, and each transaction is doing backward validation (i.e. Each transaction, when validating, checks whether it intersects its read/write sets with any transactions that have already committed. You may assume there are no other transactions in addition to the ones shown below. )

<table>
<thead>
<tr>
<th>time</th>
<th>T₁</th>
<th>T₂</th>
<th>T₃</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>READ(A)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>READ(C)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>WRITE(A)</td>
<td>READ(B)</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>WRITE(C)</td>
<td>validators?</td>
<td>READ(B)</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>7</td>
<td>validators?</td>
<td>READ(D)</td>
<td></td>
</tr>
<tr>
<td>8</td>
<td>WRITE?</td>
<td></td>
<td></td>
</tr>
<tr>
<td>9</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>10</td>
<td>WRITE(D)</td>
<td>WRITE(B)</td>
<td></td>
</tr>
<tr>
<td>11</td>
<td>WRITE(B)</td>
<td>validators?</td>
<td>READ(A)</td>
</tr>
<tr>
<td>12</td>
<td>validators?</td>
<td>WRITE?</td>
<td></td>
</tr>
<tr>
<td>13</td>
<td>WRITE(A)</td>
<td>WRITE(B)</td>
<td>validators?</td>
</tr>
<tr>
<td>14</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>15</td>
<td></td>
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<td>16</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>17</td>
<td></td>
<td></td>
<td>validators?</td>
</tr>
<tr>
<td>18</td>
<td></td>
<td></td>
<td>WRITE?</td>
</tr>
</tbody>
</table>

Figure 1: An execution schedule

(a) [4 points] Will T₁ abort?
- Yes
- No

(b) [4 points] Will T₂ abort?
- Yes
- No

(c) [4 points] Will T₃ abort?
- Yes
- No

Question 4 continues…
(d) [2 points] OCC works best when concurrent transactions access the same subset of data in a database.
   □ True □ False

(e) [2 points] Transactions can suffer from *phantom reads* in OCC.
   □ True □ False

(f) [2 points] Aborts due to OCC are wasteful because they happen after a transaction has already finished executing.
   □ True □ False