# Carnegie Mellon University <br> Computer Science Department <br> 15-445/645 - Database Systems (Fall 2023) Prof. Andy Pavlo and Jignesh Patel 

Homework \#2 (by Wiam Eddahri)
Due: Sunday September 242023 @ 11:59pm

## IMPORTANT:

- Enter all of your answers into Gradescope by 11:59pm on Sunday September 242023.
- Plagiarism: Homework may be discussed with other students, but all homework is to be completed individually.
For your information:
- Graded out of $\mathbf{1 0 0}$ points; $\mathbf{3}$ questions total
- Rough time estimate: $\approx 4-6$ hours (1-1.5 hours for each question)

Revision : 2023/09/15 10:33

| Question | Points | Score |
| :---: | :---: | :---: |
| Storage Models | 16 |  |
| Extendible Hashing | 19 |  |
| B+Tree | 65 |  |
| Total: | 100 |  |

## Question 1: Storage Models

Consider a database with a single table T (andrew_id, full_name, graduation_year, department, gpa), where andrew_id is the primary key, and all attributes are the same fixed width. Suppose T has 5,000 tuples that fit into 500 pages, Ignore any additional storage overhead for the table (e.g., page headers, tuple headers). Additionally, you should make the following assumptions:

- The DBMS does not have any additional meta-data (e.g., sort order, zone maps).
- T does not have any indexes (including for primary key andrew_id)
- None of T's pages are already in the buffer pool.
- Content-wise, the tuples of T will make each query run the longest possible (this assumption is critical for solving part (a))
- The tuples of T can be in any order (this assumption is critical for solving part (b) when you compute the minimum versus maximum number of pages that the DBMS will potentially have to read)
(a) Consider the following query:

```
SELECT MAX(gpa) FROM T
    WHERE graduation_year == 2025;
```

i. [4 points] Suppose the DBMS uses the decomposition storage model (DSM) with implicit offsets. How many pages will the DBMS potentially have to read from disk to answer this query? (Keep in mind our assumption about the contents of T !)1-100 $\square$ 101-200201-300301-500$\geq 501$
$\square$ Not possible to determine
ii. [4 points] Suppose the DBMS uses the N-ary storage model (NSM). How many pages will the DBMS potentially have to read from disk to answer this query? (Keep in mind our assumption about the contents of T !)1-200
201-300301-400 401-500$\geq 501$Not possible to determine
(b) Now consider the following query:

SELECT andrew_id, gpa FROM T
WHERE gpa = (SELECT MAX(gpa) FROM T);
i. Suppose the DBMS uses the decomposition storage model (DSM) with implicit offsets.
$\alpha$ ) [4 points] What is the minimum number of pages that the DBMS will potentially have to read from disk to answer this query?2-5100-200
201-299 $\square \geq 300$
Not possible to
determine
$\beta$ ) [4 points] What is the maximum number of pages that the DBMS will potentially have to read from disk to answer this query?
$\square 1$
2-5100-200201-299$\geq 300$Not possible to determine

## Question 2: Extendible Hashing

Consider an extendible hashing structure such that:

- Each bucket can hold up to four records.
- The hashing function uses the lowest $g$ bits, where $g$ is the global depth.
- A new extendible hashing structure is initialized with $g=0$ and one empty bucket
(a) Starting from an empty table, insert keys $0,1,4,5$.
i. [2 points] What is the global depth of the resulting table?
$\square 0$1234 None of the above
ii. [2 points] What is the local depth of the bucket containing 5?0123None of the above
(b) Starting from the result in (a), you insert keys $10,11,12,13$.
i. [2 points] What is the global depth of the resulting table?01 2$\square 4$None of the above
ii. [2 points] What is the local depth of the bucket containing 11 ?01234 None of the above
(c) Starting from the result in (c), you insert keys 2 and 14.
i. [2 points] What is the global depth of the resulting table?01234None of the above
ii. [3 points] Starting from the result in (c), you insert keys 15, 3 and 7, what is the global depth?
01 234 None of the above
iii. [3 points] Which value, if inserted, will hash to the same bucket as the bucket containing key 1 ?$3 \square 7$11All of the aboveNone of the above
(d) [3 points] Starting from the result in (c), which key(s), if inserted next, will cause a split that doubles the table's size?
$\square 33$
6429 61All of the aboveNone of the above


## Question 3: B+Tree

Consider the following B+tree.


Figure 1: B+ Tree of order $d=5$ and height $h=3$.

When answering the following questions, be sure to follow the procedures described in class and in your textbook. You can make the following assumptions:

- A left pointer in an internal node guides towards keys $<$ than its corresponding key, while a right pointer guides towards keys $\geq$.
- A leaf node underflows when the number of keys goes below $\left\lceil\frac{d-1}{2}\right\rceil$.
- An internal node underflows when the number of pointers goes below $\left\lceil\frac{d}{2}\right\rceil$.
- You should always consider redistribution before trying to merge two nodes.

Note that B+ tree diagrams for this problem omit leaf pointers for convenience. The leaves of actual $\mathrm{B}+$ trees are linked together via pointers, forming a singly linked list allowing for quick traversal through all keys.
(a) [5 points] How many tree nodes must be fetched to answer the following query: Get all records with search key smaller than 42.
$\square 6$$7 \quad \square$ 4None
(b) [5 points] Insert $40^{*}$ into the $\mathrm{B}+$ tree. Select the resulting tree.
$\square$ A)

$\square$ B)

$\square$ C)

$\square$ D)

(c) [5 points] Starting with the intial tree, delete $49^{*}$. Select the resulting tree.
$\square$ A)

B)

$\square$ C)

$\square$ D)

(d) [5 points] Starting with the intial tree, delete $62^{*}$ and delete $68^{*}$. Select the resulting tree.
$\square$ A)

$\square$ B)


## $\square$ C)


$\square$ D)



Figure 2: $\mathrm{B}+$ tree with violations
(e) The $\mathrm{B}+$ Tree shown in Figure 2 is invalid. That is, its nodes violate the correctness properties of $\mathrm{B}+$ Trees that we discussed in class. If the tree is invalid, select all the properties that are violated for each node. If the node is valid, then select 'None'. There will be no partial credit for missing violations.
Note:

- If a node's subtrees are not the same height, the balance property is violated at that node only.
- If a node's subtrees contain values not in the range specified by the node's separator keys, the separator keys property is violated at that node.
i. [2 points] Which properties are violated by Leaf 3?


## Key order property <br> Half-full property <br> Balance property

Separator keys
None
ii. [2 points] Which properties are violated by Leaf 1?Key order property
Half-full property Balance property
$\square$ Separator keys
None
iii. [2 points] Which properties are violated by Internal Node 1?

Key order property $\quad \square$ Half-full property $\quad \square$ Balance property
Separator keys $\square$ None
iv. [2 points] Which properties are violated by Internal Node 2?

Key order property $\quad \square$ Half-full property
Balance property
$\square$ Separator keys None
v. [2 points] Which properties are violated by Root?

Key order property $\quad \square$ Half-full property $\quad \square$ Balance propertySeparator keys None
(f) i. [5 points] A DBMS may potentially use separate buffer pools for a B+Tree's inner node pages and for its leaf node pages.TrueFalse
ii. [5 points] A read-only thread needs to hold at most two latches at the same time. TrueFalse
iii. [5 points] A write thread needs to hold at most three write latches at the same timeTrueFalse
iv. [5 points] A write thread might hold only one write latch (Consider a tree that has more than one leaf node)TrueFalse
v. [5 points] A read thread must release its latches in the order they were acquired (i.e., FIFO) to prevent concurrency errors.TrueFalse
vi. [5 points] A write thread must release its latches in the order they were acquired (i.e., FIFO) to prevent concurrency errors.TrueFalse
vii. [5 points] The minimum space utilization for a $B+$ tree index is 50 percent.TrueFalse

