Intro to Database Systems (15-445/645)

10 Sorting & Aggregations

Carnegie Mellon University

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ADMINISTRIVIA

Project 1 due last night.

Homework 3 available today, due Sunday, Feb 26th.

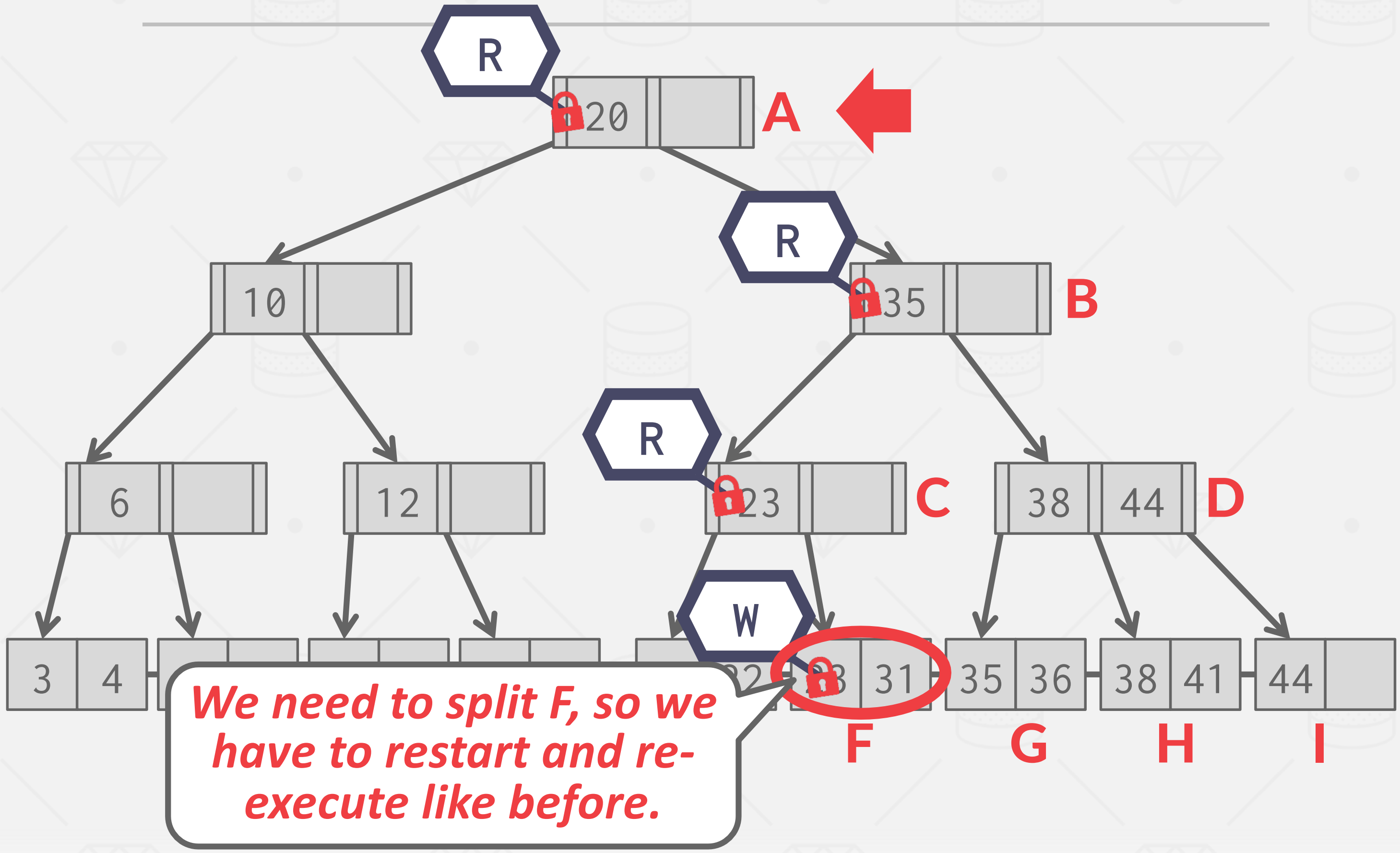
Midterm exam Wednesday, March 1st.

Project 2 available today.
→ First checkpoint due Friday, March 3rd.
→ Overall due Wednesday, March 22nd.

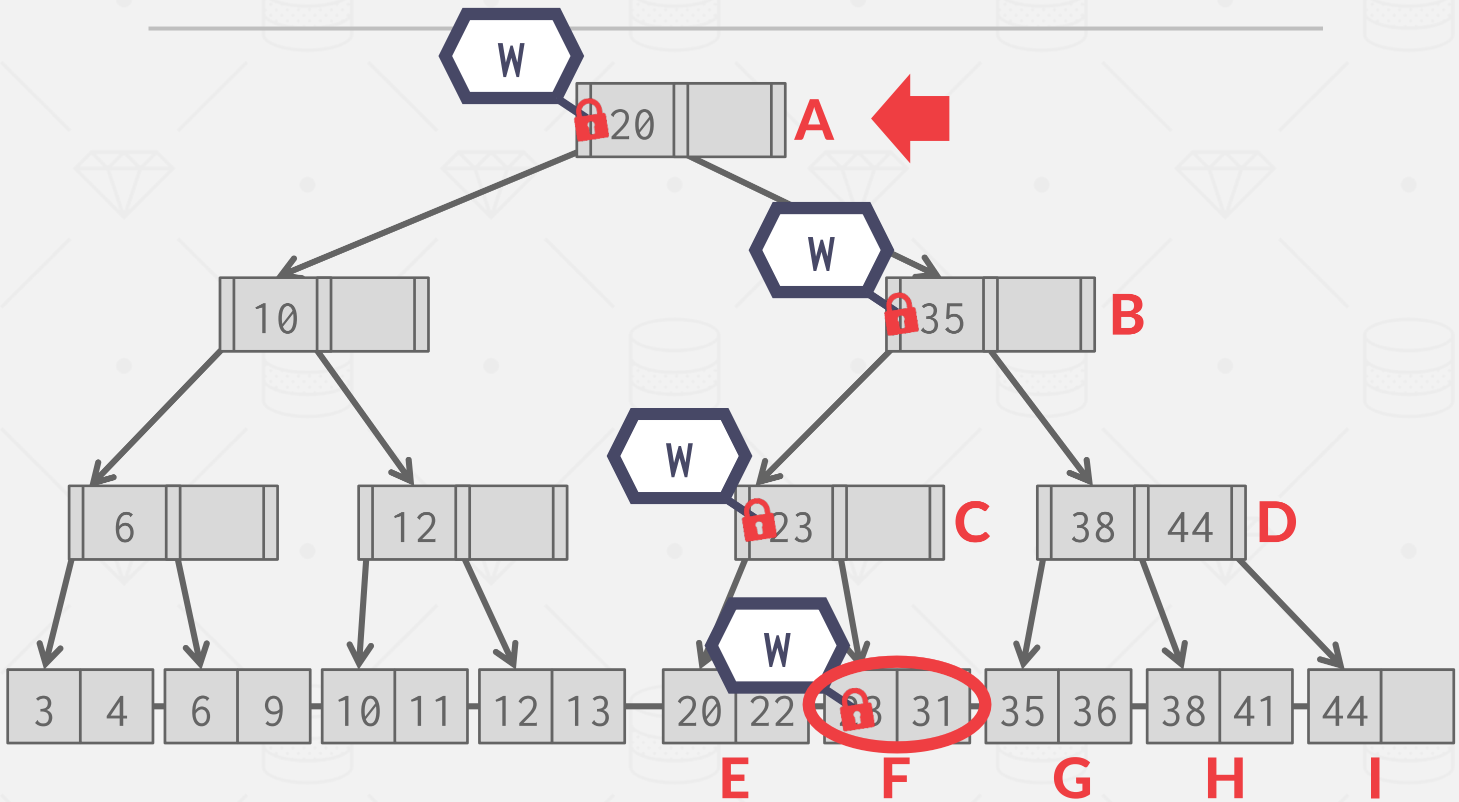
LAST TIME

Concurrent indexes
→ Concurrent hash tables
→ Concurrent B+Trees

EXAMPLE #4 – INSERT 25



EXAMPLE #4 – INSERT 25



OBSERVATION

The threads in all the examples so far have acquired latches in a "top-down" manner.

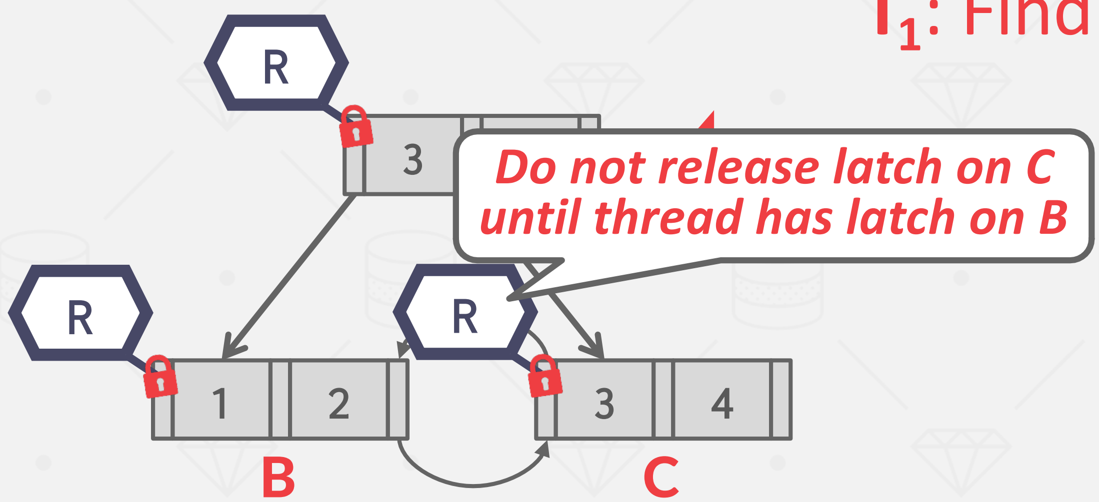
→ A thread can only acquire a latch from a node that is below its current node.

→ If the desired latch is unavailable, the thread must wait until it becomes available.

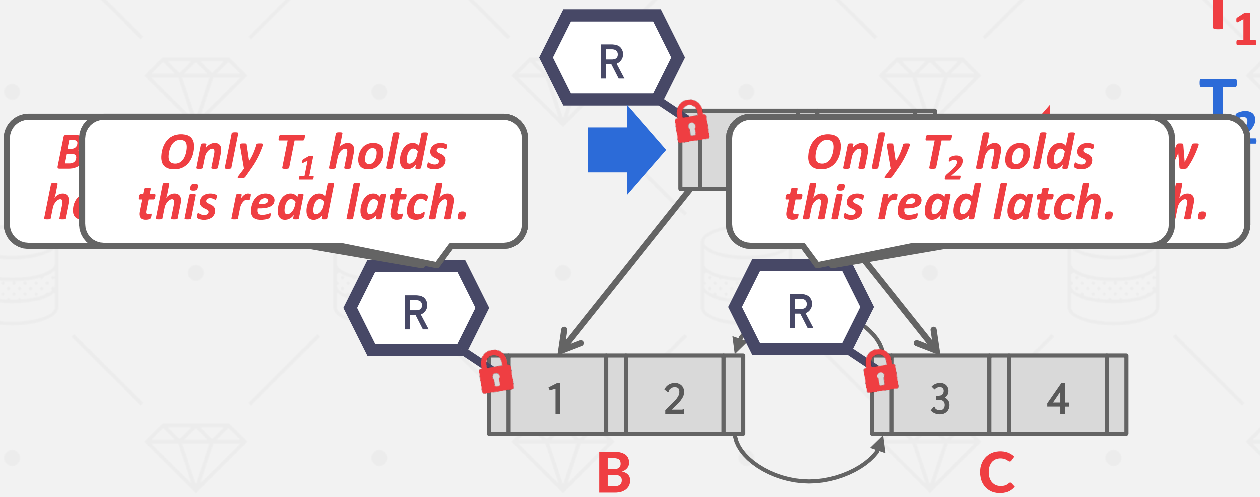
But what if threads want to move from one leaf node to another leaf node?

LEAF NODE SCAN EXAMPLE #1

T_1: Find Keys < 4



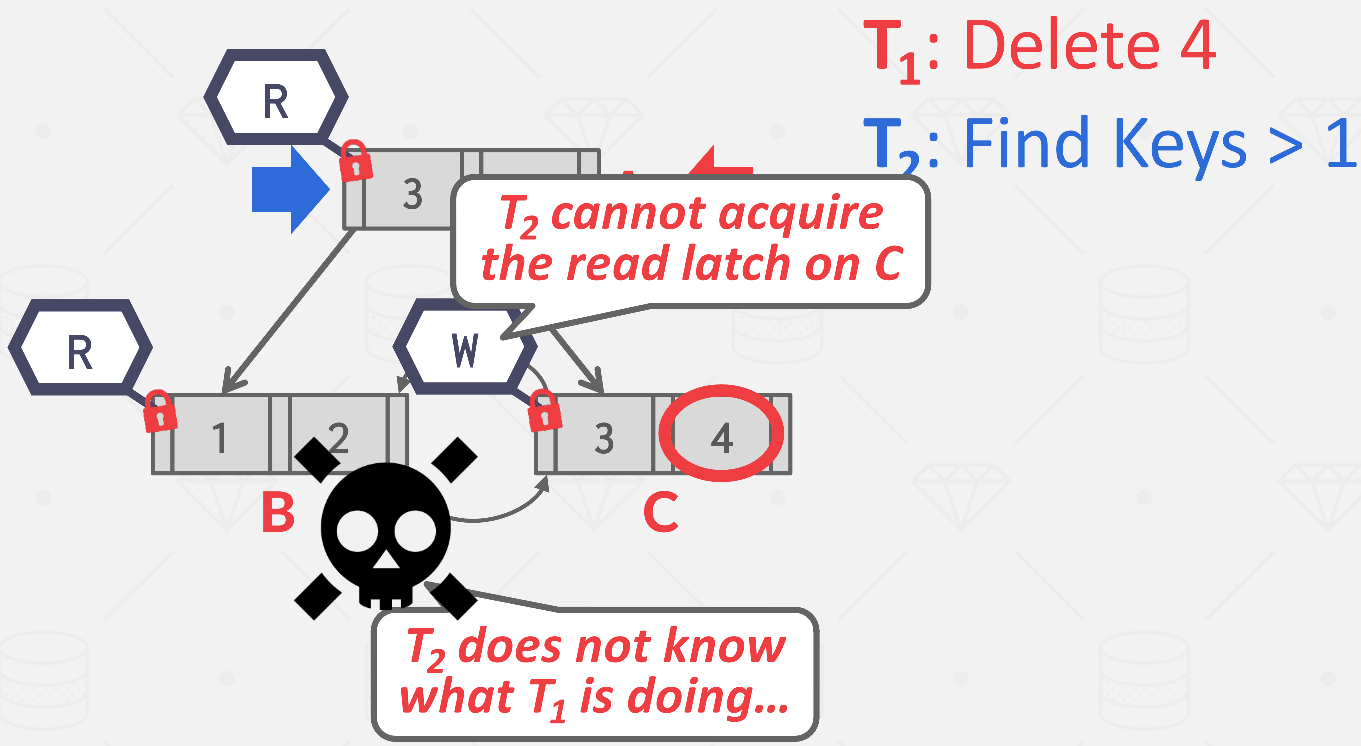
LEAF NODE SCAN EXAMPLE #2



T_1: Find Keys < 4

T_2: Find Keys > 1

LEAF NODE SCAN EXAMPLE #3



LEAF NODE SCANS

Latches do not support deadlock detection or avoidance. The only way we can deal with this problem is through coding discipline.

The leaf node sibling latch acquisition protocol must support a “no-wait” mode.

The DBMS's data structures must cope with failed latch acquisitions.

INDEX CONCURRENCY CONCLUSION

Making a data structure thread-safe is notoriously difficult in practice.

We focused on B+Trees, but the same high-level techniques are applicable to other data structures.

COURSE STATUS

We are now going to talk about how to execute queries using the DBMS components we have discussed so far.

Next four lectures:
→ Operator Algorithms
→ Query Processing Models
→ Runtime Architectures

Query Planning

Operator Execution

Access Methods

Buffer Pool Manager

Disk Manager

QUERY PLAN

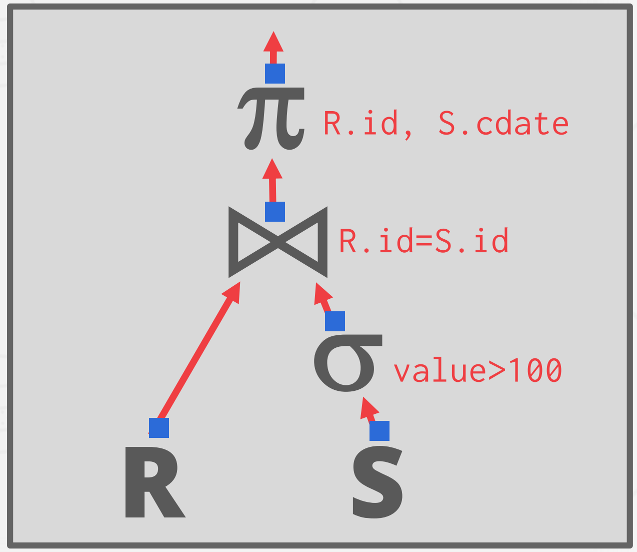
The operators are arranged in a tree.

Data flows from the leaves of the tree up towards the root.

→ We will discuss the granularity of the data movement next week.

The output of the root node is the result of the query.

```
SELECT R.id, S.cdate
  FROM R JOIN S
    ON R.id = S.id
 WHERE S.value > 100
```



DISK-ORIENTED DBMS

Just like it cannot assume that a table fits entirely in memory, a disk-oriented DBMS cannot assume that query results fit in memory.

We will use the buffer pool to implement algorithms that need to spill to disk.

We prefer algorithms that minimize I/O and maximize sequential I/O.

WHY DO WE NEED TO SORT?

Relational model/SQL is unsorted.

Queries may request that tuples are sorted in a specific way (`ORDER BY`).

But even if a query does not specify an order, we may still want to sort to do other things:

→ Trivial to support duplicate elimination (`DISTINCT`)
→ Bulk loading sorted tuples into a B+Tree index is faster
→ Aggregations (`GROUP BY`)

IN-MEMORY SORTING

If data fits in memory, then we can use a standard sorting algorithm like quicksort.

If data does not fit in memory, then we need to use a technique that is aware of the cost of reading and writing disk pages…

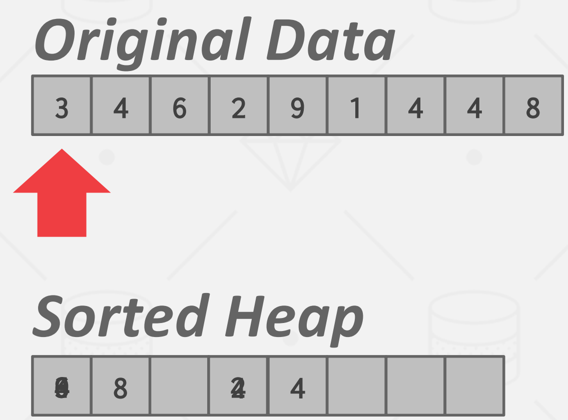
TODAY'S AGENDA

Top-K Heap Sort
External Merge Sort
Aggregations

TOP-K HEAP SORT

If a query contains an **ORDER BY** with a **LIMIT**, then the DBMS only needs to scan the data once to find the top-k elements.

```
SELECT * FROM enrolled
 ORDER BY sid
 FETCH FIRST 4 ROWS
  WITH TIES
```



Heapsort is great if the top-k elements fit in memory.

→ Scan data once, maintain an in-memory sorted priority queue.

EXTERNAL MERGE SORT

Divide-and-conquer algorithm that splits data into separate *runs*, sorts them individually, and then combines them into longer sorted runs.

Phase #1 – Sorting

→ Sort chunks of data that fit in memory and then write back the sorted chunks to a file on disk.

Phase #2 – Merging

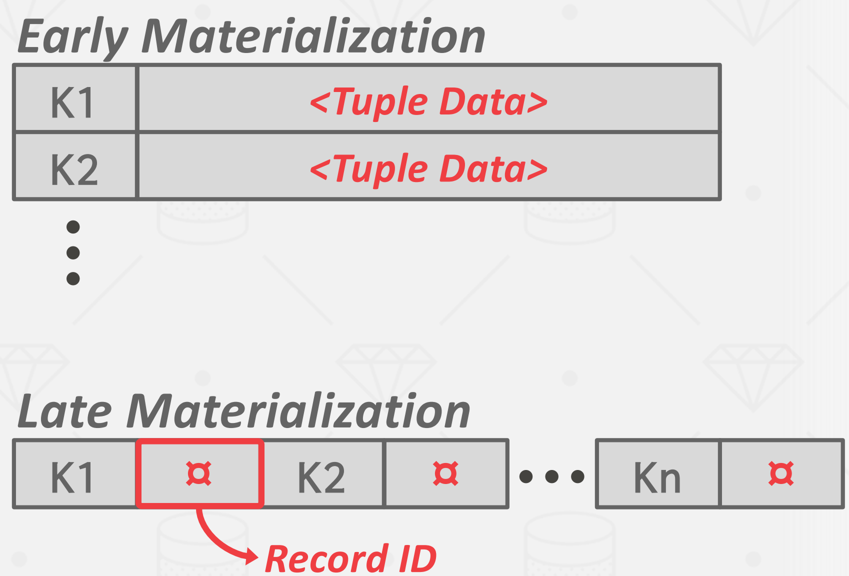
→ Combine sorted runs into larger chunks.

SORTED RUN

A run is a list of key/value pairs.

Key: The attribute(s) to compare to compute the sort order.

Value: Two choices
→ Tuple (***early materialization***).
→ Record ID (***late materialization***).



2-WAY EXTERNAL MERGE SORT

We will start with a simple example of a 2-way external merge sort.

→ "2" is the number of runs that we are going to merge into a new run for each pass.

Data is broken up into ***N*** pages.

The DBMS has a finite number of ***B*** buffer pool pages to hold input and output data.

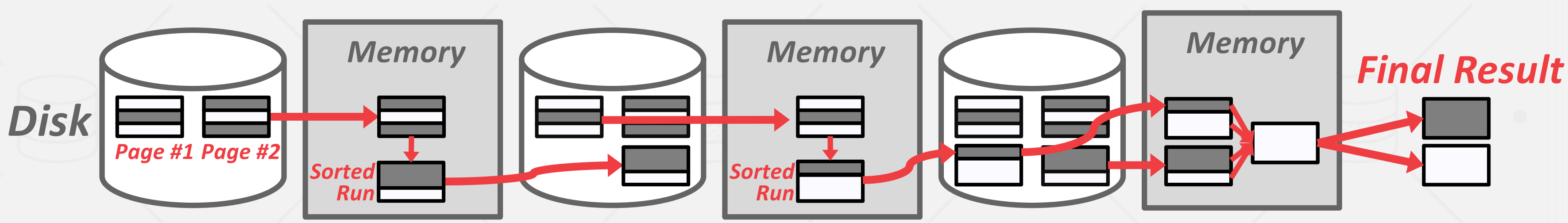
SIMPLIFIED 2-WAY EXTERNAL MERGE SORT

Pass #0

→ Read one page of the table into memory
→ Sort page into a run and write it back to disk
→ Repeat until the whole table has been sorted into runs

Pass #1,2,3,…

→ Recursively merge pairs of runs into runs twice as long
→ Need at least 3 buffer pages (2 for input, 1 for output)

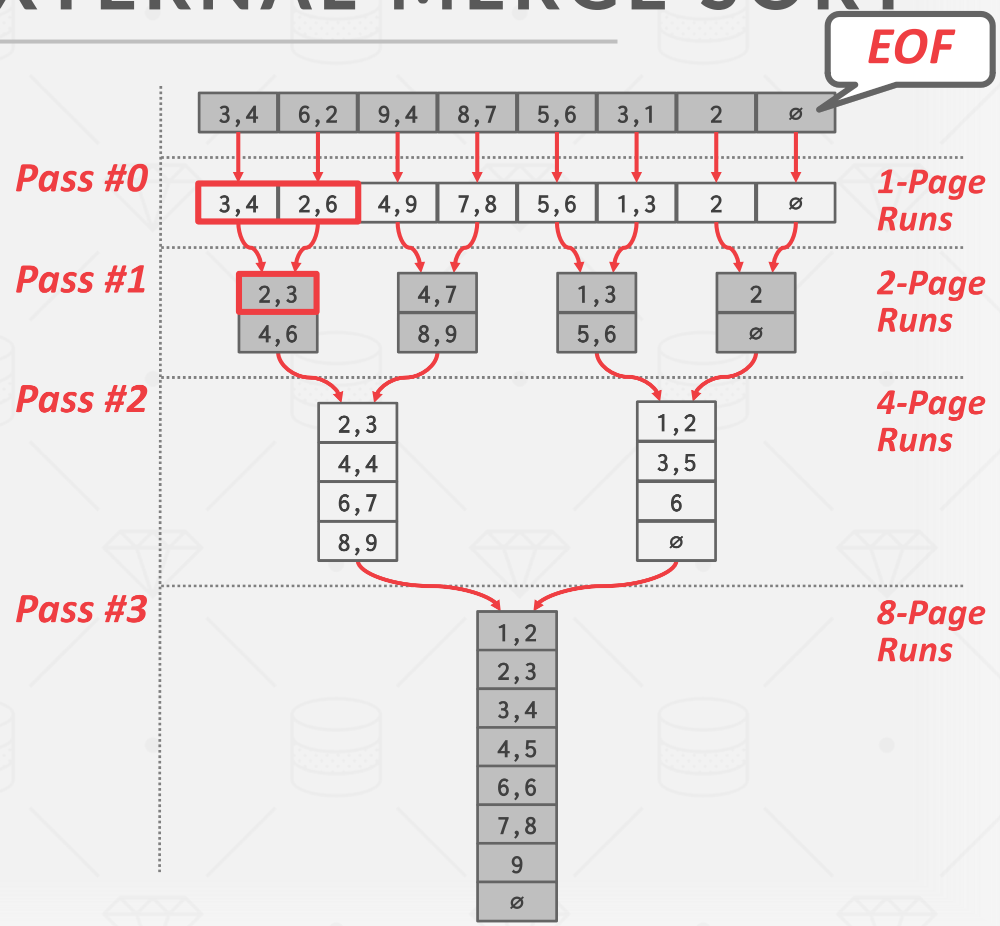


SIMPLIFIED 2-WAY EXTERNAL MERGE SORT

In each pass, we read and write every page in the file.

Number of passes
$= 1 + \lceil \log_2 N \rceil$

Total I/O cost
$= 2N \cdot (\text{\# of passes})$



SIMPLIFIED 2-WAY EXTERNAL MERGE SORT

This simplified algorithm only requires three buffer pool pages to perform the sorting (***B*=3**).

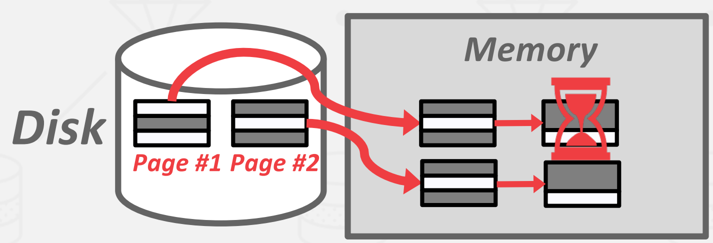
→ Two input pages, one output page

But even if we have more buffer space available (***B*>3**), it does not effectively use them if the worker must block on disk I/O…

DOUBLE BUFFERING OPTIMIZATION

Prefetch the next run in the background and store it in a second buffer while the system is processing the current run.

→ Reduces the wait time for I/O requests at each step by continuously utilizing the disk.



GENERAL EXTERNAL MERGE SORT

Pass #0

→ Use B buffer pages

→ Produce $\lceil N / B \rceil$ sorted runs of size B

Pass #1,2,3,…

→ Merge $B\text{-}1$ runs (i.e., K-way merge)

Number of passes = $1 + \lceil \log_{B-1} \lceil N / B \rceil \rceil$

Total I/O Cost = $2N \cdot$ (# of passes)

EXAMPLE

Determine how many passes it takes to sort 108 pages with 5 buffer pool pages: ***N*=108**, ***B*=5**

→ **Pass #0:** $\lceil N / B \rceil = \lceil 108 / 5 \rceil$ = 22 sorted runs of 5 pages each (last run is only 3 pages).

→ **Pass #1:** $\lceil N' / B\text{-}1 \rceil = \lceil 22 / 4 \rceil$ = 6 sorted runs of 20 pages each (last run is only 8 pages).

→ **Pass #2:** $\lceil N'' / B\text{-}1 \rceil = \lceil 6 / 4 \rceil$ = 2 sorted runs, first one has 80 pages and second one has 28 pages.

→ **Pass #3:** Sorted file of 108 pages.

$$1+\lceil \log_{B-1}\lceil N / B \rceil \rceil = 1+\lceil \log_4 22 \rceil$$
$$= 1+\lceil 2.229... \rceil$$
$$= \mathbf{4\ passes}$$

USING B+TREES FOR SORTING

If the table that must be sorted already has a B+Tree index on the sort attribute(s), then we can use that to accelerate sorting.

Retrieve tuples in desired sort order by simply traversing the leaf pages of the tree.

Cases to consider:

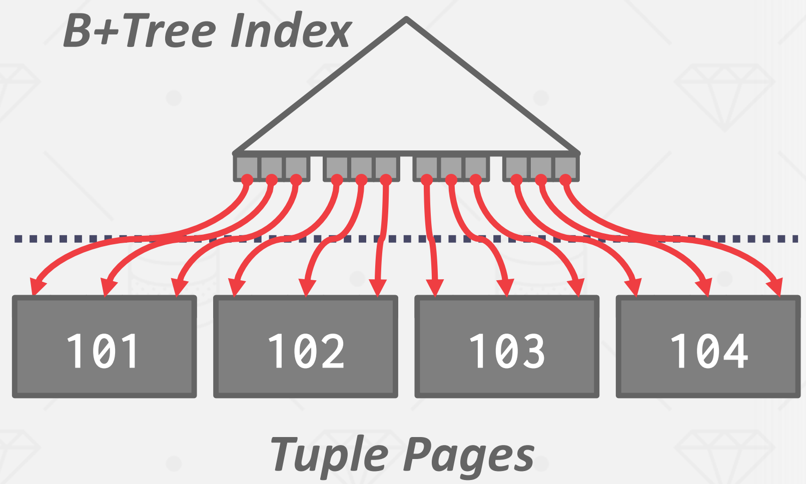
→ Clustered B+Tree

→ Unclustered B+Tree

CASE #1 – CLUSTERED B+TREE

Traverse to the left-most leaf page, and then retrieve tuples from all leaf pages.

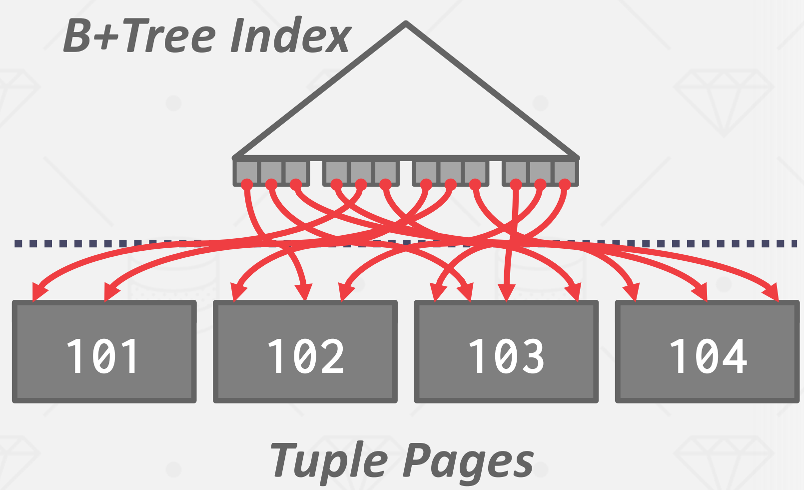
This is always better than external sorting because there is no computational cost, and all disk access is sequential.



CASE #2 – UNCLUSTERED B+TREE

Chase each pointer to the page that contains the data.

This is almost always a bad idea.
In general, one I/O per data record.



AGGREGATIONS

Collapse values for a single attribute from multiple tuples into a single scalar value.

The DBMS needs a way to quickly find tuples with the same distinguishing attributes for grouping.

Two implementation choices:
→ Sorting
→ Hashing

SORTING AGGREGATION

```
SELECT DISTINCT cid
  FROM enrolled
 WHERE grade IN ('B','C')
 ORDER BY cid
```

enrolled(sid,cid,grade)

sid	cid	grade
53666	15-445	C
53688	15-721	A
53688	15-826	B
53666	15-721	C
53655	15-445	C

Filter

sid	cid	grade
53666	15-445	C
53688	15-826	B
53666	15-721	C
53655	15-445	C

Remove Columns

cid
15-445
15-826
15-721
15-445

Sort

cid
15-445
~~15-445~~
15-721
15-826

Eliminate Dupes

ALTERNATIVES TO SORTING

What if we do not need the data to be ordered?
→ Forming groups in **GROUP BY** (no ordering)
→ Removing duplicates in **DISTINCT** (no ordering)

Hashing is a better alternative in this scenario.
→ Only need to remove duplicates, no need for ordering.
→ Can be computationally cheaper than sorting.

HASHING AGGREGATE

Populate an ephemeral hash table as the DBMS scans the table. For each record, check whether there is already an entry in the hash table:

→ `DISTINCT`: Discard duplicate

→ `GROUP BY`: Perform aggregate computation

If everything fits in memory, then this is easy.

If the DBMS must spill data to disk, then we need to be smarter…

EXTERNAL HASHING AGGREGATE

Phase #1 – Partition

→ Divide tuples into buckets based on hash key
→ Write them out to disk when they get full

Phase #2 – Rehash

→ Build in-memory hash table for each partition and compute the aggregation

PHASE #1 – PARTITION

Use a hash function h_1 to split tuples into **partitions** on disk.

→ A partition is one or more pages that contain the set of keys with the same hash value.

→ Partitions are "spilled" to disk via output buffers.

Assume that we have ***B*** buffers.

We will use ***B-1*** buffers for the partitions and ***1*** buffer for the input data.

PHASE #1 – PARTITION

```
SELECT DISTINCT cid
  FROM enrolled
 WHERE grade IN ('B','C')
```

enrolled(sid,cid,grade)

sid	cid	grade
53666	15-445	C
53688	15-721	A
53688	15-826	B
53666	15-721	C
53655	15-445	C

⋮

Filter

sid	cid	grade
53666	15-445	C
53688	15-826	B
53666	15-721	C
53655	15-445	C
53672	15-210	B

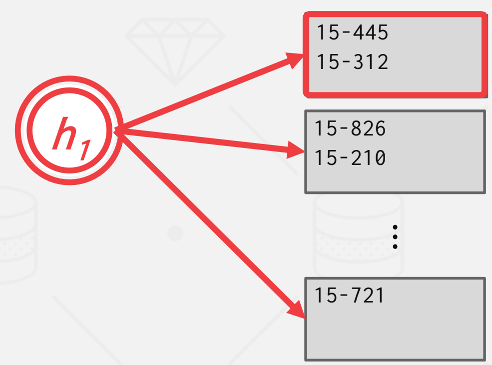
⋮

Remove Columns

cid
15-445
15-826
15-721
15-445
15-210

⋮

B-1 partitions



PHASE #2 – REHASH

For each partition on disk:
→ Read it into memory and build an in-memory hash table based on a different hash function h_2.
→ Then go through each bucket of this hash table to bring together matching tuples.

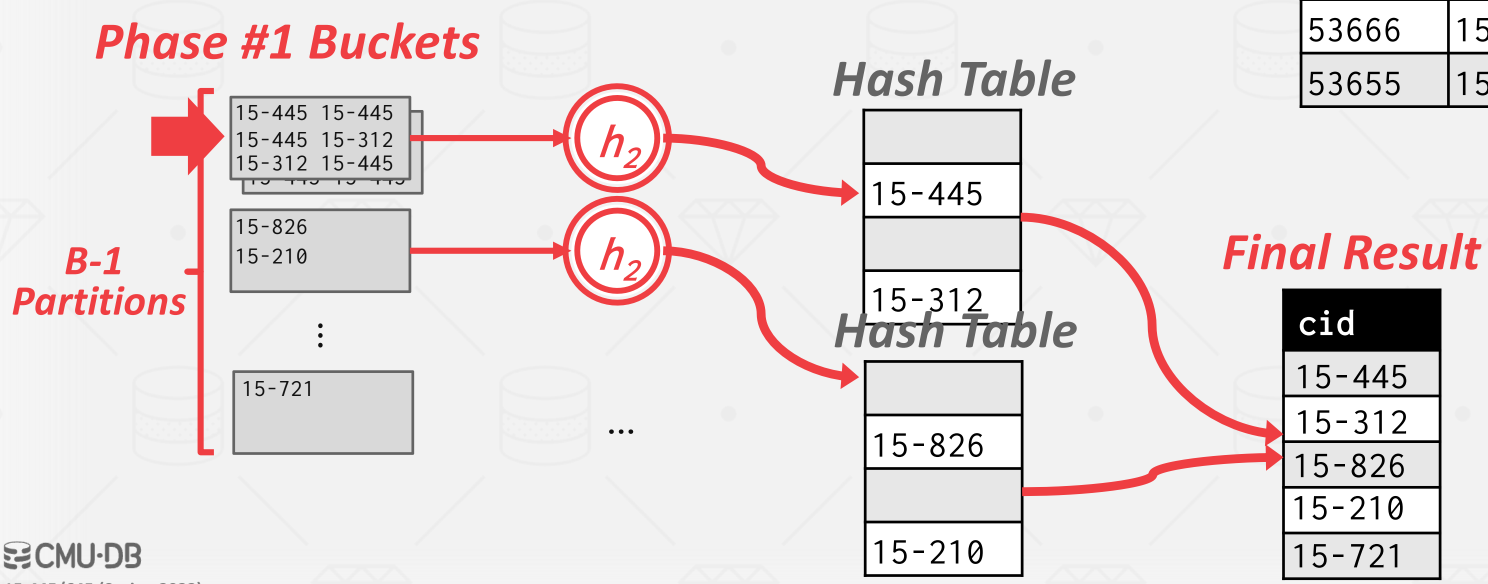
This assumes that each partition fits in memory.

PHASE #2 - REHASH

```
SELECT DISTINCT cid
  FROM enrolled
 WHERE grade IN ('B','C')
```

enrolled(sid,cid,grade)

sid	cid	grade
53666	15-445	C
53688	15-721	A
53688	15-826	B
53666	15-721	C
53655	15-445	C



HASHING SUMMARIZATION

During the rehash phase, store pairs of the form (`GroupKey→RunningVal`)

When we want to insert a new tuple into the hash table:

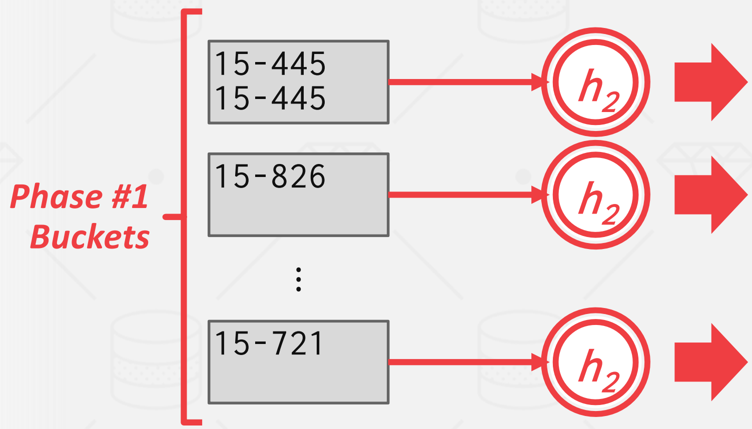
→ If we find a matching `GroupKey`, just update the `RunningVal` appropriately
→ Else insert a new `GroupKey→RunningVal`

HASHING SUMMARIZATION

```
SELECT cid, AVG(s.gpa)
  FROM student AS s, enrolled AS e
 WHERE s.sid = e.sid
 GROUP BY cid
```

Running Totals

AVG(col) → (COUNT,SUM)
MIN(col) → (MIN)
MAX(col) → (MAX)
SUM(col) → (SUM)
COUNT(col) → (COUNT)



Hash Table

key	value
15-445	(2, 7.32)
15-826	(1, 3.33)
15-721	(1, 2.89)

Final Result

cid	AVG(gpa)
15-445	3.66
15-826	3.33
15-721	2.89

CONCLUSION

Choice of sorting vs. hashing is subtle and depends on optimizations done in each case.

We already discussed the optimizations for sorting:
→ Chunk I/O into large blocks to amortize costs
→ Double-buffering to overlap CPU and I/O

NEXT CLASS

Nested Loop Join
Sort-Merge Join
Hash Join