Carnegie Mellon University Database Systems Modern SQL

TODAY'S AGENDA

Database Systems Background

Relational Model

Relational Algebra

Alternative Data Models

Q&A Session



LAST CLASS

We introduced the Relational Model as the superior data model for databases.

We then showed how Relational Algebra is the building blocks that will allow us to query and modify a relational database.



In 1971, IBM created its first relational query language called <u>SQUARE</u>.

IBM then created "SEQUEL" in 1972 for <u>IBM</u>

<u>System R</u> prototype DBMS.

 \rightarrow Structured English Query Language

IBM releases commercial SQL-based DBMSs:

→ System/38 (1979), SQL/DS (1981), and DB2 (1983).



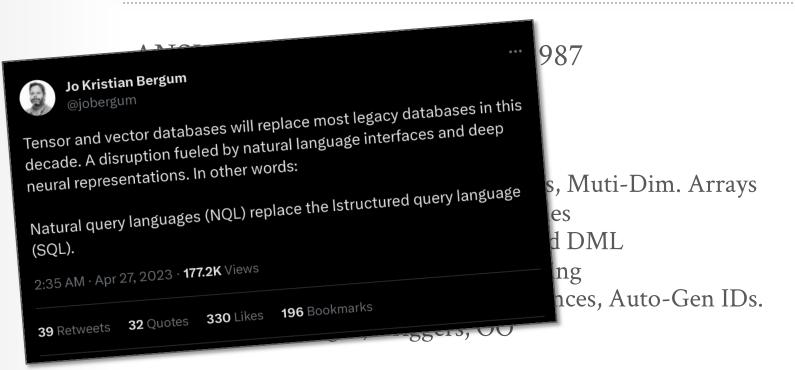
- ANSI Standard in 1986. ISO in 1987
- \rightarrow <u>S</u>tructured <u>Q</u>uery <u>L</u>anguage

Current standard is **SQL:2023**

- → **SQL:2023** → Property Graph Queries, Muti-Dim. Arrays
- \rightarrow **SQL:2016** \rightarrow JSON, Polymorphic tables
- → **SQL:2011** → Temporal DBs, Pipelined DML
- → **SQL:2008** → Truncation, Fancy Sorting
- \rightarrow **SQL:2003** \rightarrow XML, Windows, Sequences, Auto-Gen IDs.
- \rightarrow **SQL:1999** \rightarrow Regex, Triggers, OO

The minimum language syntax a system needs to say that it supports SQL is **SQL-92**.





The minimum language syntax a system needs to say that it supports SQL is **SQL-92**.



Jo Kristian Bergum @jobergum

Tensor and vector databases will replace most legar decade. A disruption fueled by natural language int neural representations. In other words:

Natural query languages (NQL) replace the Istruct (SQL).

2:35 AM · Apr 27, 2023 · **177.2K** Views

196 Bookmark 330 Likes 39 Retweets 32 Quotes

> The minimum lan say that it support



Gagan Biyani 🏛 🤣



@gaganbiyani

SQL is going to die at the hands of an Al. I'm serious.

@mayowaoshin is already doing this. Takes your company's data and ingests it into ChatGPT. Then, you can create a chatbot for the data and just ask it questions using natural language.

This video demoes the output.



10:30 AM · May 18, 2023 · **2.6M** Views

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NEWS | COMPUTING

The Rise of SQL > It's become the second programming language everyone needs to know

BY RINA DIANE CABALLAR | 23 AUG 2022 | 3 MIN READ | \[\bigcap ISTOCK

Tensor and vector day decade. A disruption neural representatio

@jobergum

Jo Kristian Bergu

Natural query langu (SQL).

32 Q **39** Retweets

sai

SHARE THIS STORY ☑ 8 ¥ f in TAGS TOP PROGRAMMING LANGUAGES

SQL dominated the jobs ranking in *IEEE Spectrum*'s interactive rankings of the top programming languages this year. Normally, the top position is occupied by Python or other mainstays, such as C, C++, Java, and JavaScript, but the sheer number of times employers said they wanted developers with SQL skills, albeit in addition to a more general-purpose language, boosted it to No. 1.

So what's behind SQL's soar to the top? The ever-increasing use of databases, for one. SQL has become the primary query language for accessing and managing data stored in such databases—specifically relational databases, which represent data in table form with rows and columns. Databases serve as the foundation of many enterprise applications and are increasingly found in other places as well, for example taking the place of traditional file systems in smartphones.

"This ubiquity means that every software developer will have to interact with databases no matter the field, and SQL is the de facto standard for interacting with databases," says Andy Pavlo, a professor specializing in database management at the Carnegie Mellon University (CMU) School of Computer Science and a member of the CMU database group.

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company's data and chatbot for the data and

RELATIONAL LANGUAGES

Data Manipulation Language (DML)
Data Definition Language (DDL)
Data Control Language (DCL)

Also includes:

- → View definition
- → Integrity & Referential Constraints
- → Transactions

Important: SQL is based on **bags** (duplicates) not **sets** (no duplicates).



TODAY'S AGENDA

Aggregations + Group By

String / Date / Time Operations

Output Control + Redirection

Window Functions

Nested Queries

Lateral Joins

Common Table Expressions



EXAMPLE DATABASE

student(sid,name,login,gpa)

sid	name	login	age	gpa
53666	RZA	rza@cs	55	4.0
53688	Taylor	swift@cs	27	3.9
53655	Tupac	shakur@cs	25	3.5

course(cid, name)

cid	name	
15-445	Database Systems	
15-721	Advanced Database Systems	
15-826	Data Mining	
15-799	Special Topics in Databases	

enrolled(sid,cid,grade)

sid cid		grade
53666	15-445	С
53688	15-721	Α
53688	15-826	В
53655	15-445	В
53666	15-721	С

Functions that return a single value from a bag of tuples:

- \rightarrow **AVG(col)** \rightarrow Return the average col value.
- \rightarrow MIN(col) \rightarrow Return minimum col value.
- \rightarrow MAX(col) \rightarrow Return maximum col value.
- \rightarrow **SUM(col)** \rightarrow Return sum of values in col.
- \rightarrow **COUNT(col)** \rightarrow Return # of values for col.



Aggregate functions can (almost) only be used in the **SELECT** output list.

```
SELECT COUNT(login) AS cnt
FROM student WHERE login LIKE '%@cs'
```



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```



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```
SELECT COUNT(*) AS cnt

SELECT COUNT(*) AS cnt

FROM student WHERE login LIKE '%@cs'
```



Aggregate functions can (almost) only be used in the **SELECT** output list.

```
SELECT COUNT(*) AS cnt

SELECT COUNT(1) AS cnt

FROM student WHERE login LIKE '%@cs'
```



Aggregate functions can (almost) only be used in the **SELECT** output list.

```
SELECT COUNT(*) AS cnt

SELECT COUNT(1) AS cnt

SELECT COUNT(1+1+1) AS cnt

FROM student WHERE login LIKE '%@cs'
```



MULTIPLE AGGREGATES

Get the number of students and their average GPA that have a "@cs" login.

	AVG(gpa)	COUNT(sid)	1
SELECT AVG(gpa), COUNT(sid)	3.8	3]
FROM student WHERE login LIKE	'%@cs'		



Output of other columns outside of an aggregate is undefined.

Get the average GPA of students enrolled in each course.

	AVG(s.gpa)	e.cid
SELECT AVG(s.gpa), e.cid	3.86	???
FROM enrolled AS e JOIN student AS s		
ON e.sid = s.sid		

Output of other columns outside of an aggregate is undefined.

Get the average GPA of students enrolled in each course.

		AVG(s.gpa)	e.cid
SELECT	AVG(s.gpa), e d	3.86	???
FROM	enrolled AS JON student AS s		
ON	e.sid = s.sid		



Output of other columns outside of an aggregate is undefined.

Get the average GPA of students enrolled in each course.

		AVG(s.g	pa) e.cid
SELECT	AVG(s.gpa), e d	3.86	???
FROM	enrolled AS JON student AS s		
ON	e.sid = s.sid		

		AVG(s.gpa)	e.cid
SELECT	AVG(s.gpa), ANY_VALUE(e.cid)	3.86	15-445
FROM	enrolled AS e JOIN student AS s		
ON	ON e.sid = s.sid		



Project tuples into subsets and calculate aggregates against each subset.

```
SELECT AVG(s.gpa), e.cid
  FROM enrolled AS e JOIN student AS s
   ON e.sid = s.sid
GROUP BY e.cid
```

e.sid	s.sid	s.gpa	e.cid
53435	53435	2.25	15-721
53439	53439	2.70	15-721
56023	56023	2.75	15-826
59439	59439	3.90	15-826
53961	53961	3.50	15-826
58345	58345	1.89	15-445



Project tuples into subsets and calculate aggregates against each subset.

SELECT AVG(s.gpa), e.cid
 FROM enrolled AS e JOIN student AS s
 ON e.sid = s.sid
GROUP BY e.cid

e.sid	s.sid	s.gpa	e.cid
53435	53435	2.25	15-721
53439	53439	2.70	15-721
56023	56023	2.75	15-826
59439	59439	3.90	15-826
53961	53961	3.50	15-826
58345	58345	1.89	15-445



Project tuples into subsets and calculate aggregates against each subset.

SELECT AVG(s.gpa), e.cid
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e.sid	s.sid	s.gpa	e.cid
53435	53435	2.25	15-721
53439	53439	2.70	15-721
56023	56023	2.75	15-826
59439	59439	3.90	15-826
53961	53961	3.50	15-826
58345	58345	1.89	15-445



AVG(s.gpa)	e.cid
2.46	15-721
3.39	15-826
1.89	15-445



Non-aggregated values in **SELECT** output clause must appear in **GROUP BY** clause.

```
SELECT AVG(s.gpa), e.cid, s.name
  FROM enrolled AS e, student AS s
WHERE e.sid = s.sid
GROUP BY e.cid
```



Non-aggregated values in **SELECT** output clause must appear in **GROUP BY** clause.

```
SELECT AVG(s.gpa), e.cid, s. te

FROM enrolled AS e, student A s
WHERE e.sid = s.sid
GROUP BY e.cid
```



Non-aggregated values in **SELECT** output clause must appear in **GROUP BY** clause.

```
SELECT AVG(s.gpa), e.cid, s.name
  FROM enrolled AS e JOIN student AS s
  ON e.sid = s.sid
  GROUP BY e.cid, s.name
```



Filters results based on aggregation computation.

```
SELECT AVG(s.gpa) AS avg_gpa, e.cid
  FROM enrolled AS e, student AS s
WHERE e.sid = s.sid
  AND avg_gpa > 3.9
GROUP BY e.cid
```



Filters results based on aggregation computation.

```
SELECT AVG(s.gpa) AS avg_gpa, e.cid
  FROM enrolled AS e, student AS s
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  AND avg_gpa > 3.9
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SELECT AVG(s.gpa) AS avg_gpa, e.cid
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WHERE e.sid = s.sid
GROUP BY e.cid
HAVING avg_gpa > 3.9;
```



Filters results based on aggregation computation.

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```
SELECT AVG(s.gpa) AS avg_gpa, e.cid
  FROM enrolled AS e, student AS s
WHERE e.sid = s.sid
GROUP BY e.cid
HAVING AVG(s.gpa) > 3.9;
```

AVG(s.gpa)	e.cid
3.75	15-415
3.950000	15-721
3.900000	15-826



avg_gpa	e.cid
3.950000	15-721



STRING OPERATIONS

	String Case	String Quotes
SQL-92	Sensitive	Single Only
Postgres	Sensitive	Single Only
MySQL	Insensitive	Single/Double
SQLite	Sensitive	Single/Double
MSSQL	Sensitive	Single Only
Oracle	Sensitive	Single Only
<pre>WHERE UPPER(name) = UPPER('TuPaC')</pre> <pre>SQL-92</pre>		
WHERE name = "TuPaC"		



STRING OPERATIONS

LIKE is used for string matching.

String-matching operators

- → '%' Matches any substring (including empty strings).
- → '_' Match any one character

```
SELECT * FROM enrolled AS e WHERE e.cid LIKE '15-%'
```

```
SELECT * FROM student AS s
WHERE s.login LIKE '%@c_'
```

STRING OPERATIONS

SQL-92 defines string functions.

→ Many DBMSs also have their own unique functions

Can be used in either output and predicates:

```
SELECT SUBSTRING(name,1,5) AS abbrv_name
FROM student WHERE sid = 53688
```

```
SELECT * FROM student AS s
WHERE UPPER(s.name) LIKE 'KAN%'
```



STRING OPERATIONS

SQL standard defines the | operator for concatenating two or more strings together.

```
SELECT name FROM student
WHERE login = LOWER(name) || '@cs'

SELECT name FROM student
WHERE login = LOWER(name) + '@cs'

SELECT name FROM student
WHERE login = CONCAT(LOWER(name), '@cs')
```



DATE/TIME OPERATIONS

Operations to manipulate and modify **DATE/TIME** attributes.

Can be used in both output and predicates.

Support/syntax varies wildly...

Demo: Get the # of days since the beginning of the year.



OUTPUT REDIRECTION

Store query results in another table:

- \rightarrow Table must not already be defined.
- → Table will have the same # of columns with the same types as the input.

```
SELECT DISTINCT cid INTO CourseIds
FROM enrolled;

SELECT DISTINCT cid
INTO TEMPORARY CourseIds
FROM enrolled;

CREATE TABLE CourseIds (
SELECT DISTINCT cid FROM enrolled);
```



OUTPUT REDIRECTION

Insert tuples from query into another table:

- → Inner **SELECT** must generate the same columns as the target table.
- → DBMSs have different options/syntax on what to do with integrity violations (e.g., invalid duplicates).

```
INSERT INTO CourseIds
(SELECT DISTINCT cid FROM enrolled);
```

ORDER BY <column*> [ASC|DESC]

CELECT aid grade EDOM ennelled	sid	grade
	53123	Α
WHERE cid = '15-721'	53334	Α
ORDER BY grade	53650	В
Office of the second of the se	53666	D

ORDER BY <column*> [ASC|DESC]

```
SELECT sid, grade FROM enrolled

WH SELECT sid, grade FROM enrolled

OF WHERE cid = '15-721'

ORDER BY 2
```

ORDER BY <column*> [ASC|DESC]

```
SELECT sid, grade FROM enrolled

WH SELECT sid, grade FROM enrolled

OF WHERE cid = '15-721'

ORDER BY 2
```

```
SELECT sid FROM enrolled
WHERE cid = '15-721'
ORDER BY grade DESC, sid ASC
```



ORDER BY <column*> [ASC|DESC]

```
SELECT sid, grade FROM enrolled

WH SELECT sid, grade FROM enrolled

OF WHERE cid = '15-721'

ORDER BY 2
```

```
SELECT sid FROM enrolled

WHERE cid = '15-721'

ORDER BY grade DESC, sid ASC

53650

53123

53334
```



FETCH {FIRST|NEXT} <count> ROWS OFFSET <count> ROWS

- → Limit the # of tuples returned in output.
- → Can set an offset to return a "range"

```
SELECT sid, name FROM student
WHERE login LIKE '%@cs'
FETCH FIRST 10 ROWS ONLY;
```



FETCH {FIRST|NEXT} <count> ROWS OFFSET <count> ROWS

- \rightarrow Limit the # of tuples returned in output.
- → Can set an offset to return a "range"

```
SELECT sid, name FROM student
WHERE login LIKE '%@cs'
FETCH FIRST 10 ROWS ONLY;
```

```
SELECT sid, name FROM student
WHERE login LIKE '%@cs'
ORDER BY gpa
OFFSET 10 ROWS
FETCH FIRST 10 ROWS WITH TIES;
```



FETCH {FIRST|NEXT} <count> ROWS OFFSET <count> ROWS

- → Limit the # of tuples returned in output.
- → Can set an offset to return a "range"

```
SELECT sid, name FROM student
WHERE login LIKE '%@cs'
FETCH FIRST 10 ROWS ONLY;
```

```
SELECT sid, name FROM student
WHERE login LIKE '%@cs'
ORDER BY gpa
OFFSET 10 ROWS
FETCH FIRST 10 ROWS WITH TIES;
```

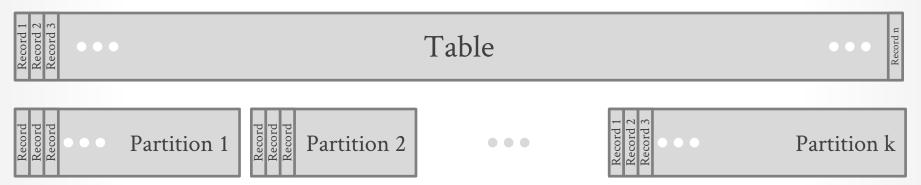


Conceptual execution: (optional) Partition data → (optional) sort each partition

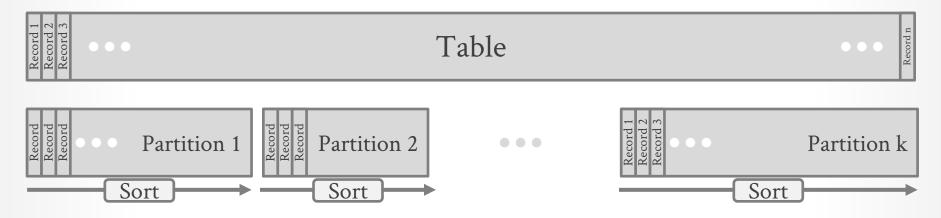
Conceptual execution: (optional) Partition data → (optional) sort each partition



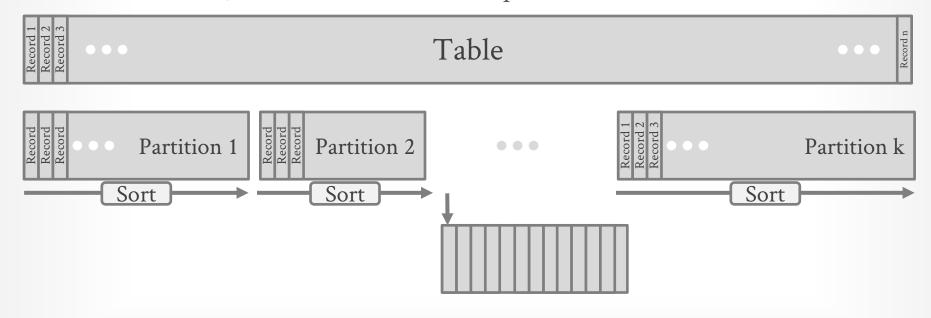
Conceptual execution: (optional) Partition data → (optional) sort each partition



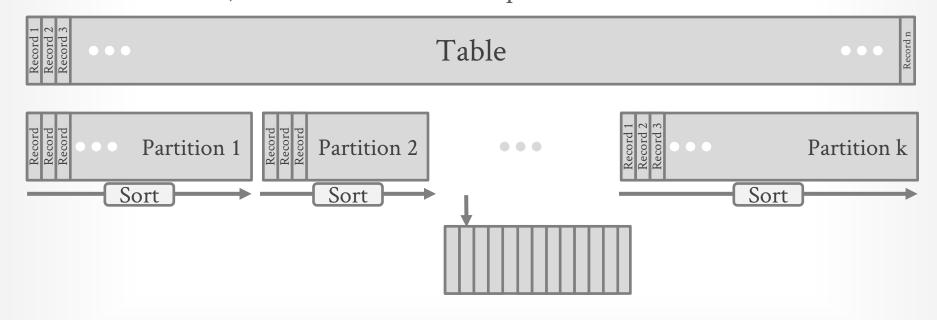




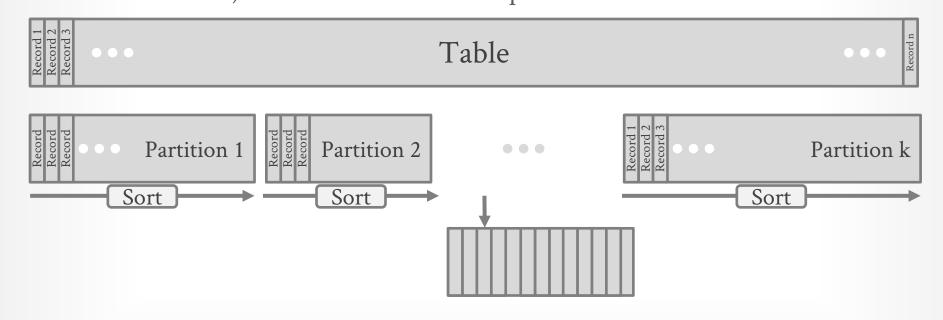






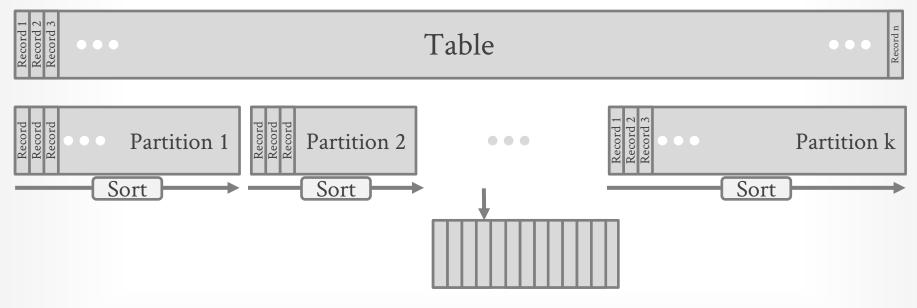




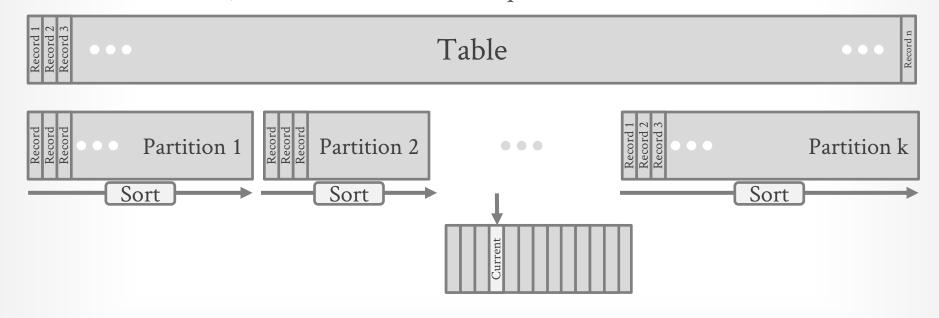




Conceptual execution: (optional) Partition data → (optional) sort each partition

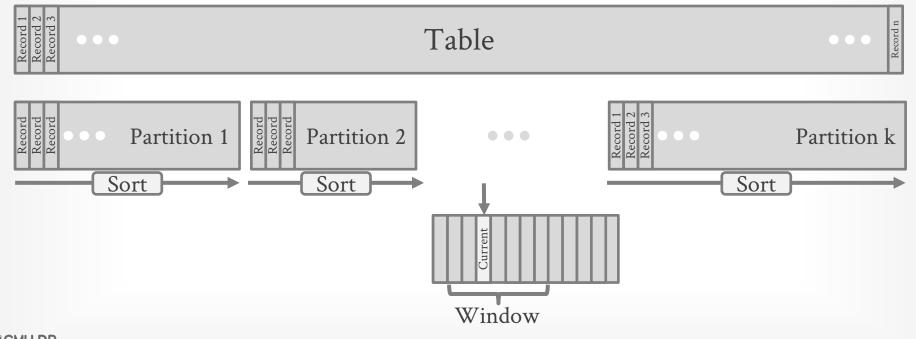






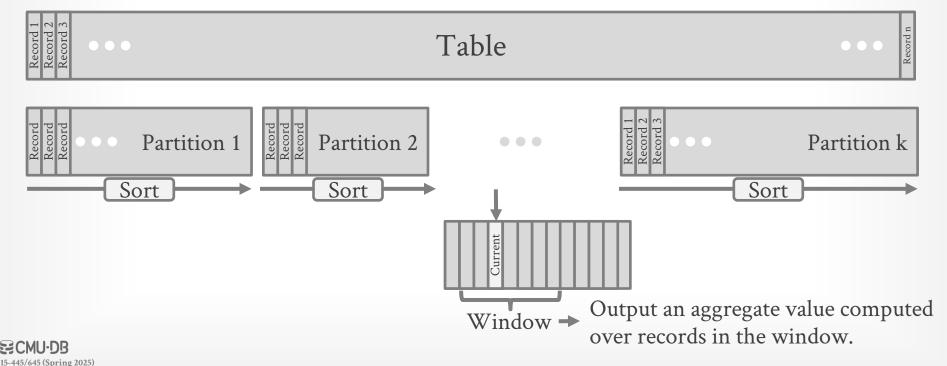


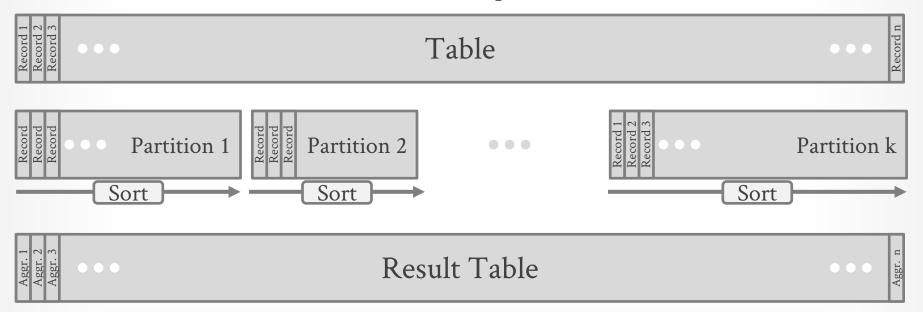
Conceptual execution: (optional) Partition data → (optional) sort each partition





Conceptual execution: (optional) Partition data → (optional) sort each partition







Performs a calculation across a set of tuples that are related to the current tuple, without collapsing them into a single output tuple, to support running totals, ranks, and moving averages.

→ Like an aggregation but tuples are not grouped into a single output tuples.

```
SELECT FUNC-NAME(...) OVER (...)
FROM tableName
```



Performs a calculation across a set of tuples that are related to the current tuple, without collapsing them into a single output tuple, to support running totals, ranks, and moving averages.

→ Like an aggregation but tuples are not grouped into a single output tuples.

How to "slice" up data
Can also sort tuples

```
SELECT FUNC-NAME(...) OVER (...)
FROM tableName
```

Aggregation Functions
Special Functions



Aggregation functions:

→ Anything that we discussed earlier

Special window functions:

- \rightarrow **ROW_NUMBER()** \rightarrow # of the current row
- → RANK() → Order position of the current row.

```
SELECT *, ROW_NUMBER() OVER () AS row_num
FROM enrolled
```



Aggregation functions:

→ Anything that we discussed earlier

Special window functions:

- \rightarrow **ROW_NUMBER()** \rightarrow # of the current row
- → RANK()→ Order position of the current row.

sid	cid	grade	row_num
53666	15-445	С	1
53688	15-721	Α	2
53688	15-826	В	3
53655	15-445	В	4
53666	15-721	С	5

```
SELECT *, ROW_NUMBER() OVER () AS row_num
FROM enrolled
```



The **OVER** keyword specifies how to group together tuples when computing the window function.

Use **PARTITION BY** to specify group.

```
SELECT cid, sid,
ROW_NUMBER() OVER (PARTITION BY cid)
FROM enrolled
ORDER BY cid
```



The **OVER** keyword specifies how to group together tuples when computing the window function.

Use **PARTITION BY** to specify group.

cid	sid	row_number
15-445	53666	1
15-445	53655	2
15-721	53688	1
15-721	53666	2
15-826	53688	1

```
SELECT cid, sid,
ROW_NUMBER() OVER (PARTITION BY cid)
FROM enrolled
ORDER BY cid
```



The **OVER** keyword specifies how to group together tuples when computing the window function.

Use **PARTITION BY** to specify group.

cid	sid	row_number
15-445	53666	1
15-445	53655	2
15-721	53688	1
15-721	53666	2
15-826	53688	1

```
SELECT cid, sid,
ROW_NUMBER() OVER (PARTITION BY cid)
FROM enrolled
ORDER BY cid
```



You can also include an **ORDER BY** in the window grouping to sort entries in each group.

```
SELECT *,

ROW_NUMBER() OVER (ORDER BY cid)

FROM enrolled
ORDER BY cid
```



Find the student with the <u>second</u> highest grade for each course.

```
SELECT * FROM (
    SELECT *, RANK() OVER (PARTITION BY cid
          ORDER BY grade ASC) AS rank
    FROM enrolled) AS ranking
WHERE ranking.rank = 2
```



Find the student with the <u>second</u> highest grade for each course.

Group tuples by cid
Then sort by grade

```
SELECT * FROM (
    SELECT *, RANK() OVER (PARTITION BY cid
          ORDER BY grade ASC) AS rank
    FROM enrolled) AS ranking
WHERE ranking.rank = 2
```



Find the student with the <u>second</u> highest grade for each course.

Group tuples by cid Then sort by grade

```
SELECT * FROM (
SELECT *, RANK() OVER (PARTITION BY cid
ORDER BY grade ASC) AS rank
FROM enrolled) AS ranking
WHERE ranking.rank = 2
```



NESTED QUERIES

Invoke a query inside of another query to compose more complex computations.

→ Inner queries can appear (almost) anywhere in query.

Outer Query

SELECT name FROM student WHERE

sid IN (SELECT sid FROM enrolled) Inner Query



Invoke a query inside of another query to compose more complex computations.

 \rightarrow Inner queries can appear (almost) anywhere in query.

```
SELECT name FROM student WHERE

sid IN (SELECT sid FROM enrolled)

SELECT sid,

(SELECT name FROM student AS s

WHERE s.sid = e.sid) AS name

FROM enrolled AS e;
```



Invoke a query inside of another query to compose more complex computations.

 \rightarrow Inner queries can appear (almost) anywhere in query.

```
SELECT name FROM student WHERE

sid IN (SELECT sid FROM enrolled)

SELECT sid,

(SELECT name FROM student AS s

WHERE s.sid = e.sid) AS name

FROM enrolled AS e;
```



```
SELECT name FROM student
WHERE ...

sid in the set of people that take 15-445
```



```
SELECT name FROM student
WHERE ...
SELECT sid FROM enrolled
WHERE cid = '15-445'
```



```
SELECT name FROM student
WHERE sid IN (
   SELECT sid FROM enrolled
   WHERE cid = '15-445'
)
```

```
SELECT name FROM student
WHERE sid in (
   SELECT sid FROM enrolled
   WHERE cid = '15-445'
)
```

ALL→ Must satisfy expression for all rows in the sub-query.

ANY→ Must satisfy expression for at least one row in the sub-query.

IN→ Equivalent to '=ANY()'.

EXISTS→ At least one row is returned without comparing it to an attribute in outer query.



```
SELECT name FROM student
WHERE sid = ANY(
    SELECT sid FROM enrolled
    WHERE cid = '15-445'
)
```



Find student record with the highest id that is enrolled in at least one course.

```
SELECT MAX(e.sid), s.name
  FROM enrolled AS e, student AS s
WHERE e.sid = s.sid;
```

This won't work in SQL-92. It runs in SQLite, but not Postgres or MySQL (v8 with strict mode).



Find student record with the highest id that is enrolled in at least one course.

```
SELECT sid, name FROM student
WHERE ...
```



Find student record with the highest id that is enrolled in at least one course.

```
SELECT sid, name FROM student
WHERE ...
```

"Is the highest enrolled sid"



Find student record with the highest id that is enrolled in at least one course.

```
SELECT sid, name FROM student
WHERE sid =
   SELECT MAX(sid) FROM enrolled
```



Find student record with the highest id that is enrolled in at least one course.

SELECT sid, name FROM student
WHERE sid =
SELECT MAX(sid) FROM enrolled

sid name
53688 Bieber



Find all courses that have no students enrolled in it.

```
SELECT * FROM course
WHERE ...
```

"with no tuples in the enrolled table"

cid	name
15-445	Database Systems
15-721	Advanced Database Systems
15-826	Data Mining
15-799	Special Topics in Databases

sid	cid	grade
53666	15-445	С
53688	15-721	Α
53688	15-826	В
53655	15-445	В
53666	15-721	С



Find all courses that have no students enrolled in it.

```
SELECT * FROM course
WHERE NOT EXISTS(
tuples in the enrolled table
)
```



Find all courses that have no students enrolled in it.

```
SELECT * FROM course
WHERE NOT EXISTS(
    SELECT * FROM enrolled
    WHERE course.cid = enrolled.cid
)
```

cid	name
15-799	Special Topics in Databases



Find all courses that have no students enrolled in it.

```
SELECT * FROM course
WHERE NOT EXISTS(
    SELECT * FROM enrolled
    WHERE course.cid = enrolled.cid
)
```

cid	name
15-799	Special Topics in Databases



The **LATERAL** operator allows a nested query to reference attributes in other nested queries that precede it.

→ You can think of it like a **for** loop that allows you to invoke another query for each tuple in a table.

```
SELECT * FROM

(SELECT 1 AS x) AS t1,

LATERAL (SELECT t1.x+1 AS y) AS t2;
```



Calculate the number of students enrolled in each course and the average GPA. Sort by enrollment count in descending order.

SELECT * **FROM** course **AS** c,

For each course:

→ Compute the # of enrolled students

For each course:

→ Compute the average gpa of enrolled students



```
SELECT * FROM course AS c,

LATERAL (SELECT COUNT(*) AS cnt FROM errolled

WHERE enrolled.cid = c.cid) S t1,

LATERAL (SELECT AVG(gpa) AS avg FROM student AS s

JOIN enrolled AS e ON s.sid = e.sid

WHERE e.cid = c.cid) AS t2;
```



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WHERE e.cid = c.cid) AS t2;
```





	cid	name	cnt	avg
	15-445	Database Systems	2	3.75
CELECT IN EDOM COLUMN AC			2	3.95
SELECT * FROM course AS c,	15-826	Data Mining	1	3.9
LATERAL (SELECT COUNT(*)	15-799	Special Topics in Databases	0	null
WHERE enrolled.cid = c.cid) AS t1,				
LATERAL (SELECT AVG(gpa)				
		ON s.sid = e.sid		
WHERE e.cid = c.cid AS t2;				



Specify a temporary result set that can then be referenced by another part of that query.

 \rightarrow Think of it like a temp table just for one query.

Alternative to nested queries, views, and explicit temp tables.

```
WITH cteName AS (
SELECT 1
)
SELECT * FROM cteName
```



Specify a temporary result set that can then be referenced by another part of that query.

 \rightarrow Think of it like a temp table just for one query.

Alternative to nested queries, views, and explicit temp tables.

```
WITH cteName AS (
SELECT 1
)
SELECT * FROM cteName
```



You can bind/alias output columns to names before the AS keyword.

```
WITH cteName (col1, col2) AS (
    SELECT 1, 2
)
SELECT col1 + col2 FROM cteName
```

```
WITH cteName (colXXX, colXXX) AS (         Postgres
         SELECT 1, 2
)
SELECT * FROM cteName
```



Find student record with the highest id that is enrolled in at least one course.

```
WITH cteSource (maxId) AS (
    SELECT MAX(sid) FROM enrolled
)
SELECT name FROM student, cteSource
WHERE student.sid = cteSource.maxId
```



OTHER THINGS TO NOTE

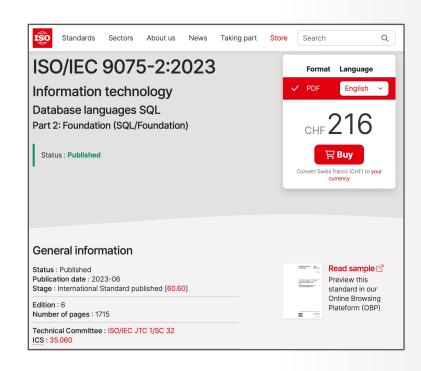
Identifiers (e.g. table and column names) are case-insensitive.

→ Makes it harder for applications that care about case (e.g., use CamelCased names).

One often sees quotes around names:

→ SELECT "ArtistList.firstName"

You have to pay cash money to get the standard documents.

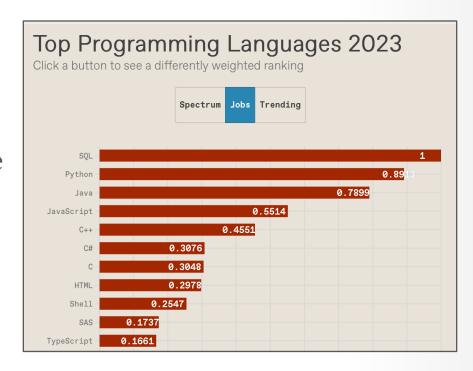


CONCLUSION

SQL is a hot language.

→ Lots of NL2SQL tools, but writing SQL is not going away.

You should (almost) always strive to compute your answer as a single SQL statement.



HOMEWORK #1

Write SQL queries to perform basic data analysis.

- → Write the queries locally using SQLite + DuckDB.
- → Submit them to Gradescope
- \rightarrow You can submit multiple times and use your best score.

Due: Wednesday Jan. 29th @ 11:59pm

https://15445.courses.cs.cmu.edu/spring2025/homework1



NEXT CLASS

We will begin our journey to understanding the internals of database systems starting with Storage!

